

High Throughput Secure MPC Over Small Population in Hybrid Networks

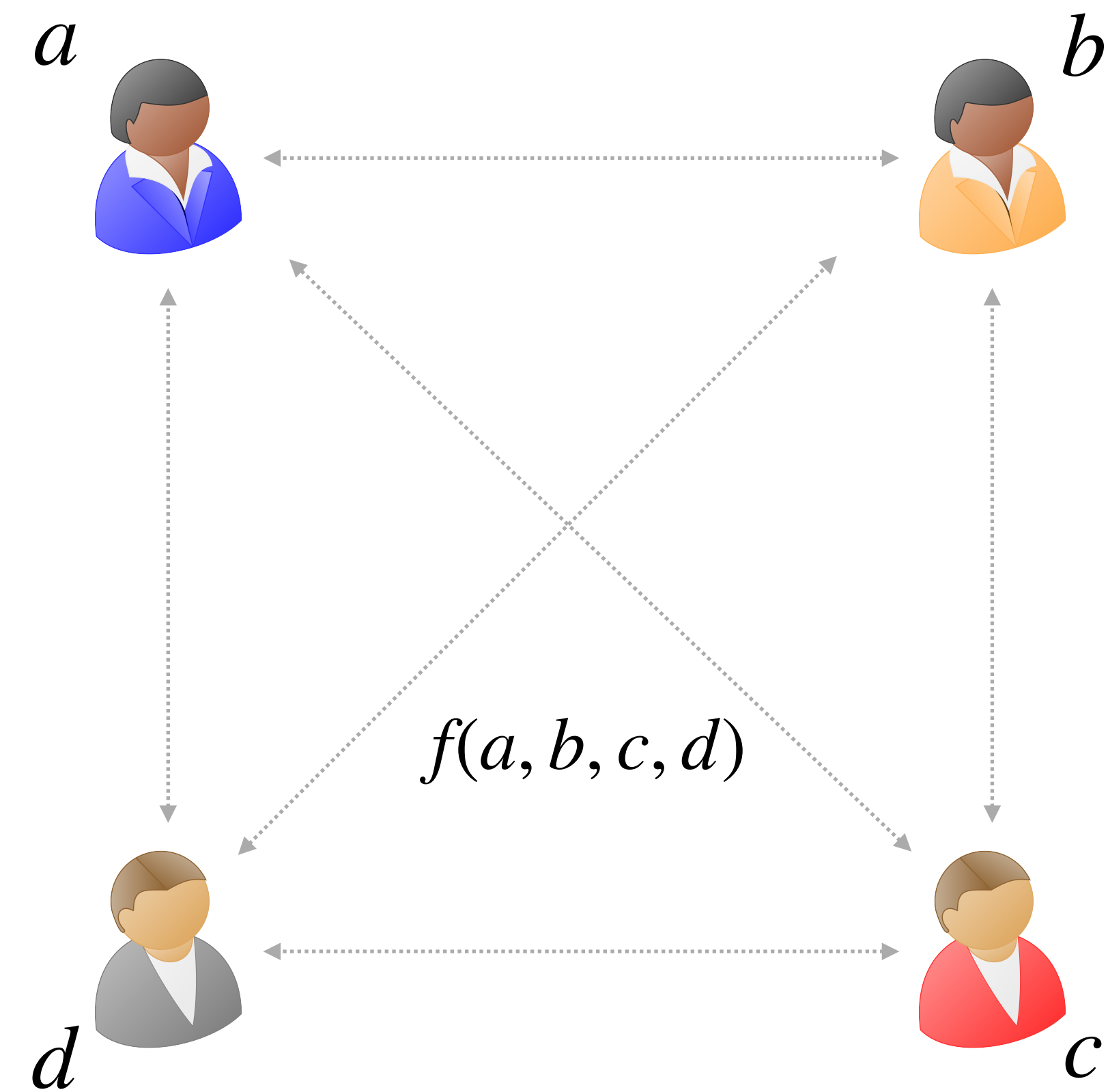
Ashish Choudhury, Aditya Hegde

INDOCRYPT 2020



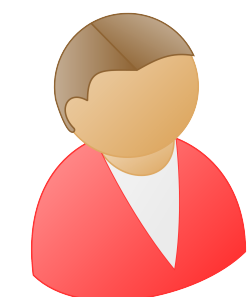
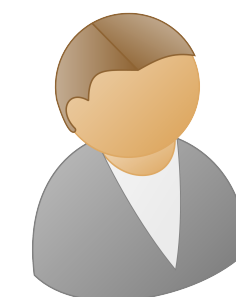
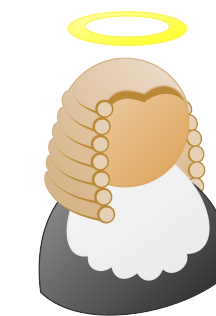
Secure Multi-Party Computation (MPC)

- **Distrusting** parties compute a function on **private** inputs



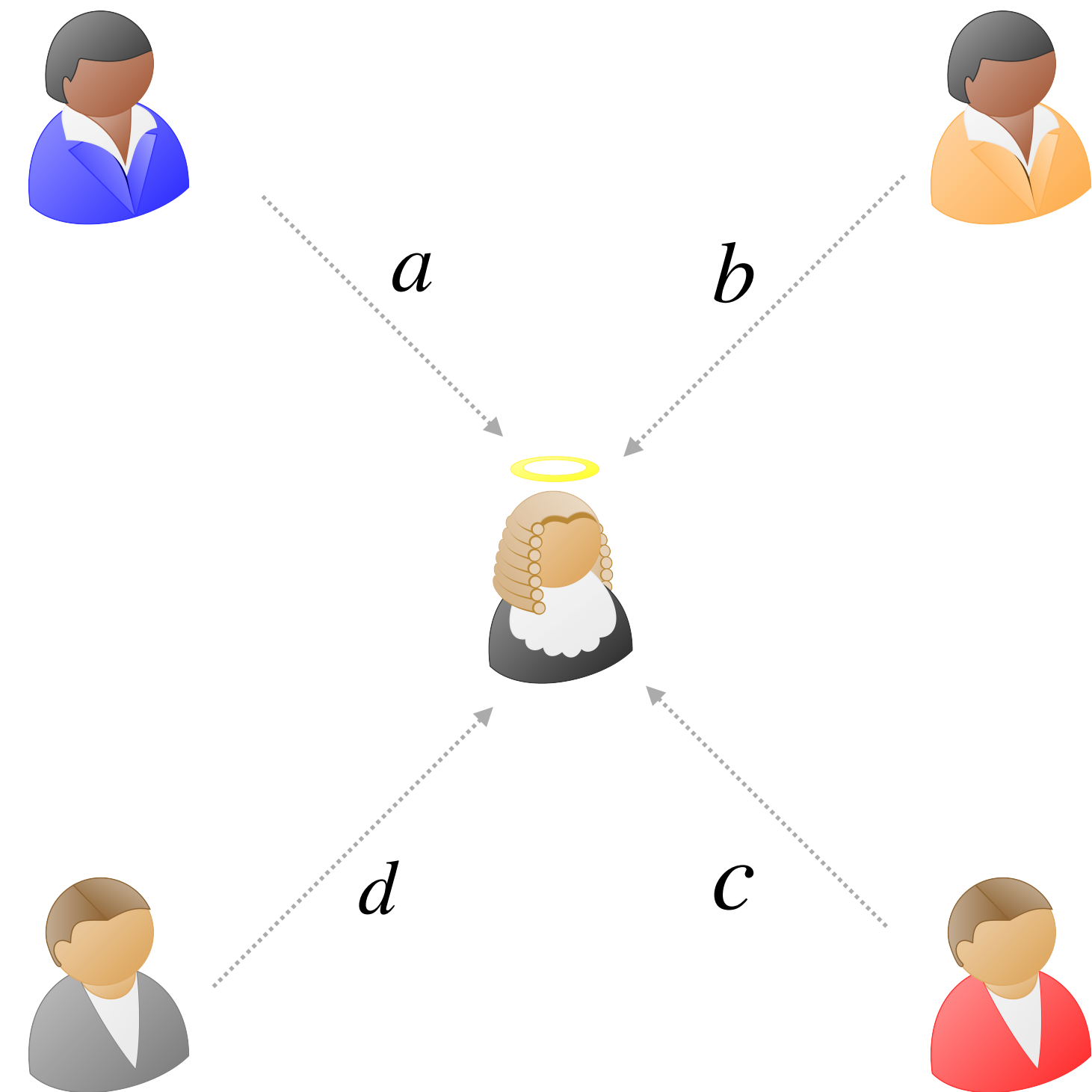
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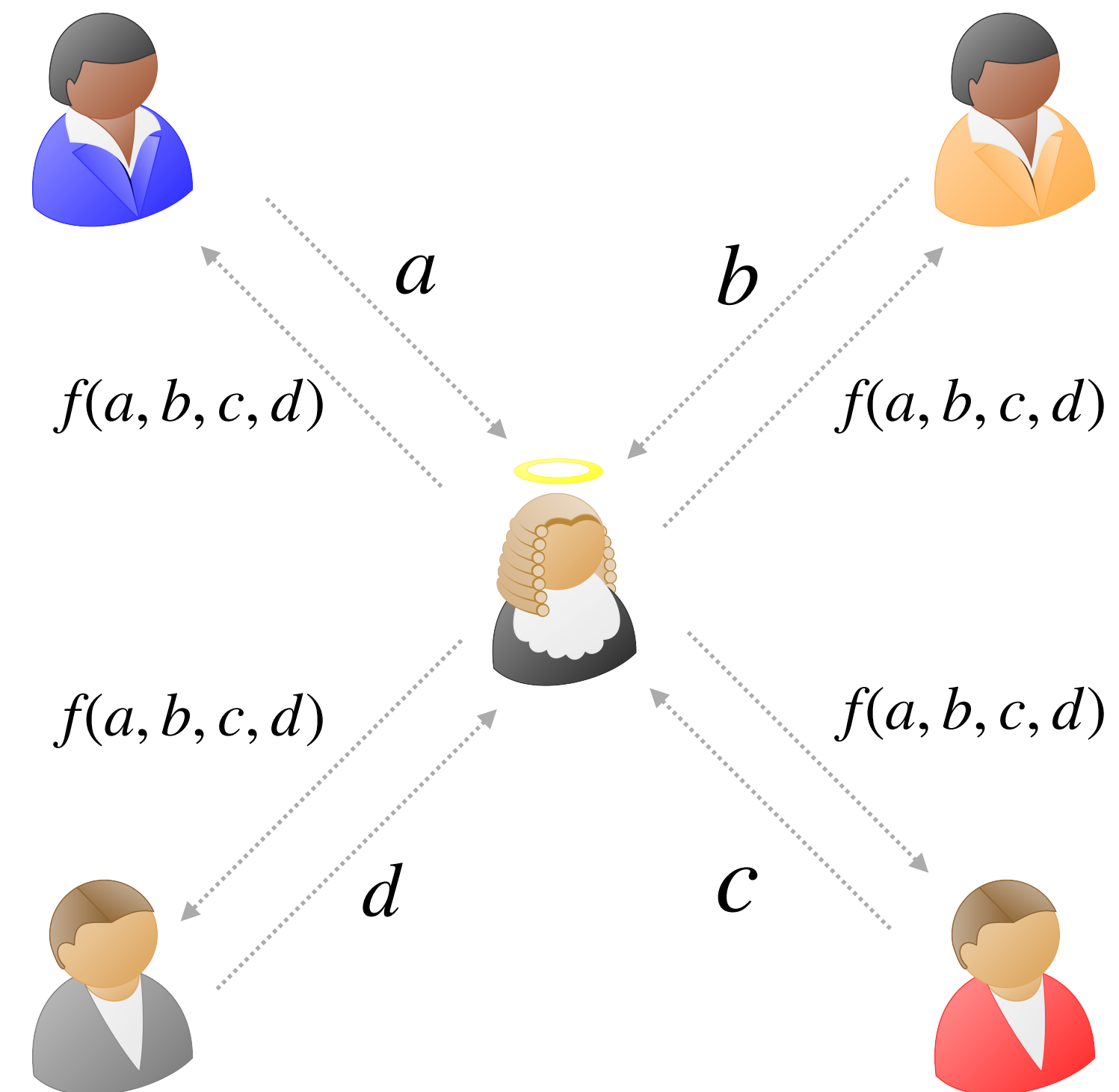
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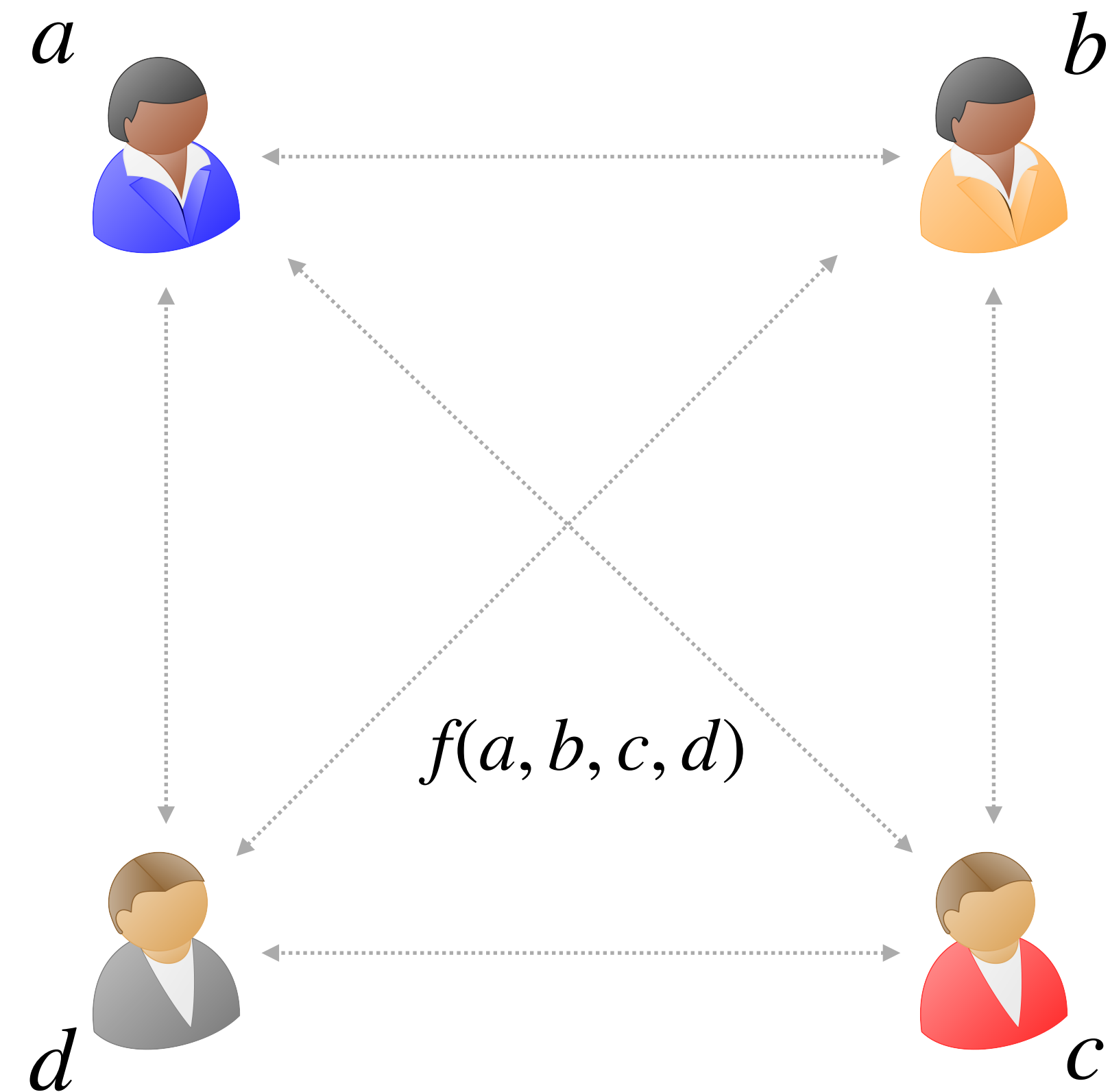
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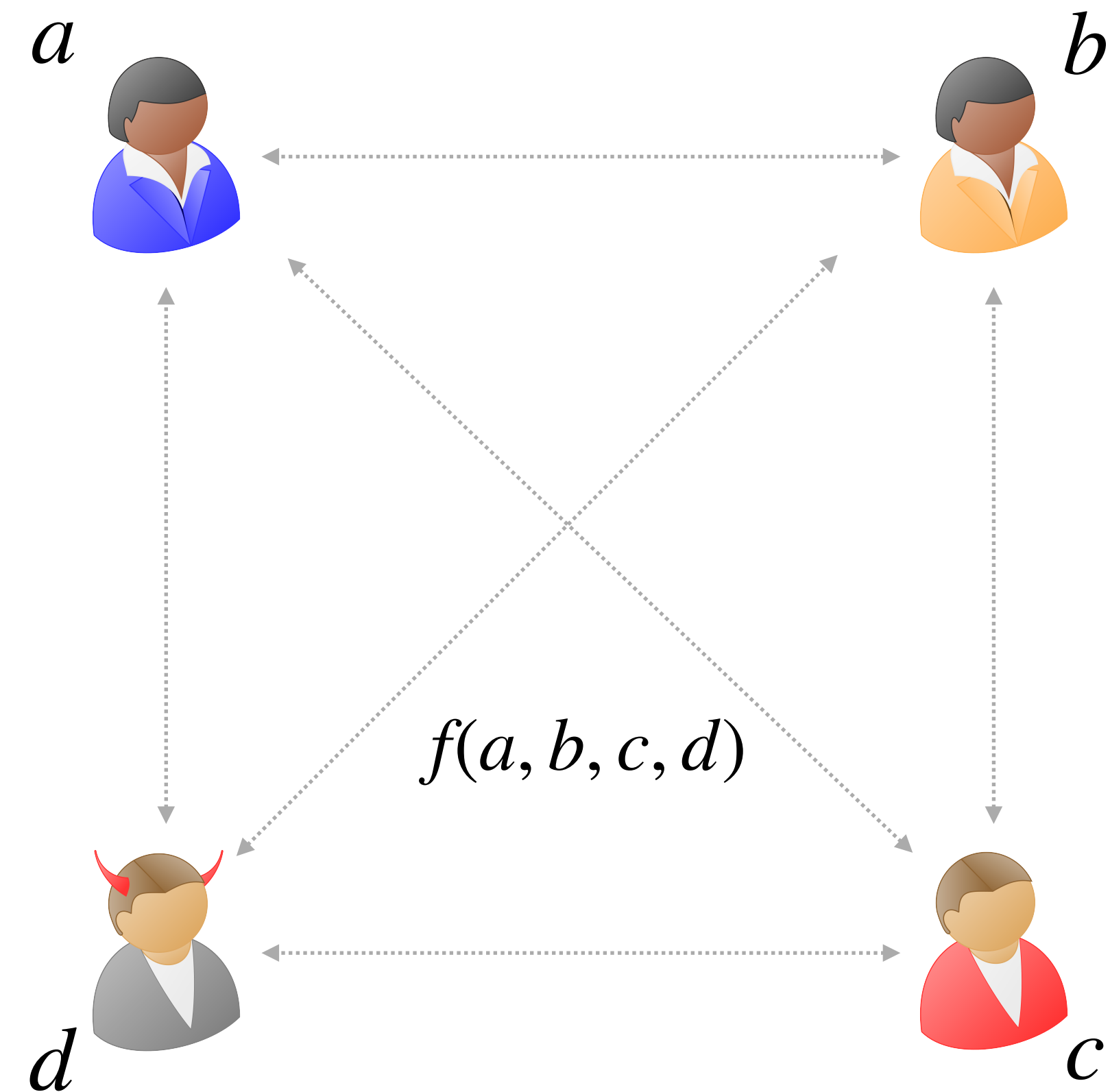
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- Setting
 - $n = 4, t = 1$
 - Malicious adversary

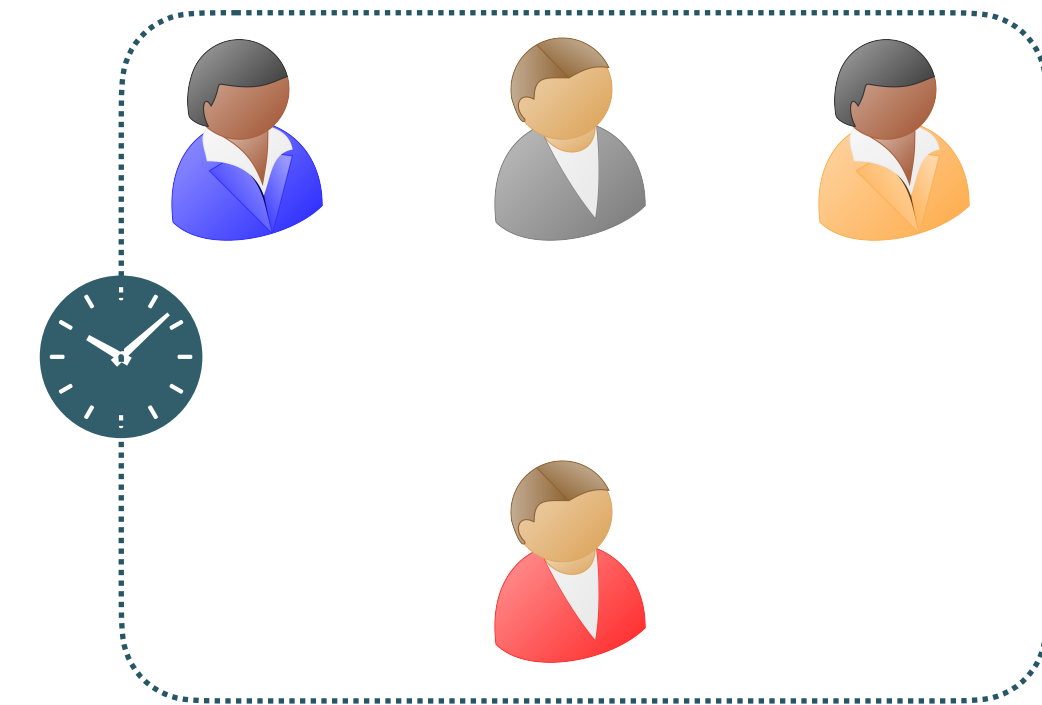


Communication Model - Synchronous and Asynchronous Networks

- Pairwise [private](#) and [authentic](#) channels

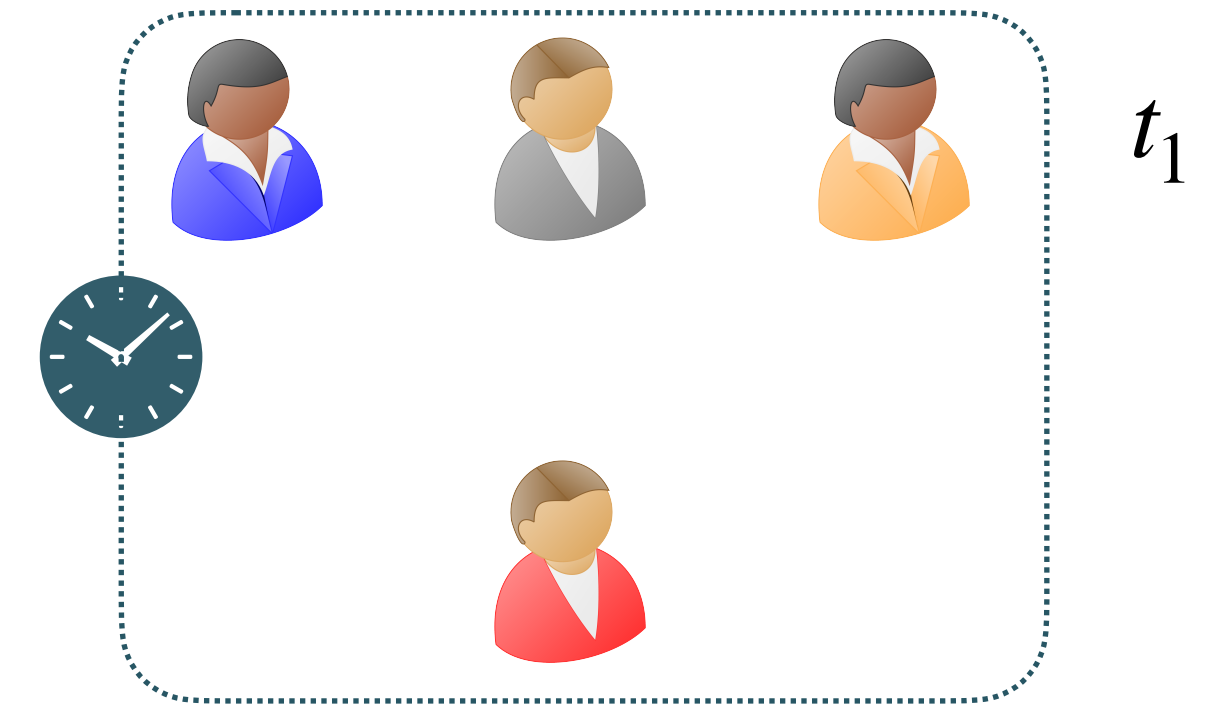
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 - Publicly known upper bound on message delay



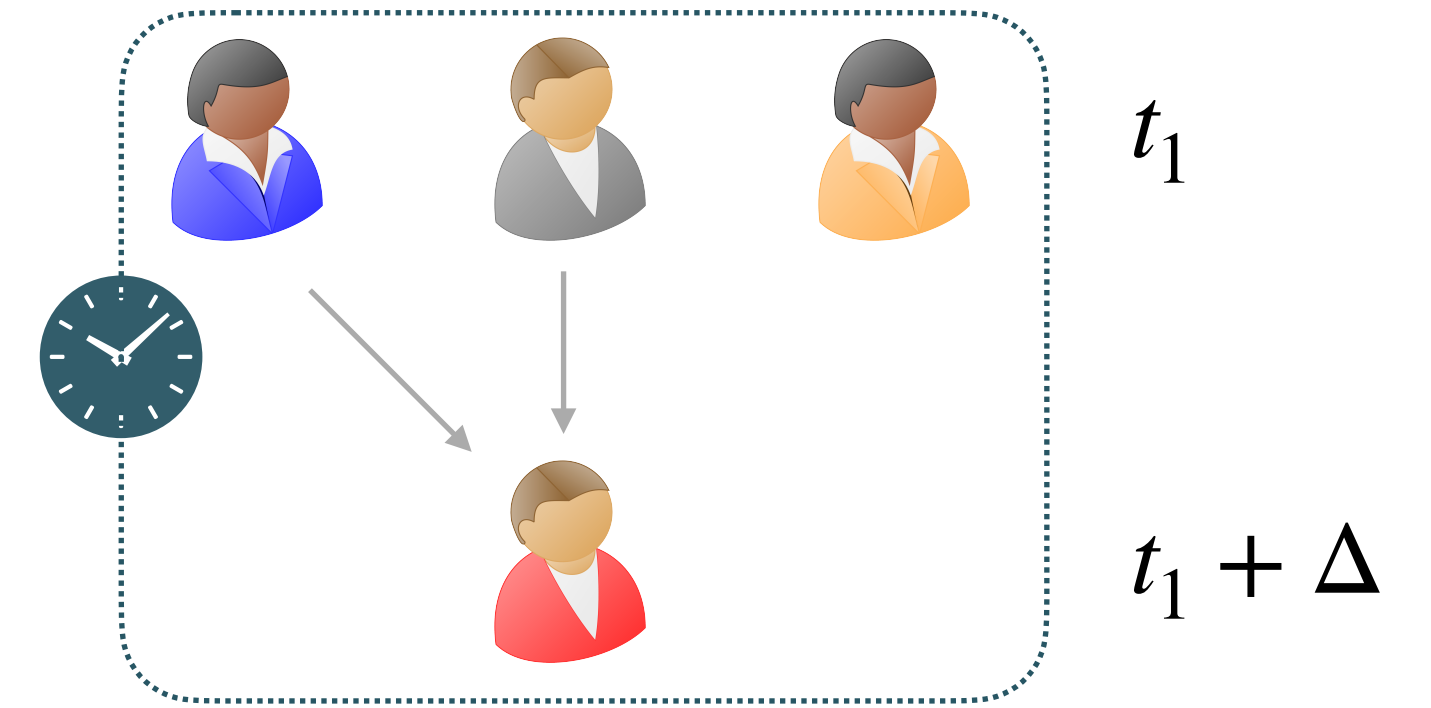
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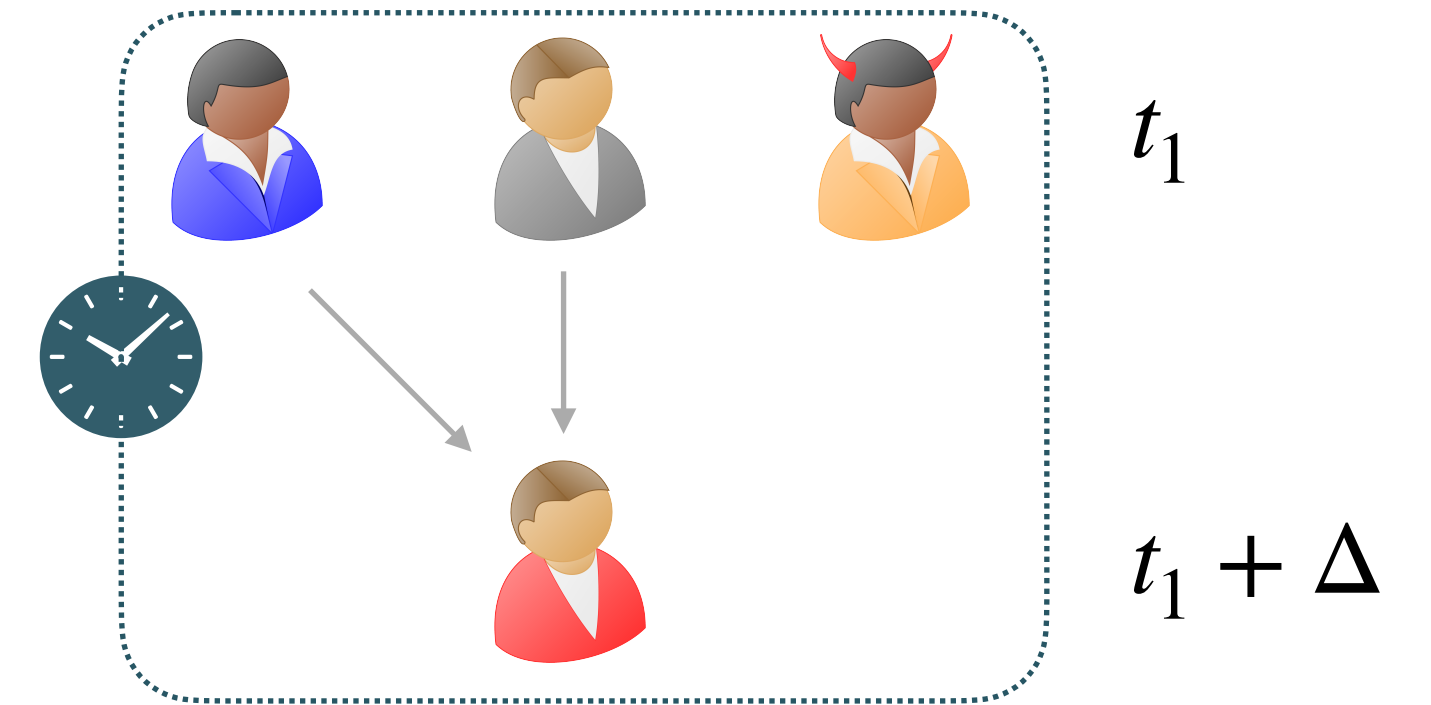
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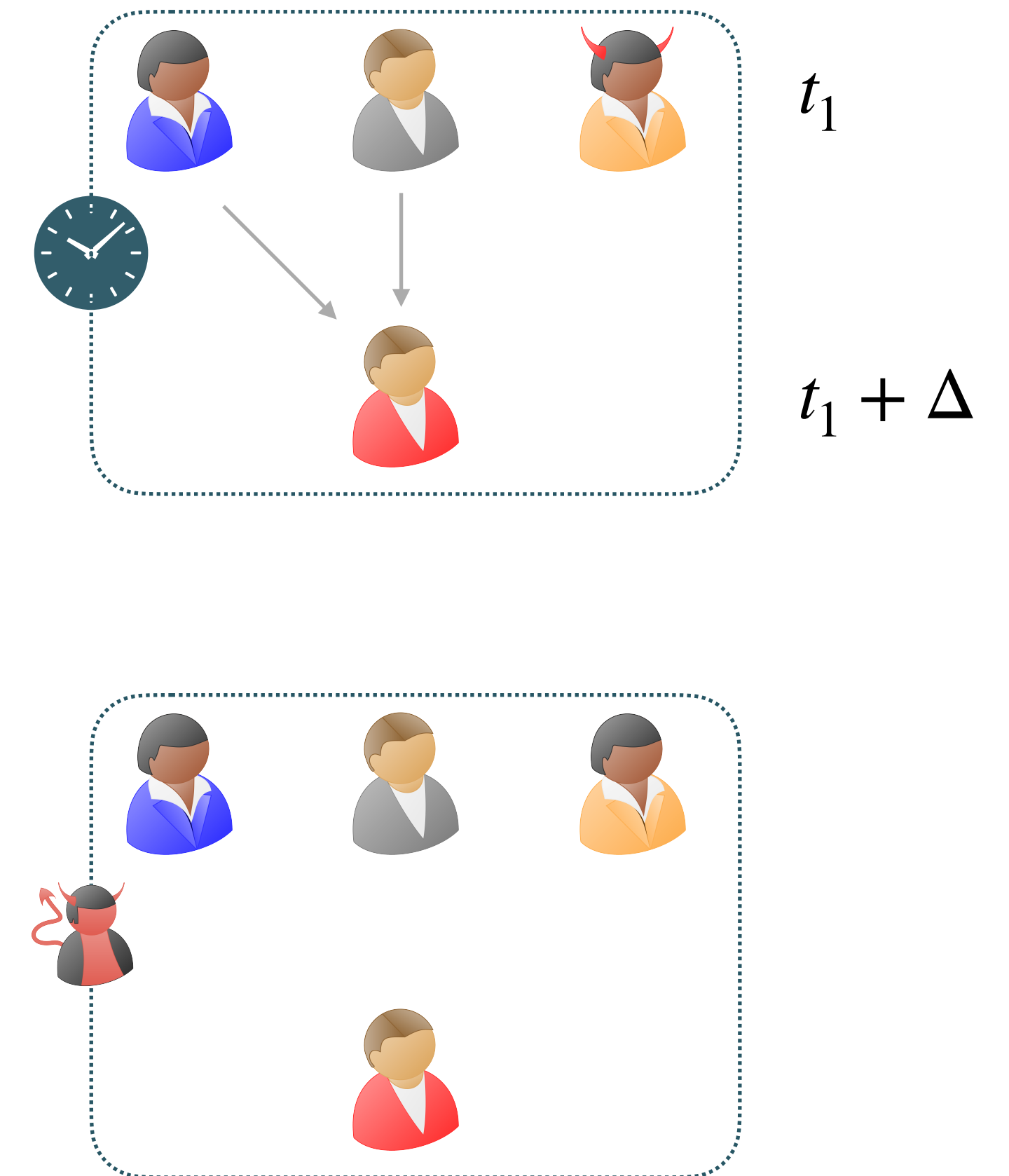
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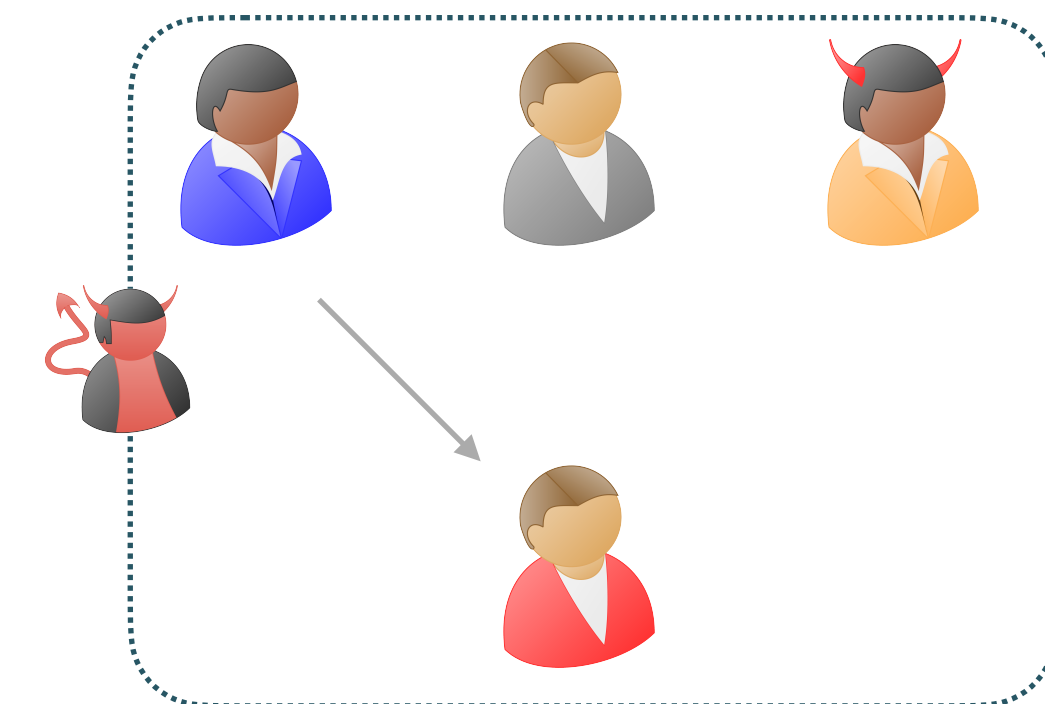
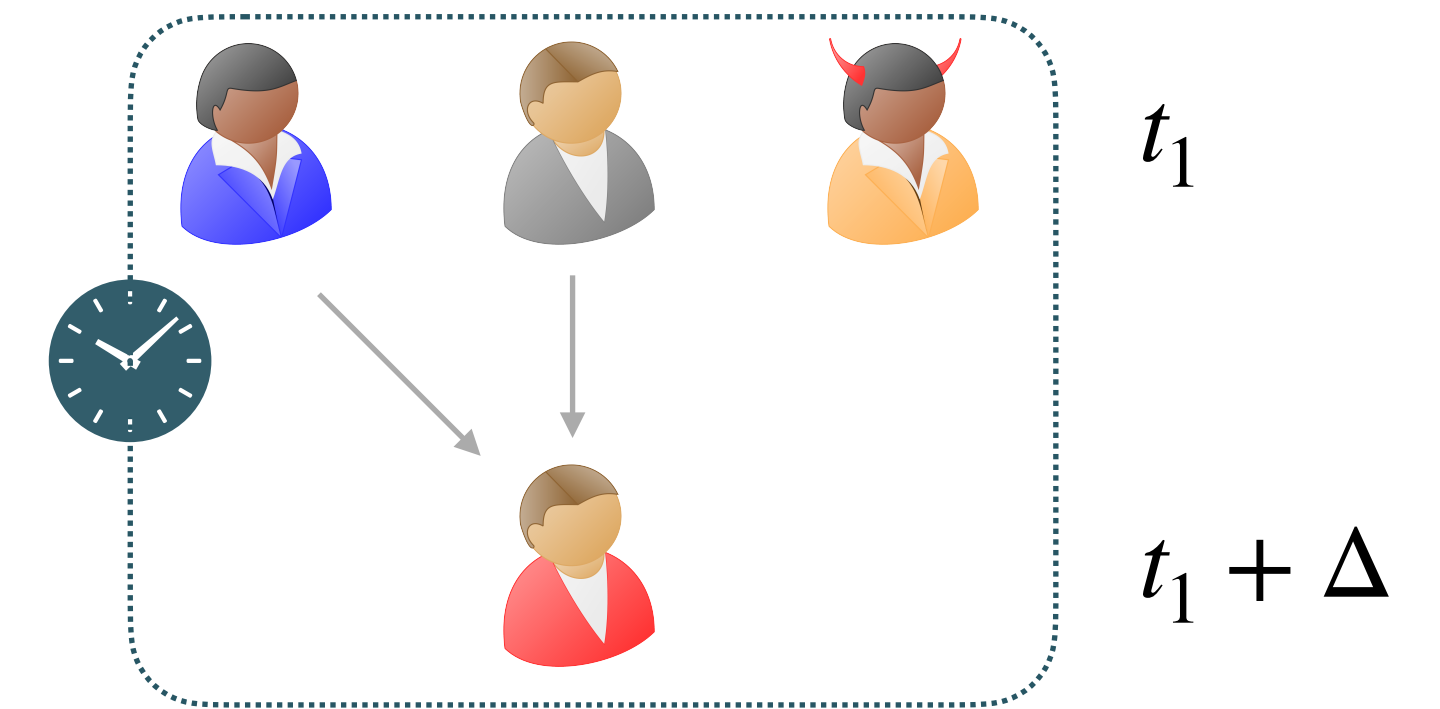
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 - No synchronisation
 - Adversary schedules messages
 - Eventual delivery



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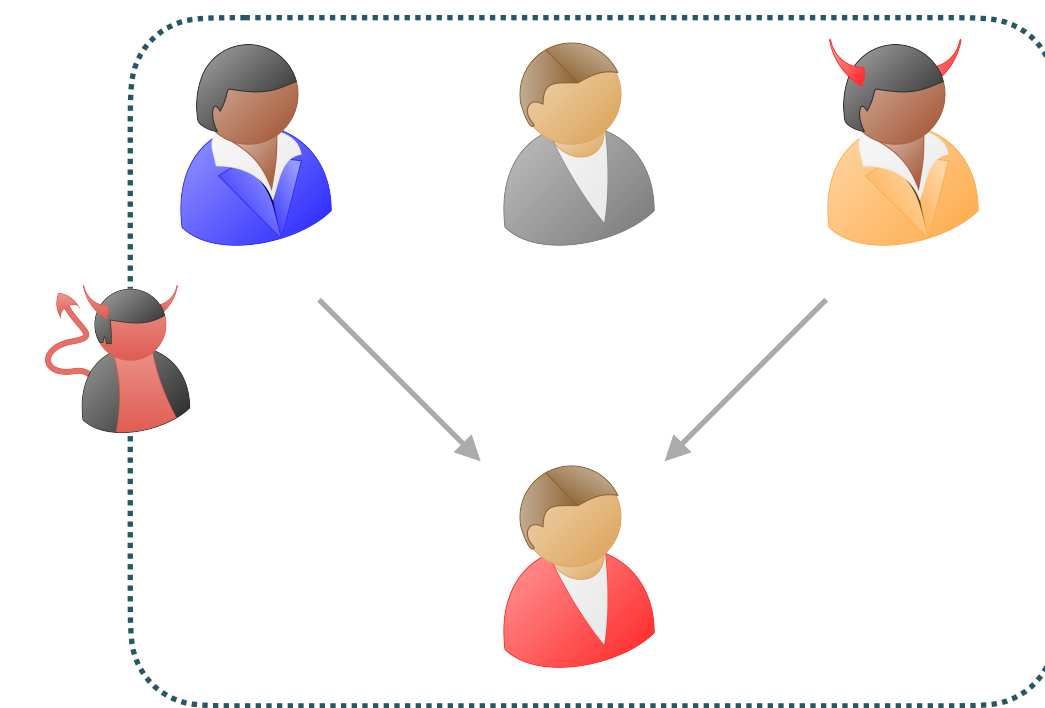
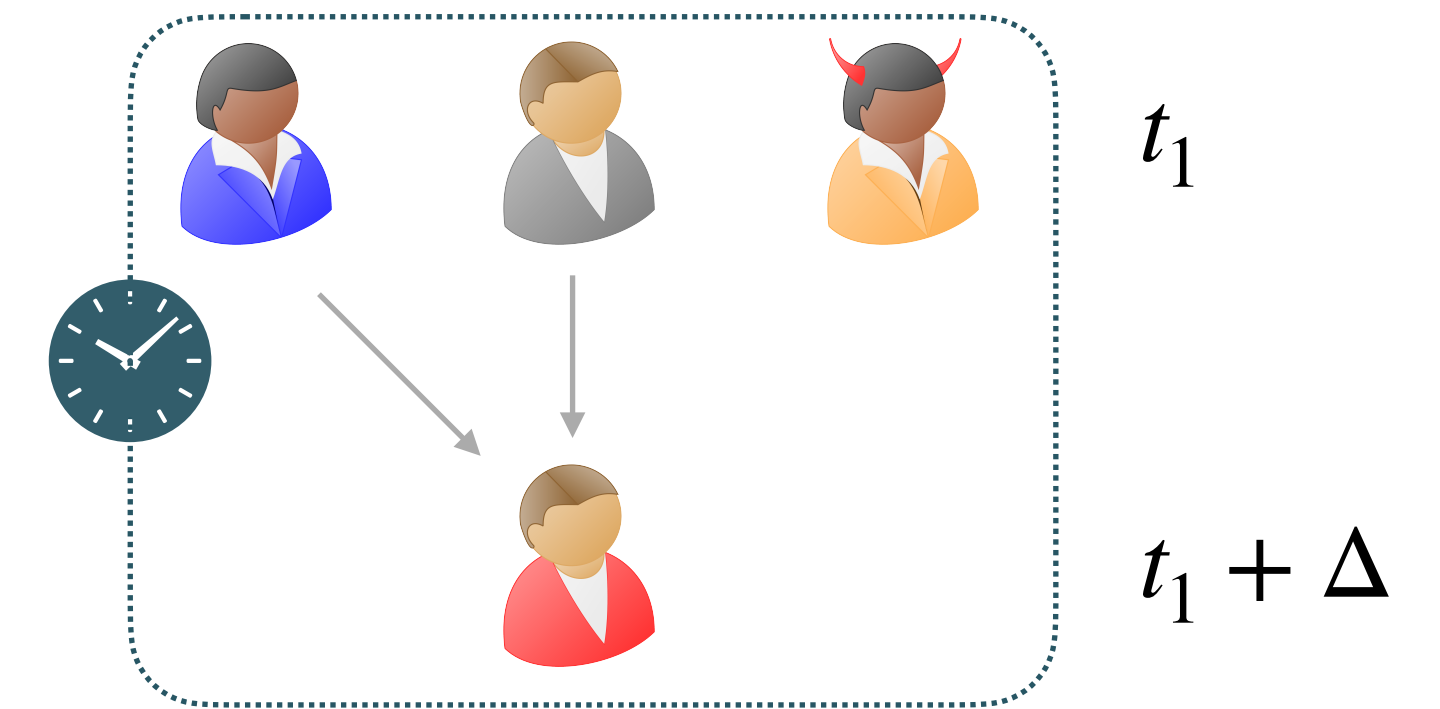
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Cannot distinguish
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Can only wait for
 $n - t$ messages at
each step

Communication Model - Hybrid Networks

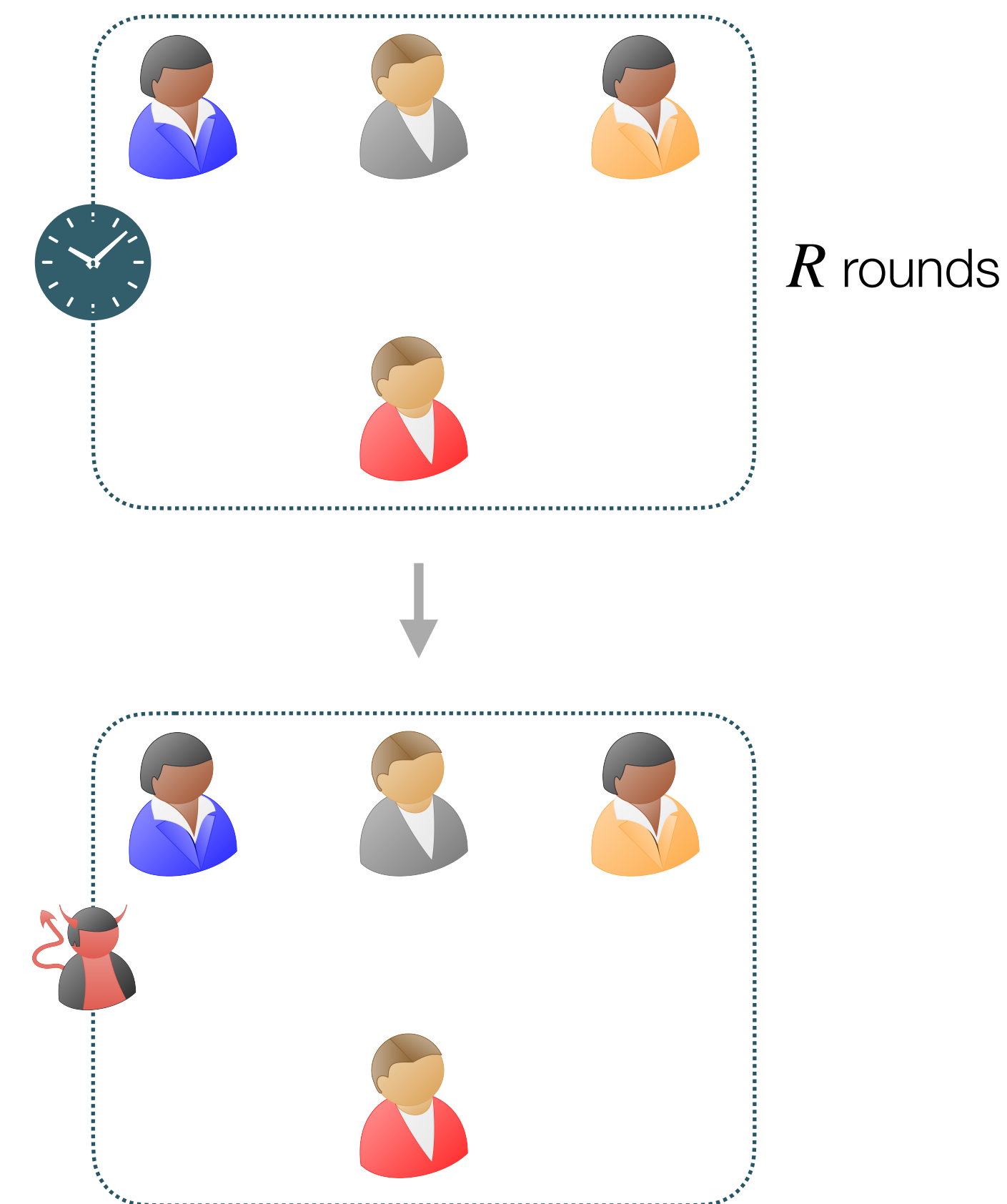
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 - Better model for real-world networks
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- **Hybrid networks:** R initial synchronous rounds followed by asynchronous computation [BHN10,CHP13,PR18]
 - Assume synchronous broadcast channel in first R rounds



Our Contributions

- **Perfectly** secure MPC protocol over **hybrid** network with $R = 2$
 - First protocol in this setting
- **Cryptographically** secure MPC protocol over **hybrid** network with $R = 1$
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Rely only on symmetric key primitives

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Optimal resilience and Guaranteed Output Delivery

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 - Implementation and benchmarks

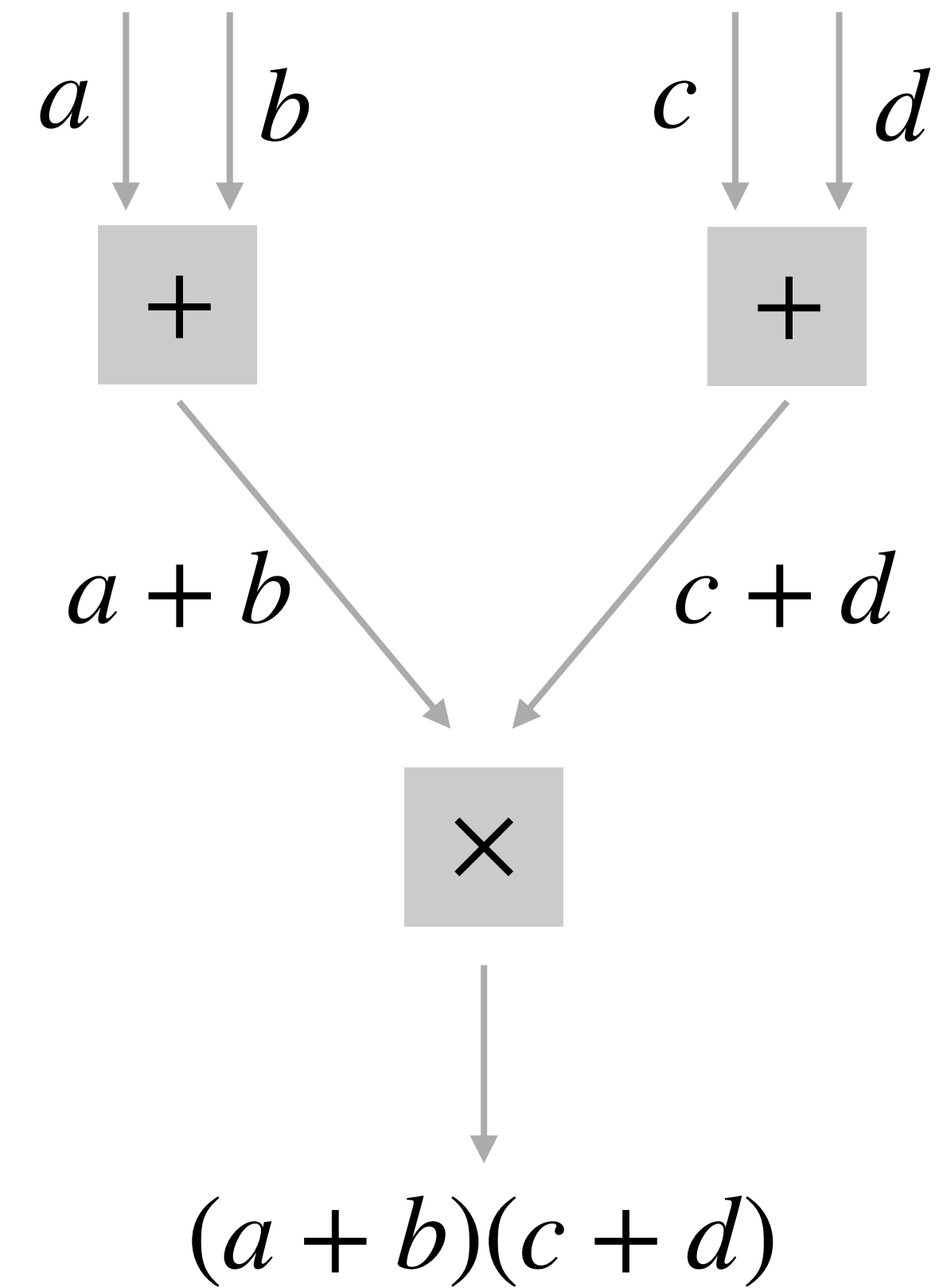
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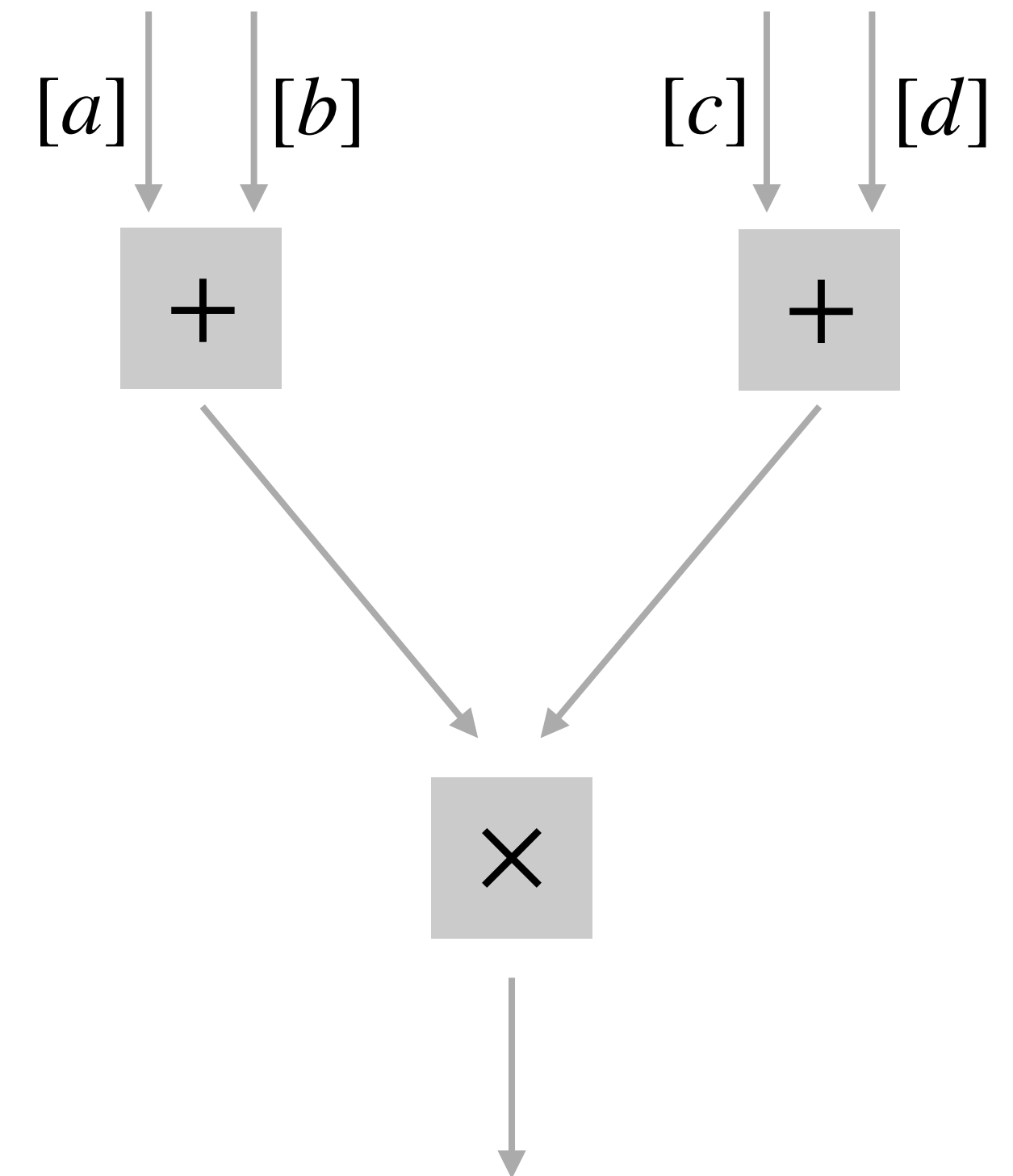
Overview - Circuit Evaluation

- f represented as [arithmetic circuit](#) over finite field



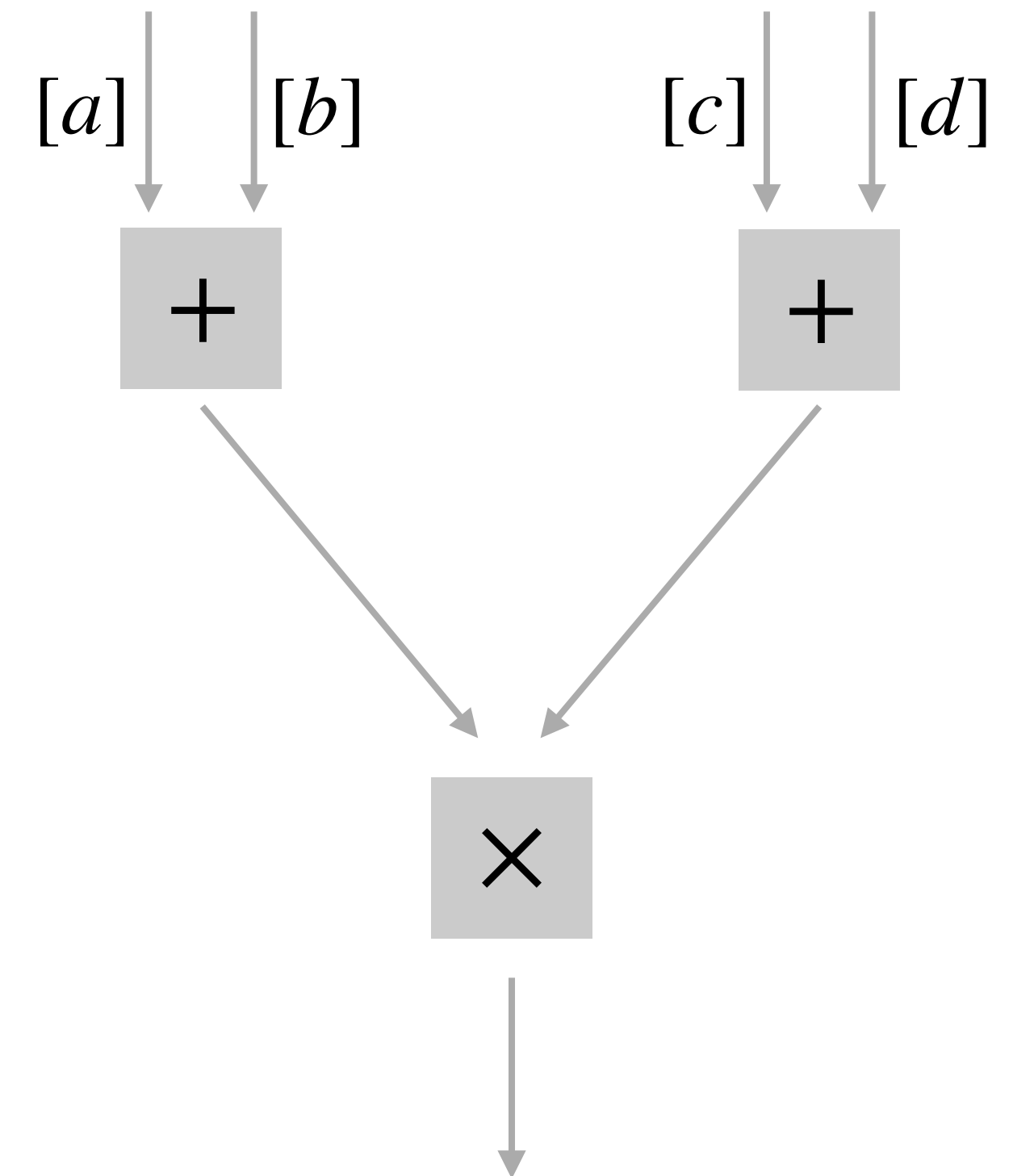
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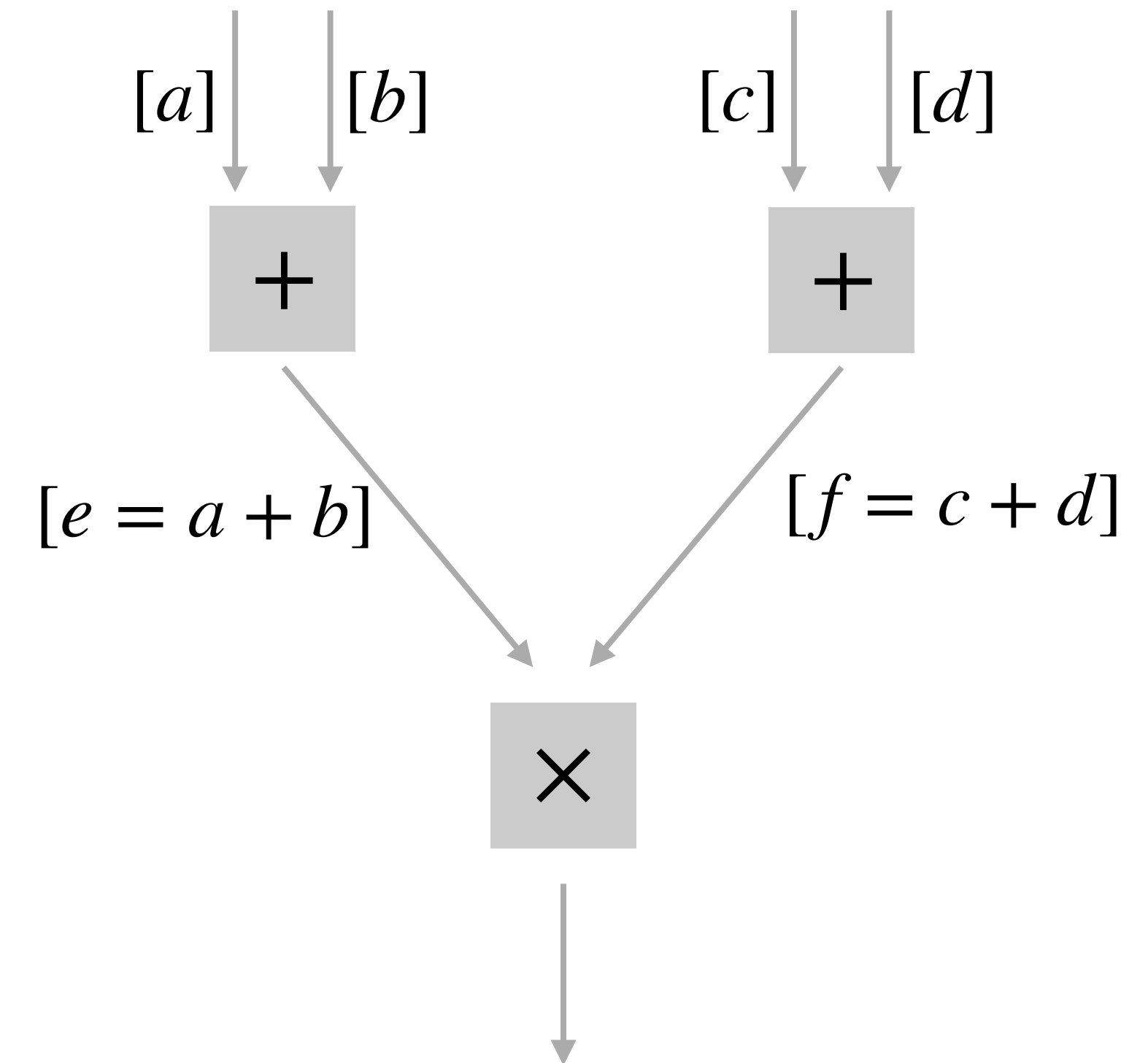
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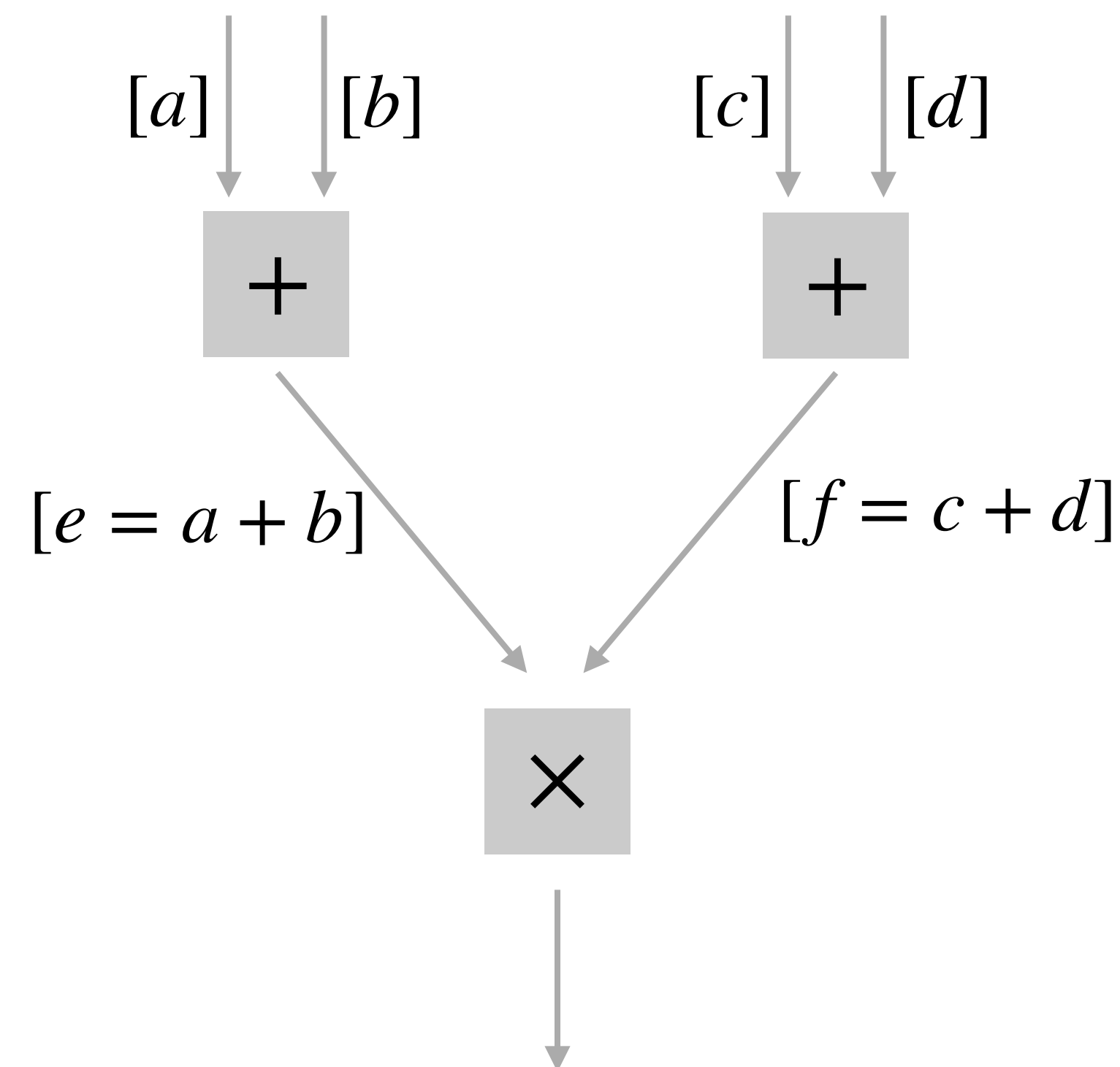
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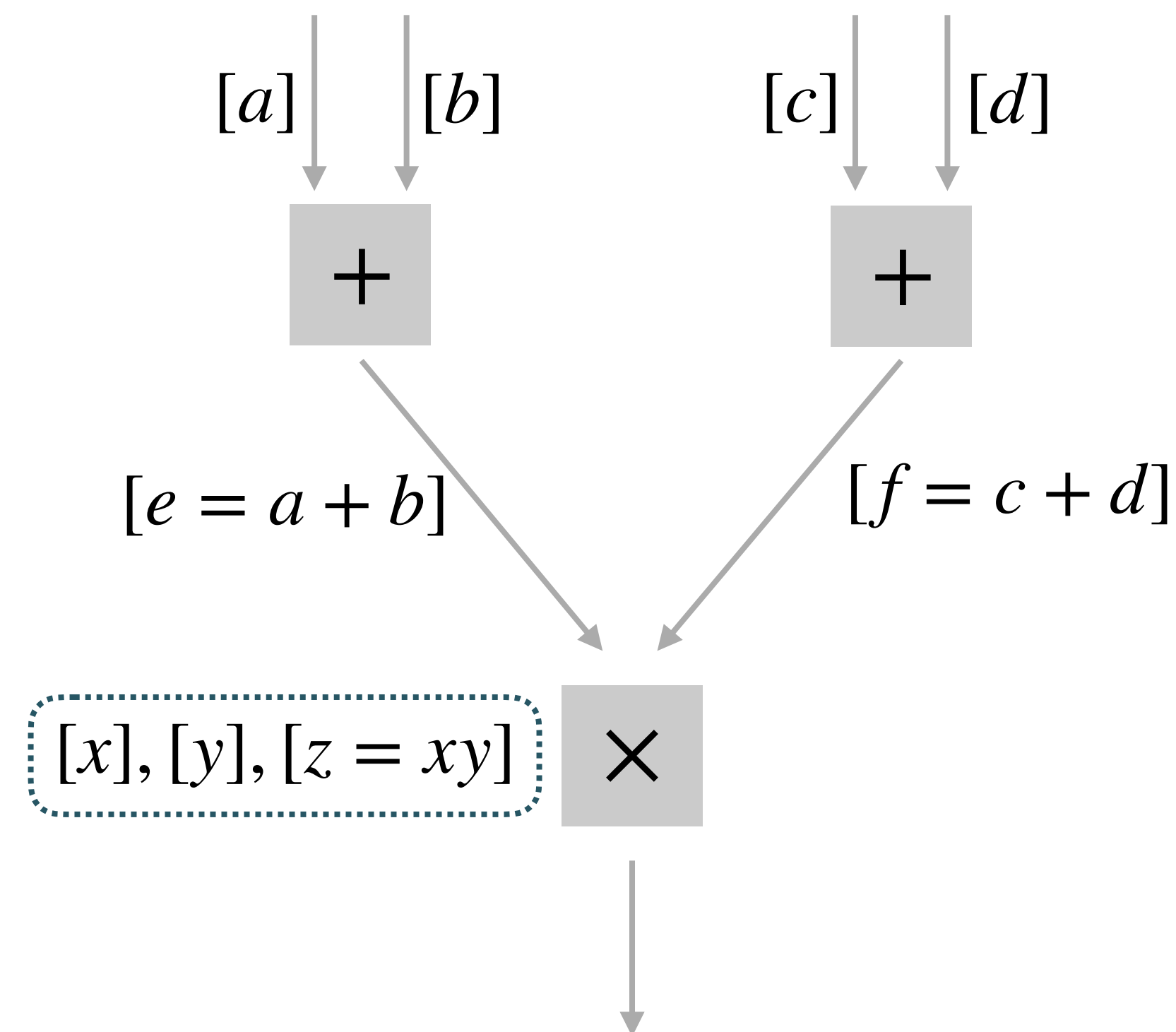
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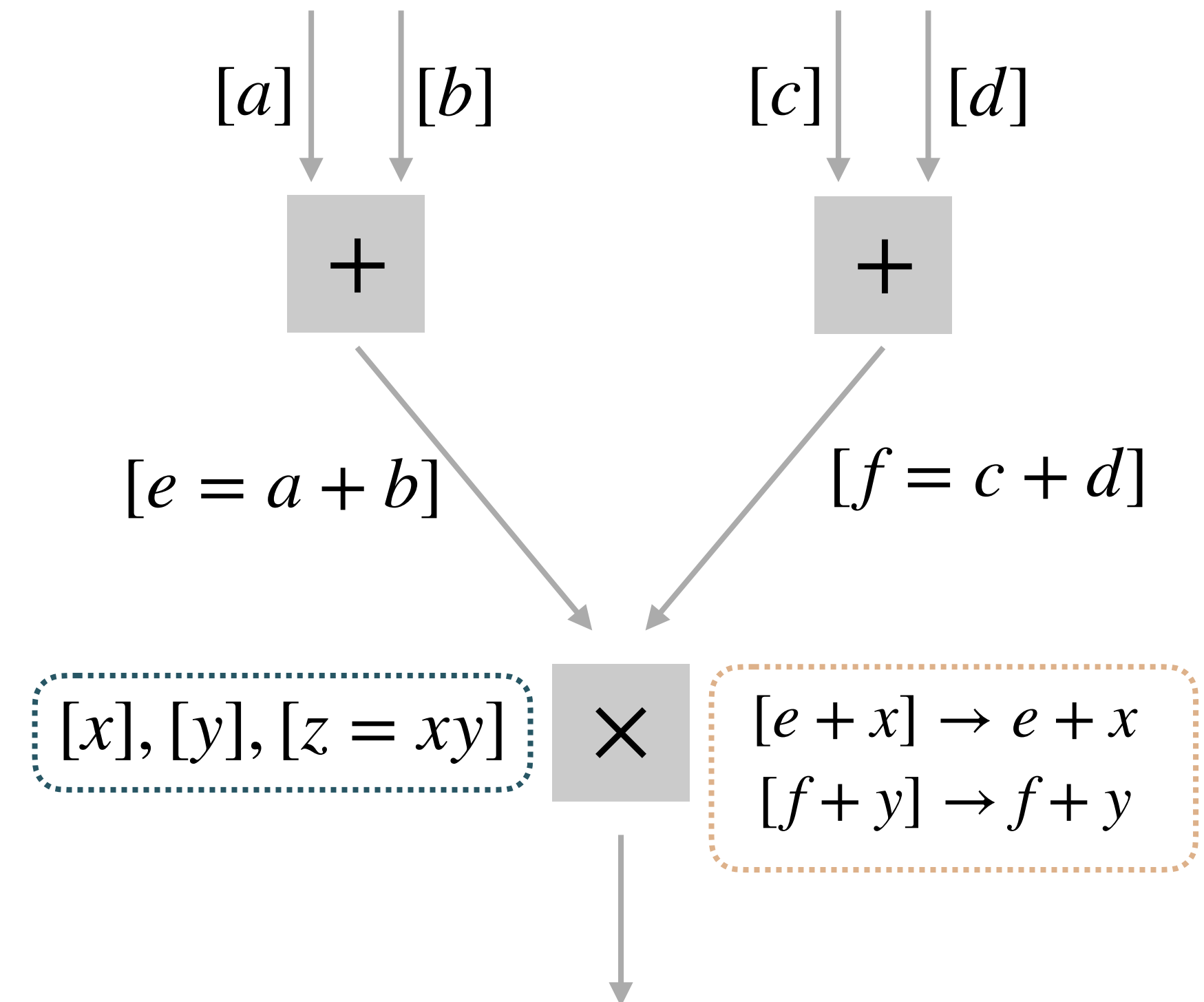
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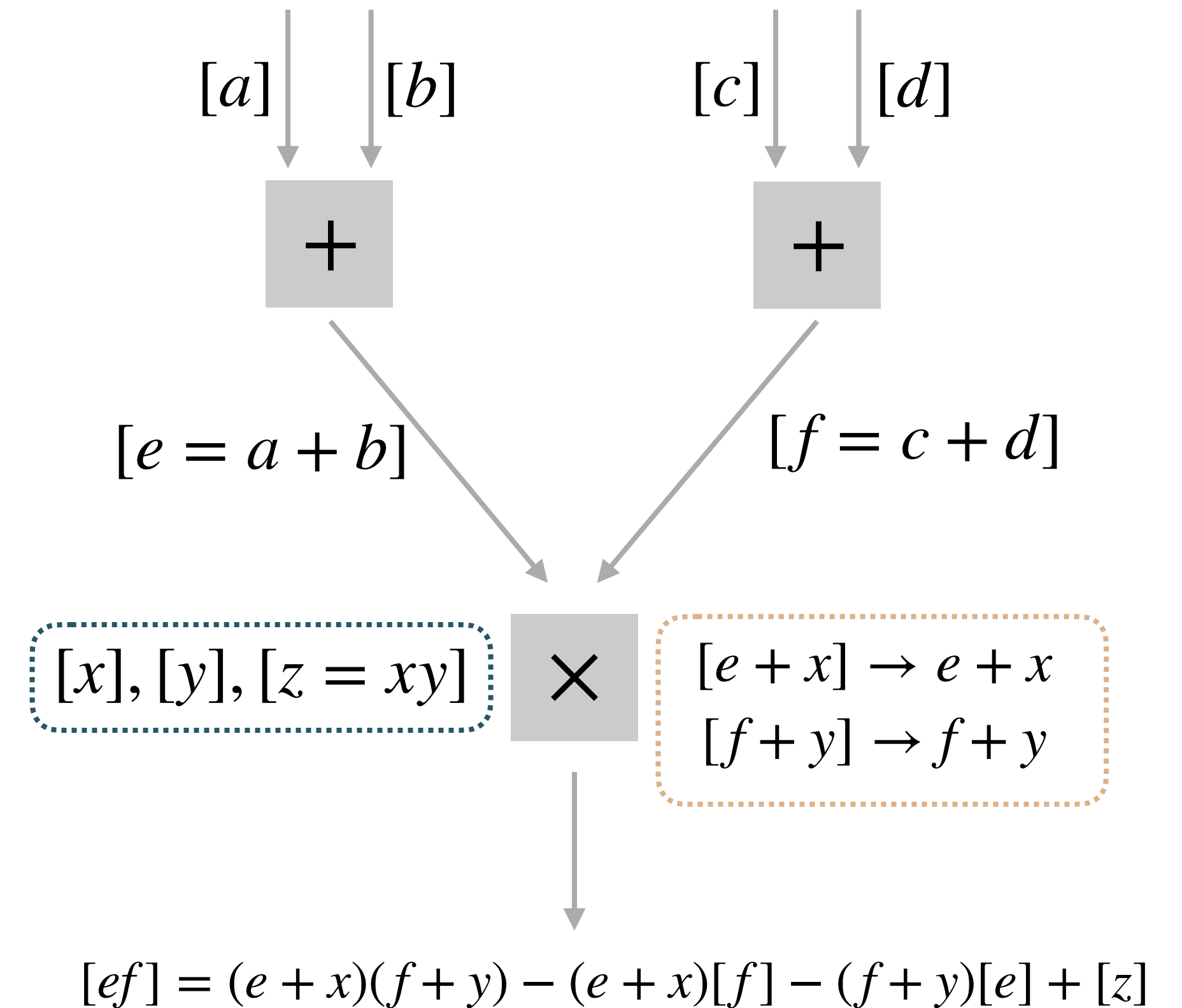
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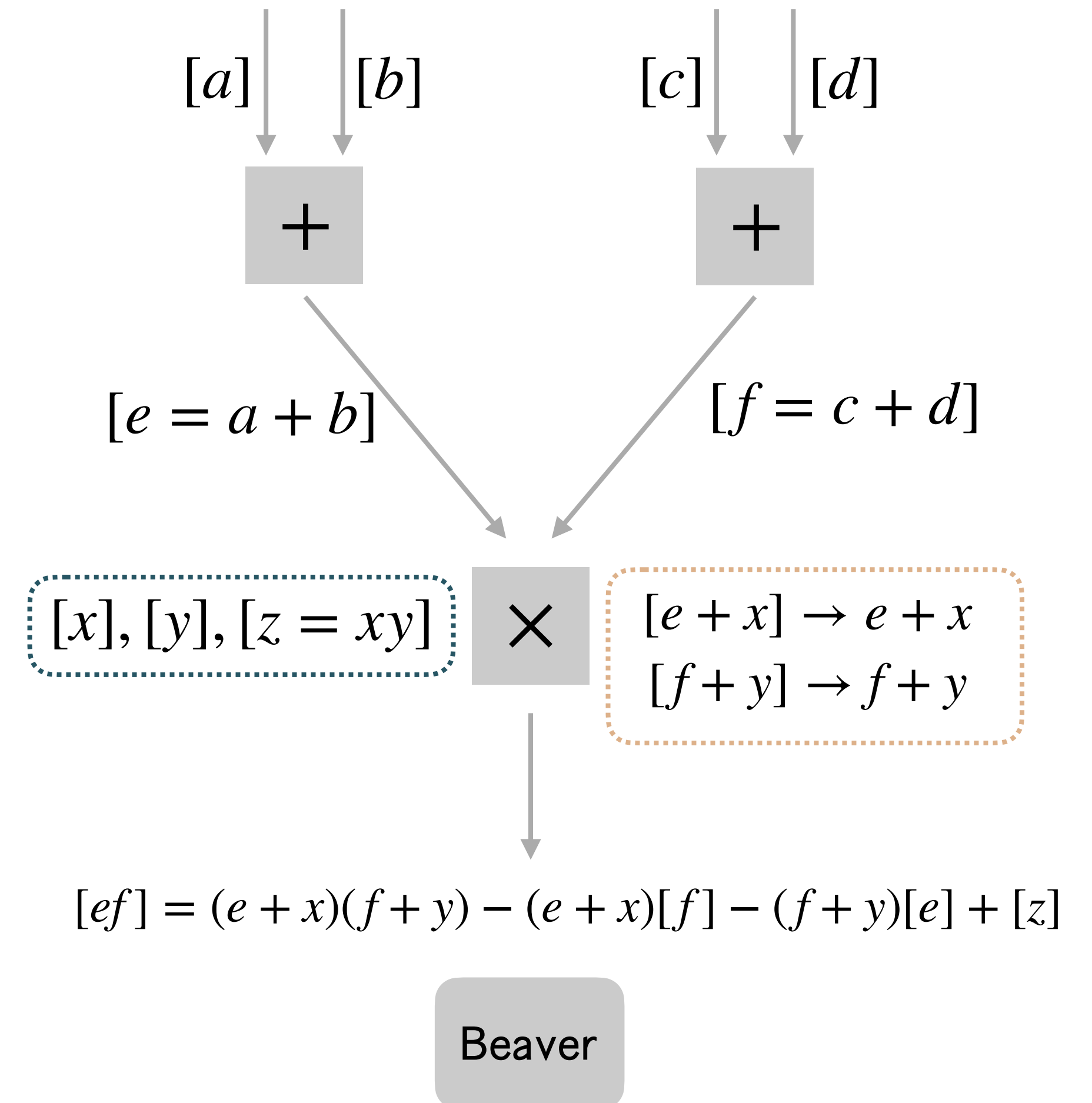
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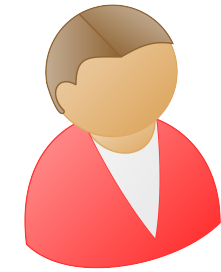
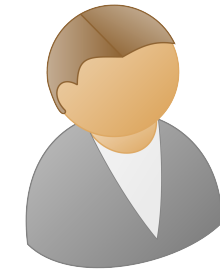
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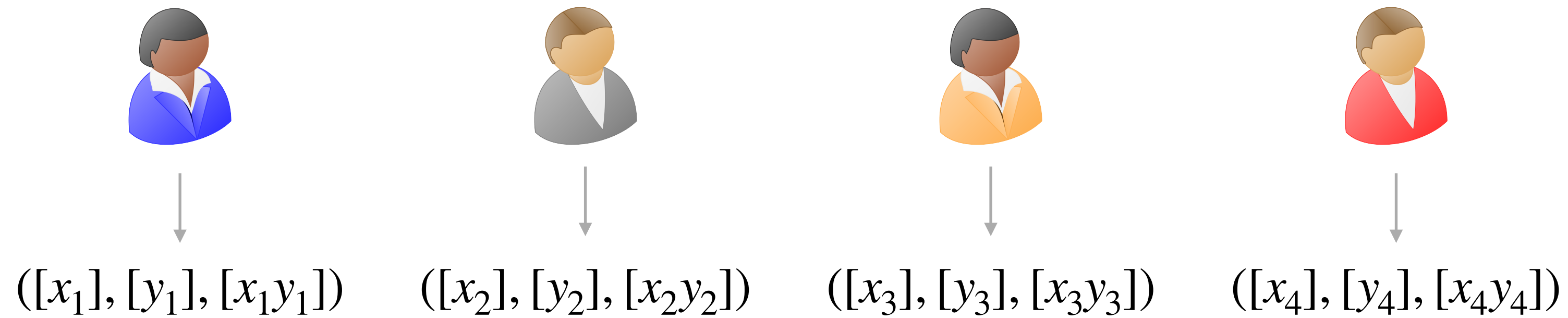
Overview - Generating Multiplication Triples

- Triple generation framework of [CP17]



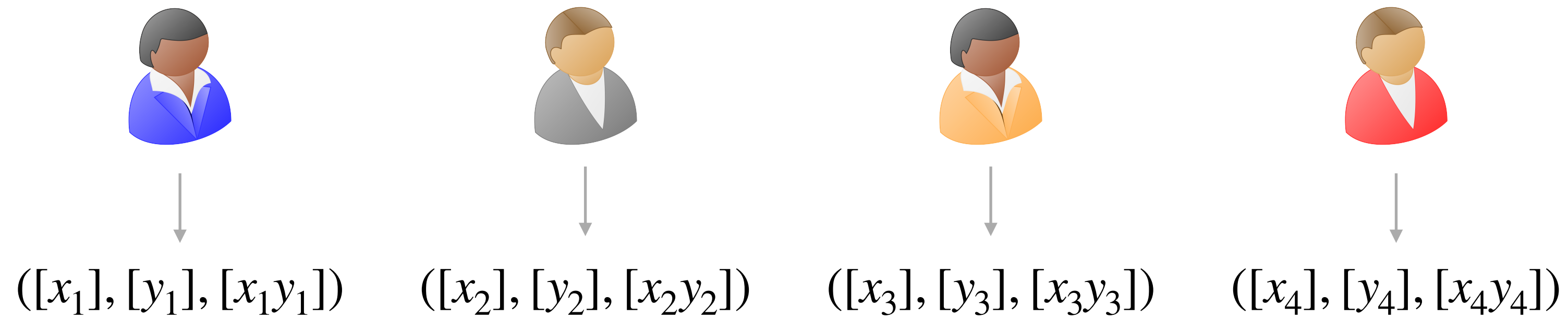
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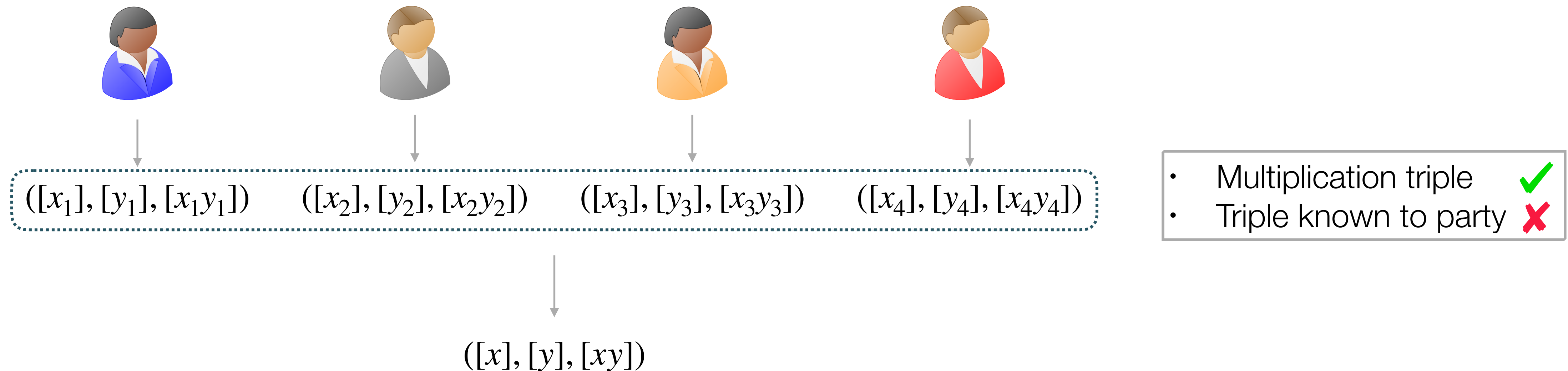
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- | | |
|-------------------------|---|
| • Multiplication triple | ✓ |
| • Triple known to party | ✗ |

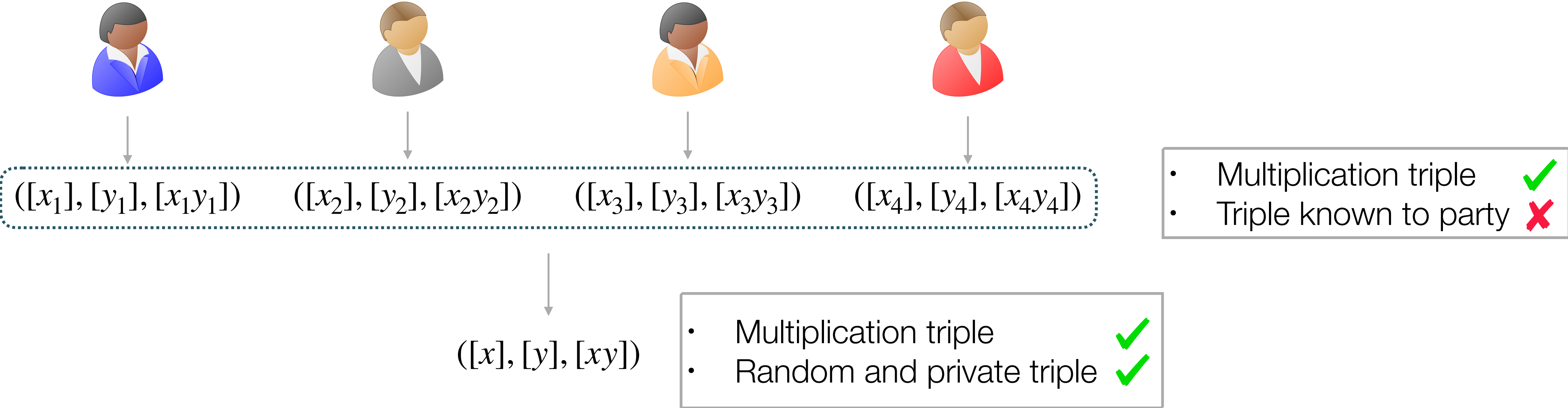
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Perfect HMPC

- Open Problem [PR18]: Perfectly secure MPC protocol over hybrid network
 - Two synchronous rounds
 - Tolerating $t < n/3$ corruptions
 - With synchronous broadcast channel
 - Guaranteed output delivery
- Input provision impossible in this setting [PR18]

Perfect HMPC - Linear Secret Sharing Scheme

- Replicated Secret Sharing [ISN89]

- $[s] = (s_1, s_2, s_3, s_4)$

- $s = s_1 + s_2 + s_3 + s_4$

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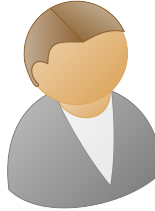
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- P_i does not have s_i



P_1

(s_2, s_3, s_4)



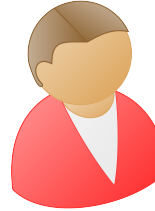
P_2

(s_3, s_4, s_1)



P_3

(s_4, s_1, s_2)



P_4

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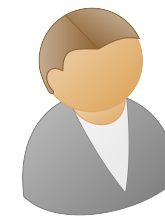
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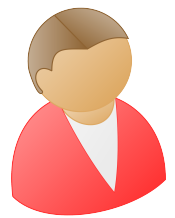
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P_3

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P_4

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- P_i does not have s_i

- All other parties except P_i have s_i

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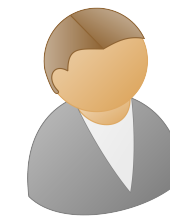
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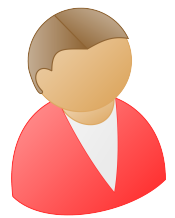
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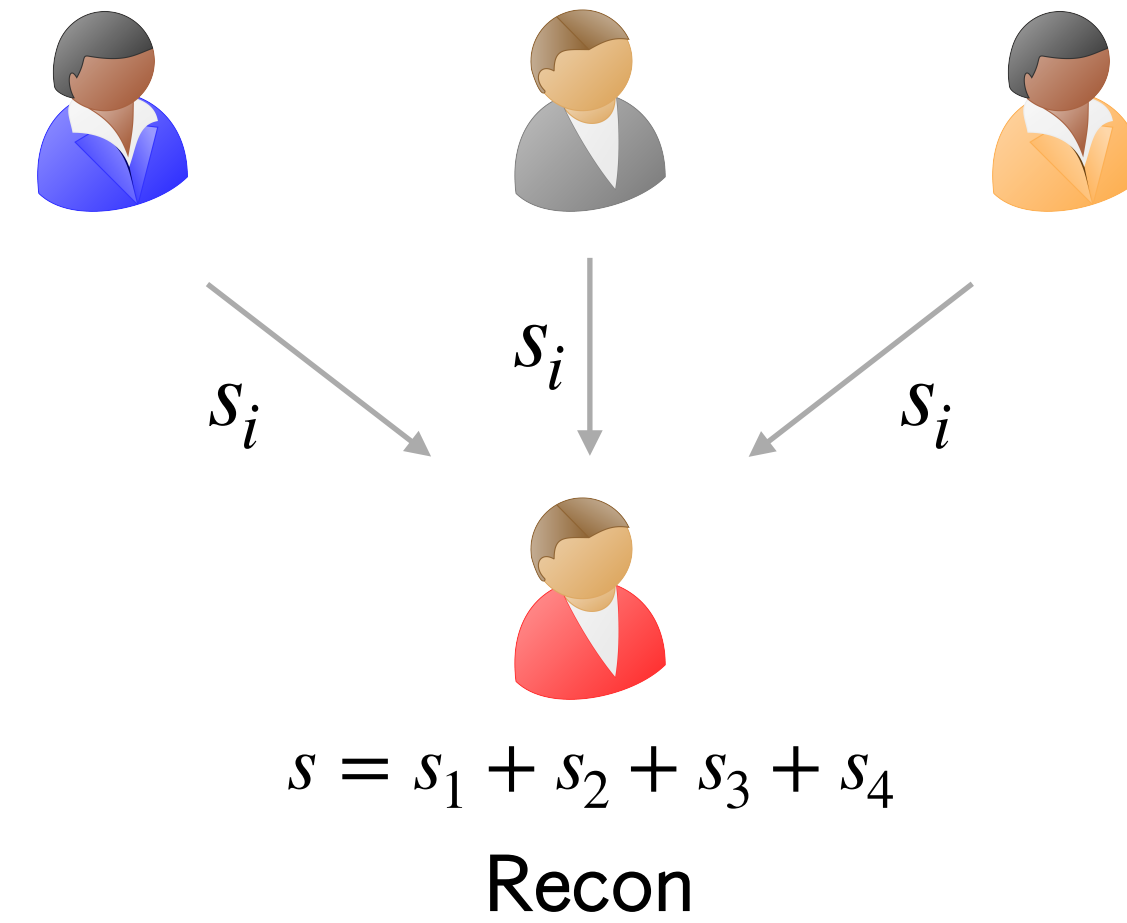
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$$[cs] = (cs_1, cs_2, cs_3, cs_4)$$

$$[s + s'] = (s_1 + s'_1, s_2 + s'_2, s_3 + s'_3, s_4 + s'_4)$$

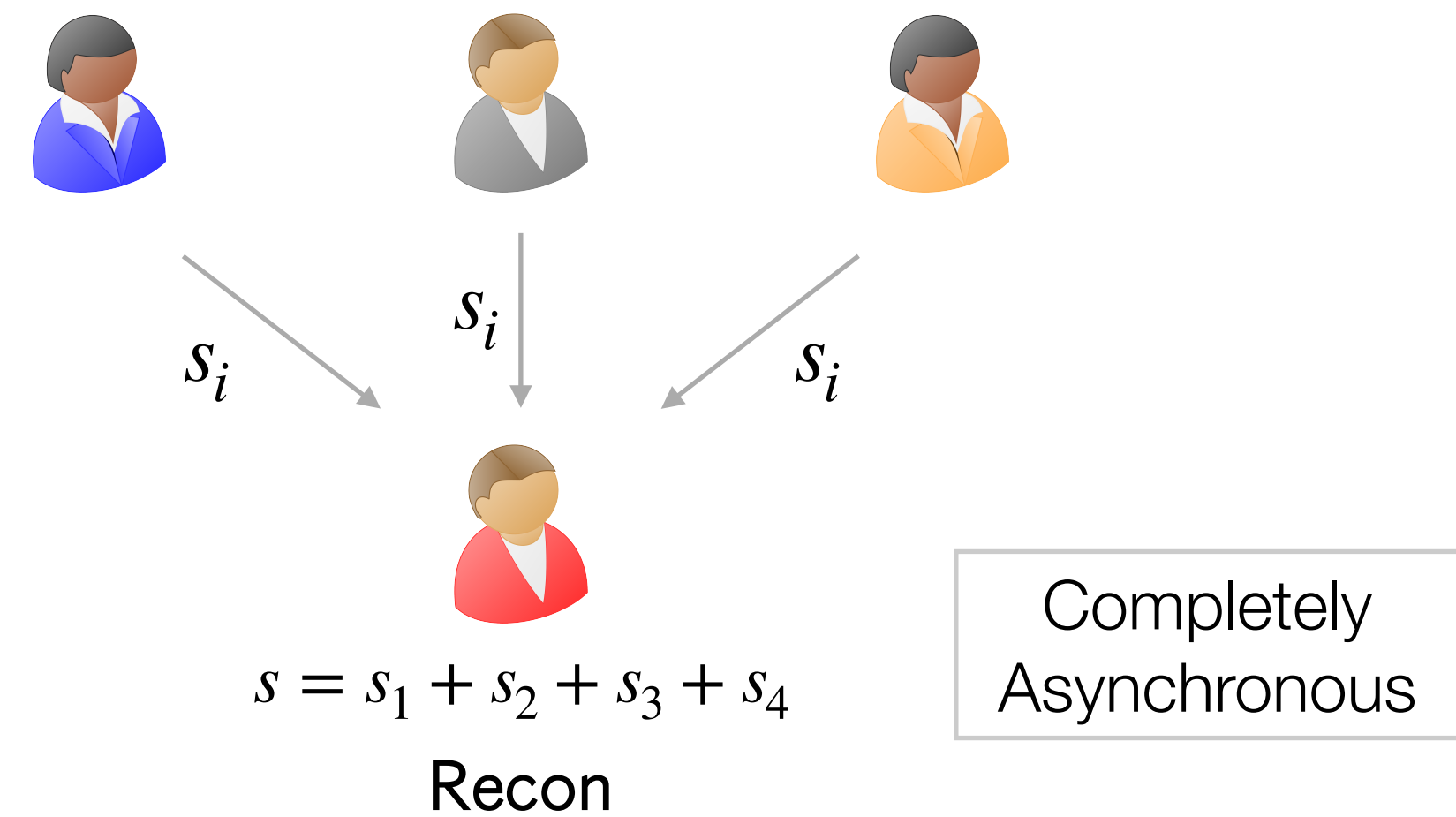
Perfect HMPC - Linear Secret Sharing Scheme

- Reconstruction
 - P_j sends s_i to P_i
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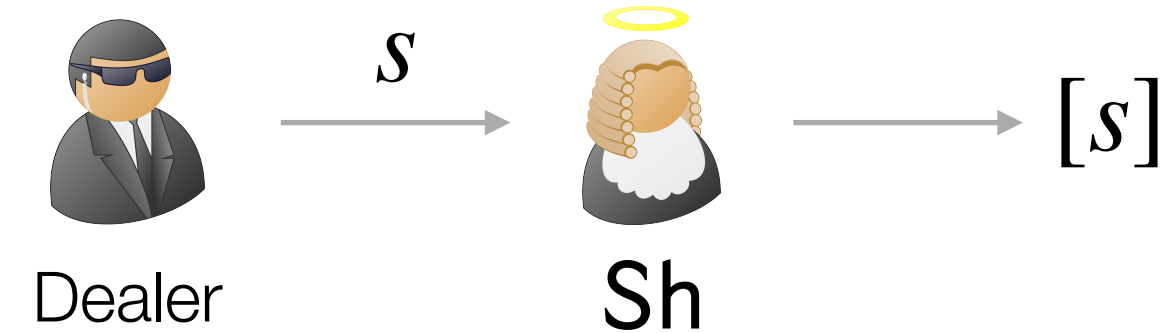
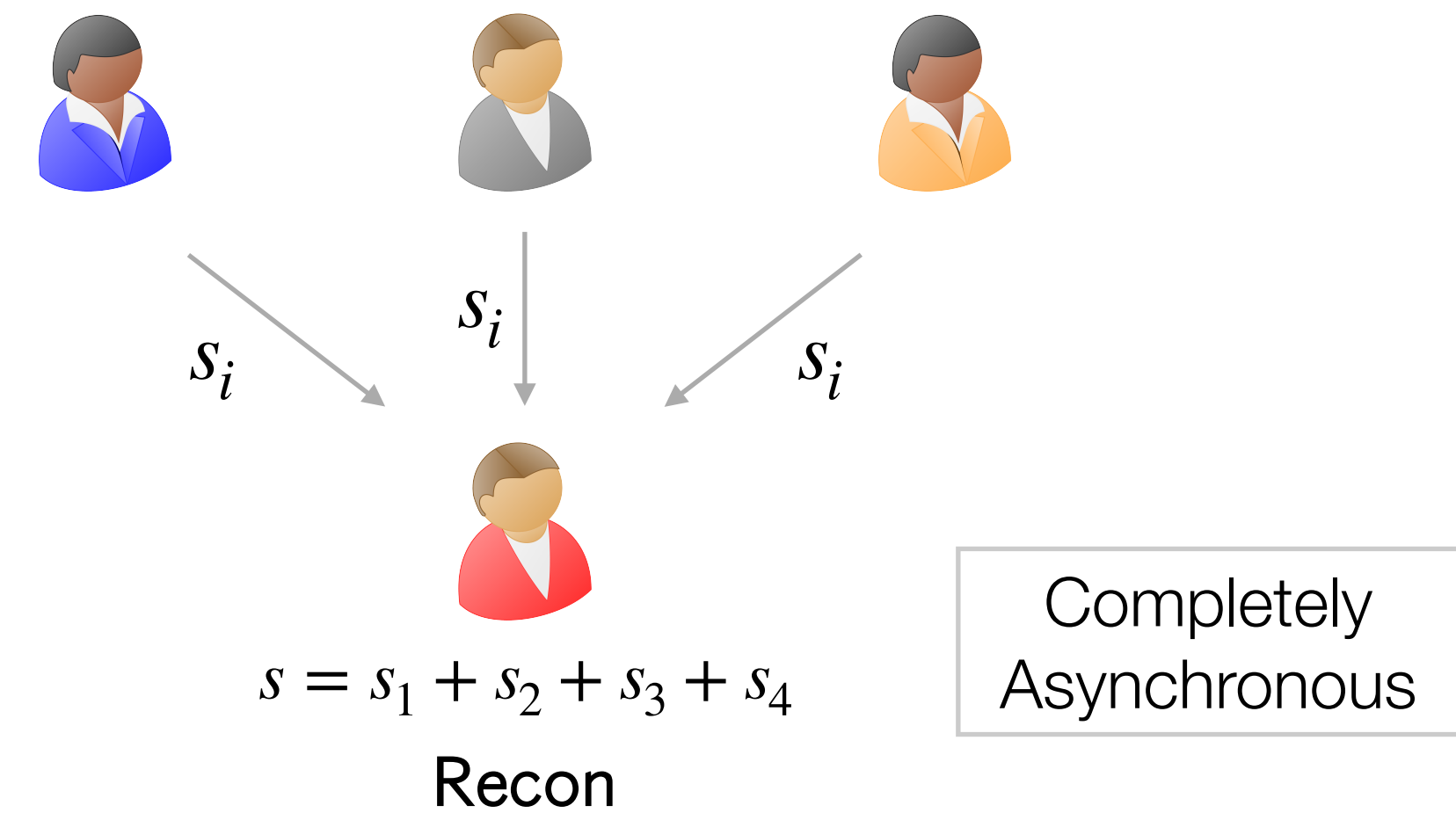
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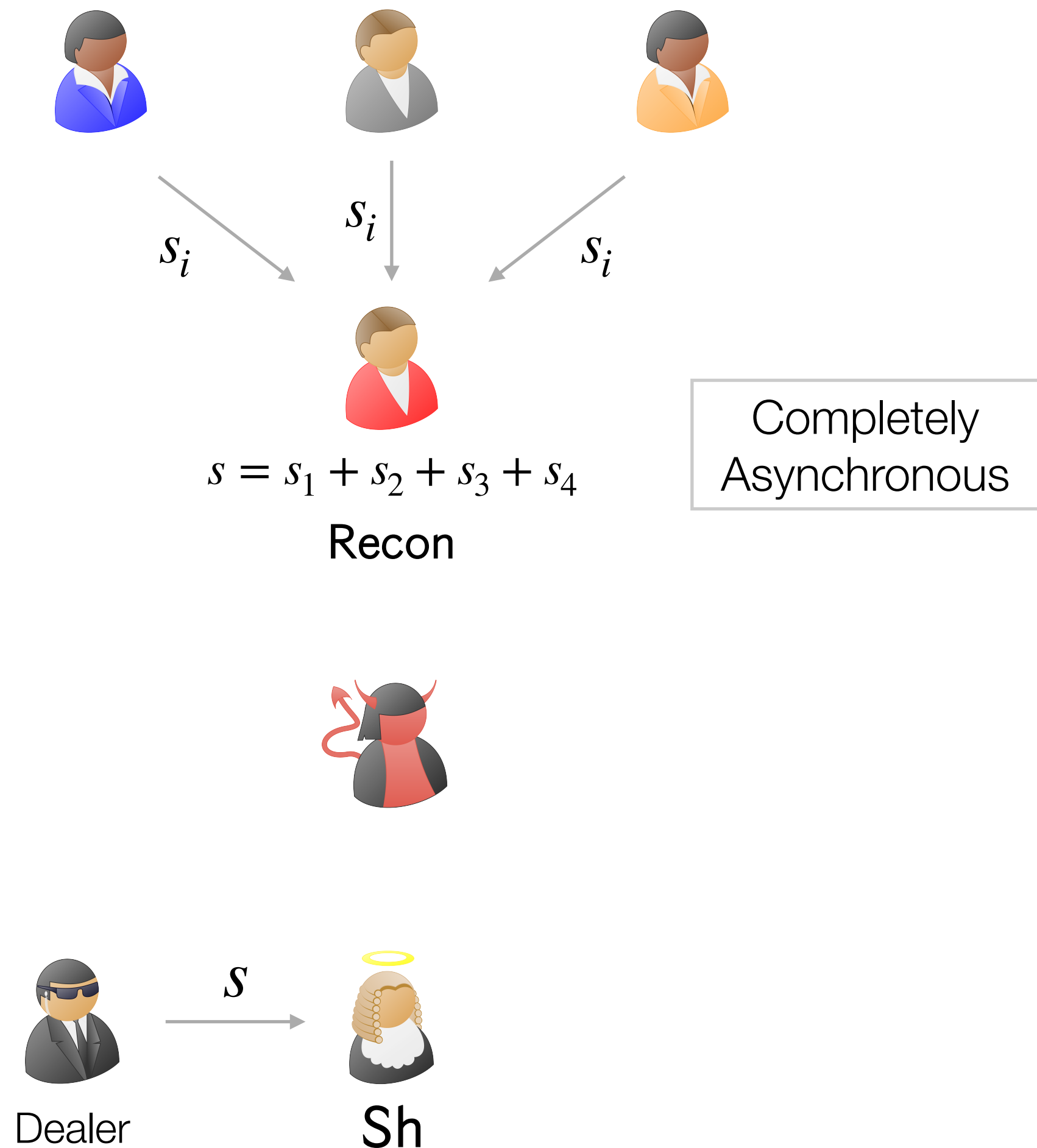
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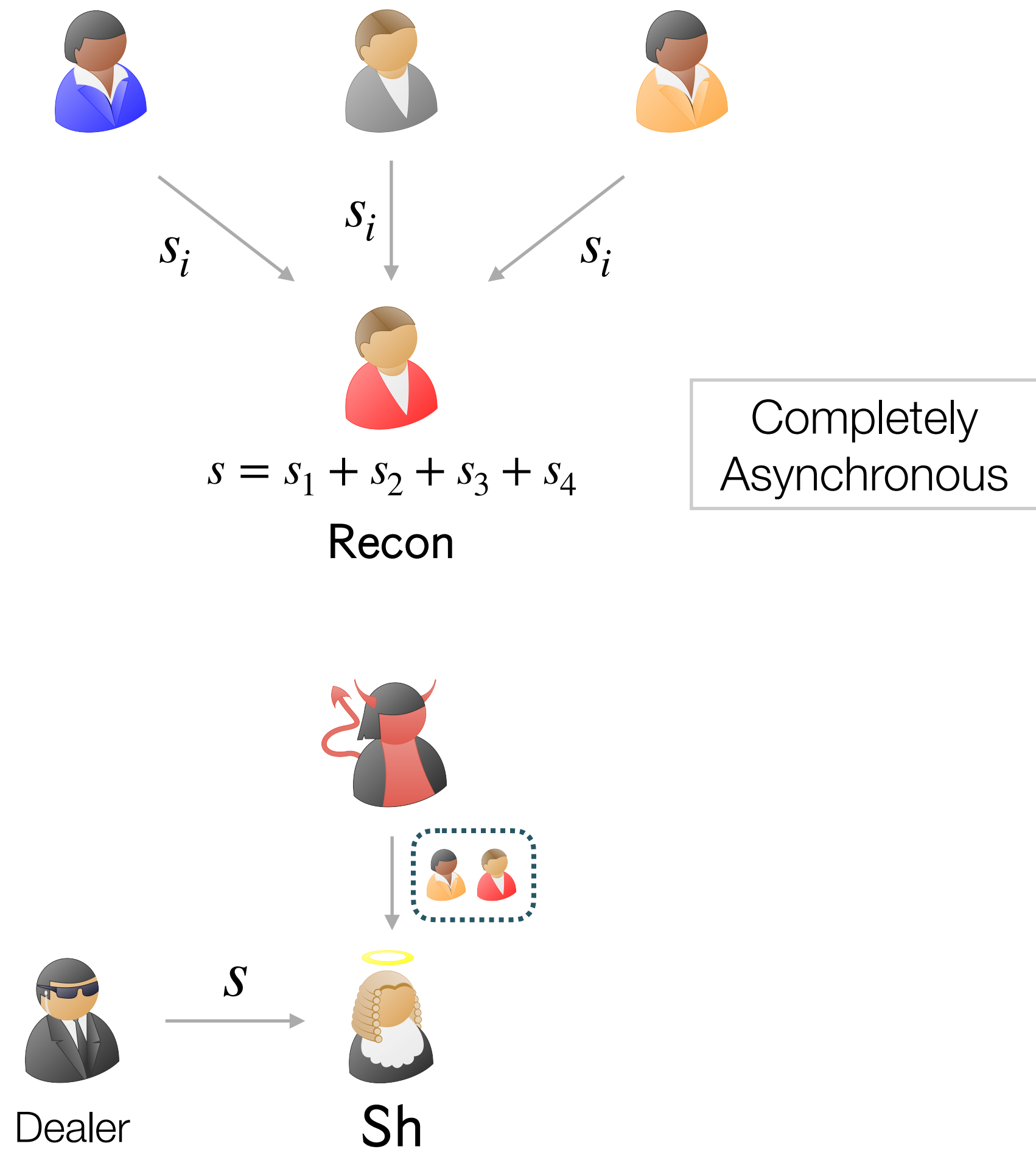
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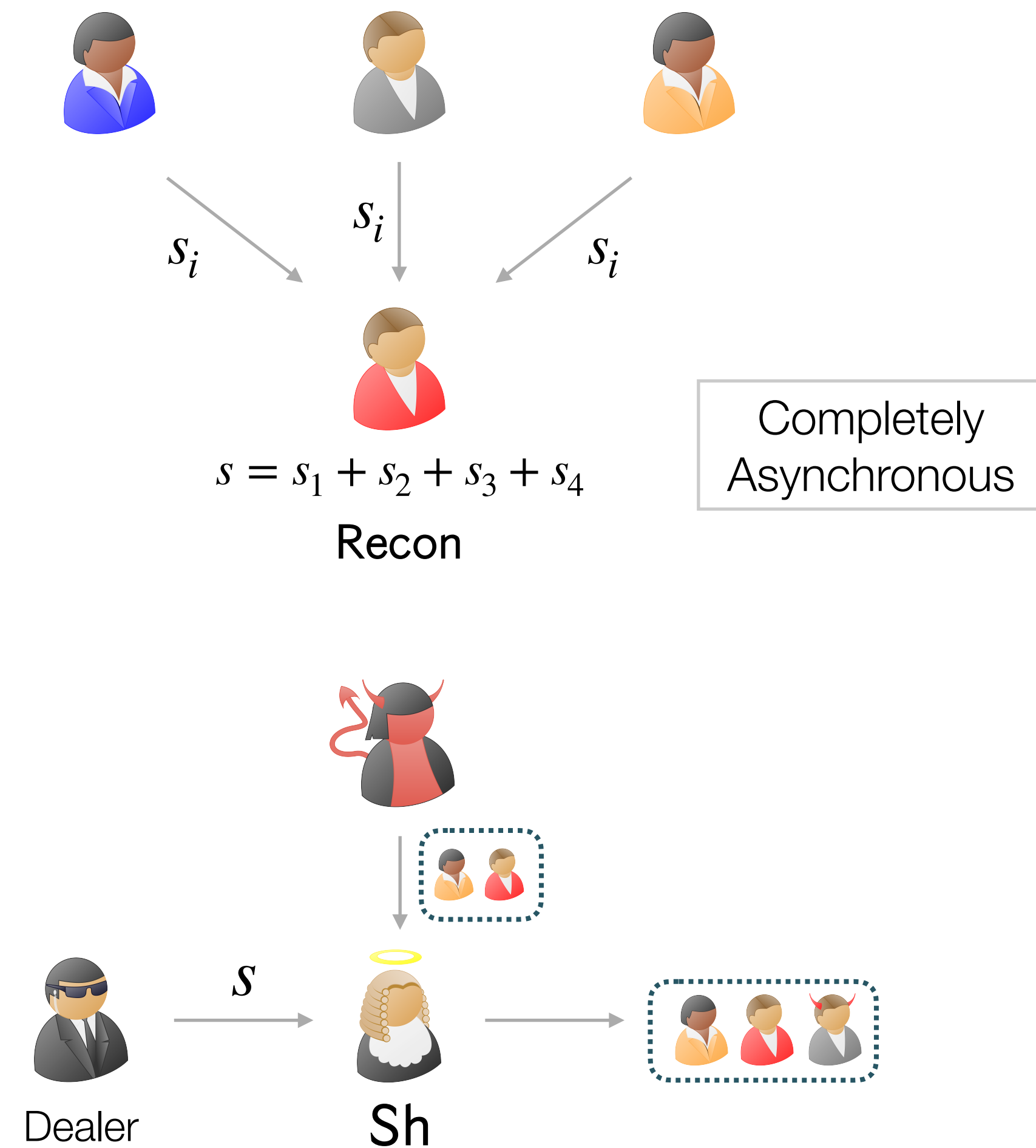
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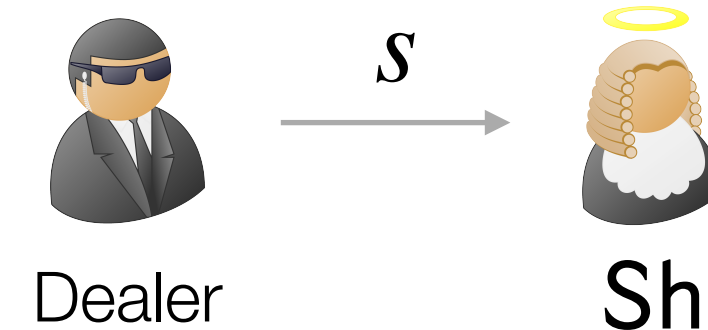
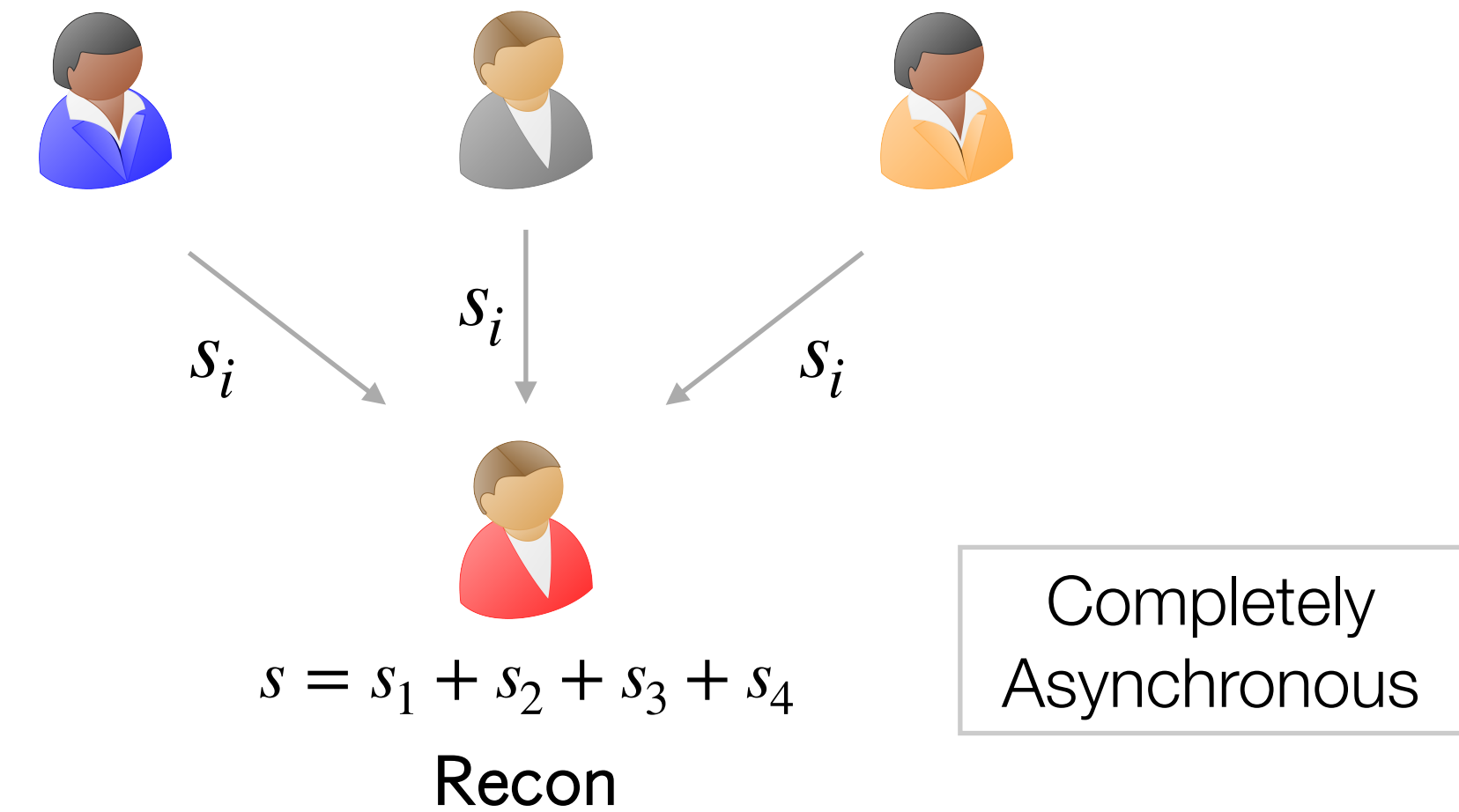
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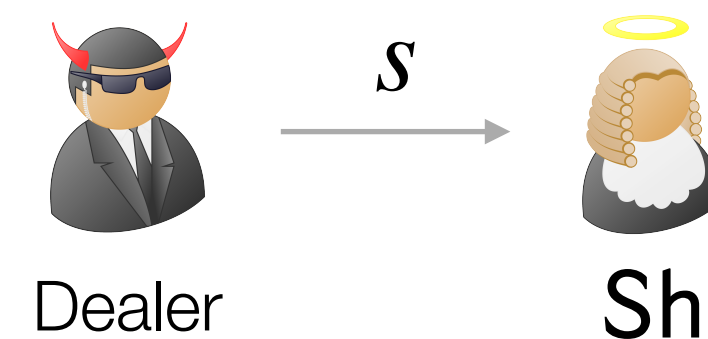
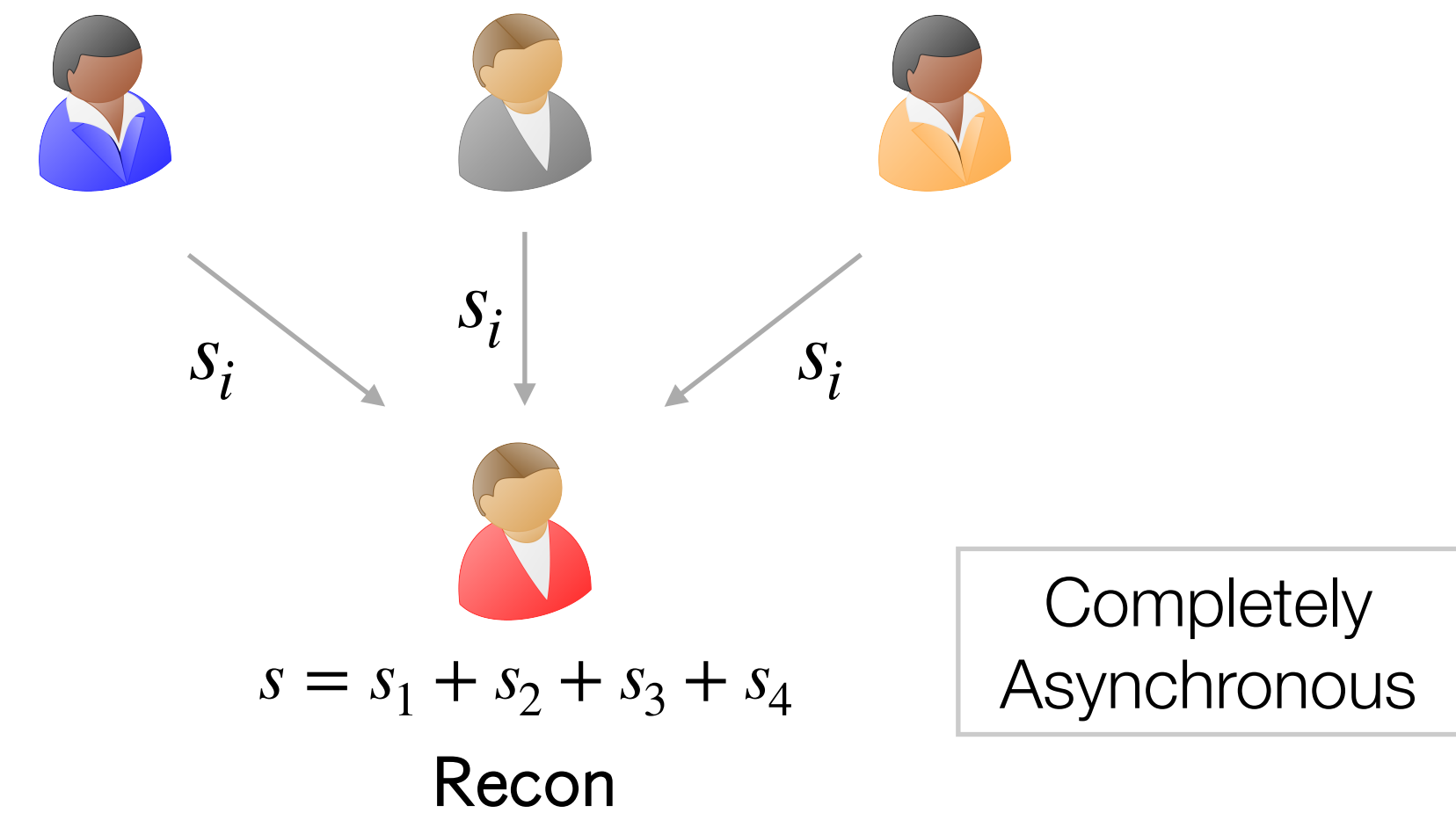
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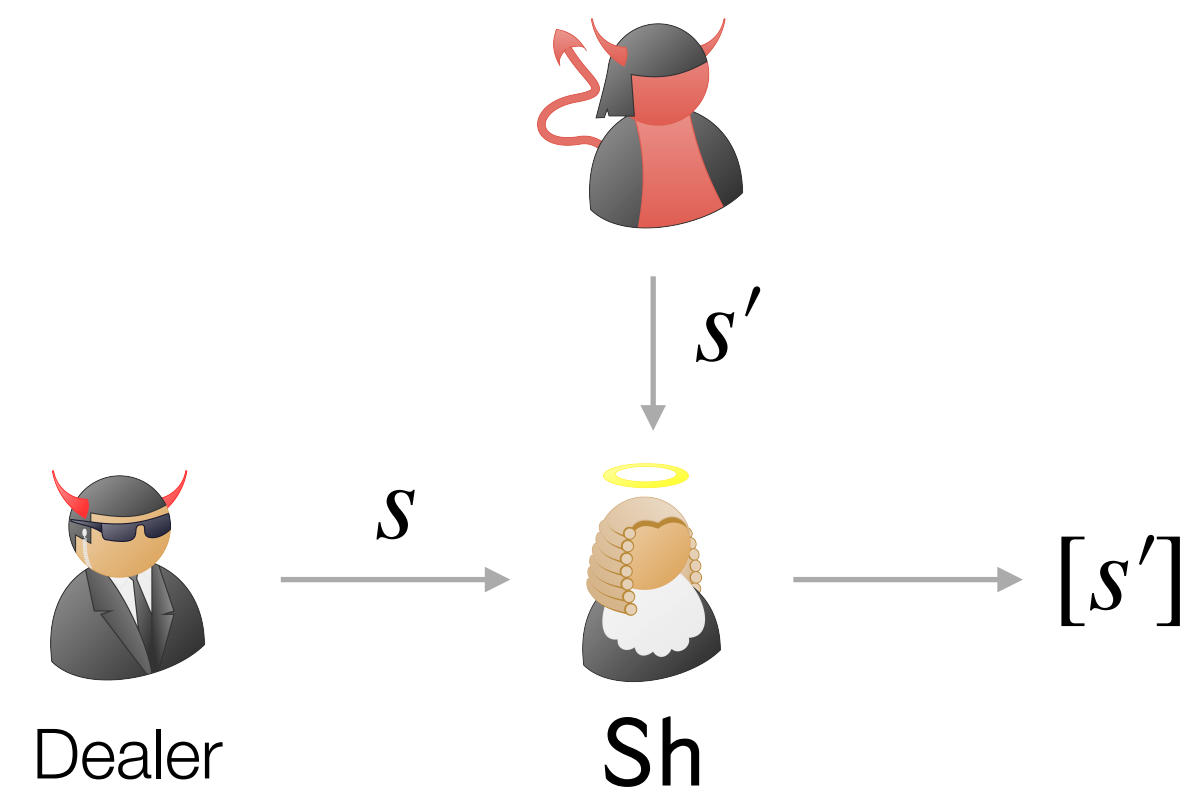
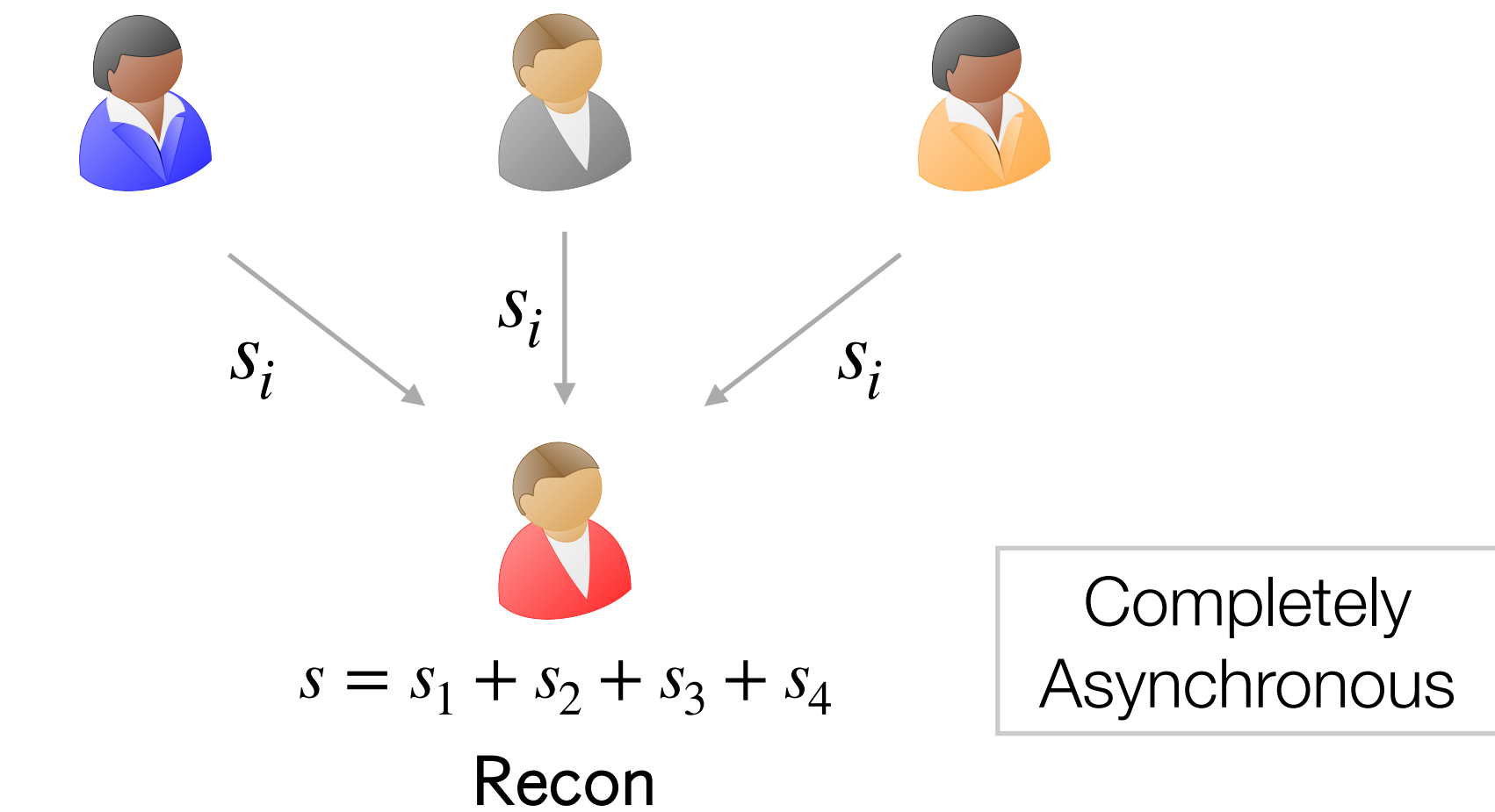
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Perfect HMPC - VSS with Party Elimination

- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



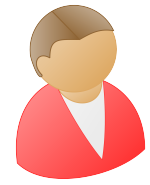
P_1



P_2



P_3

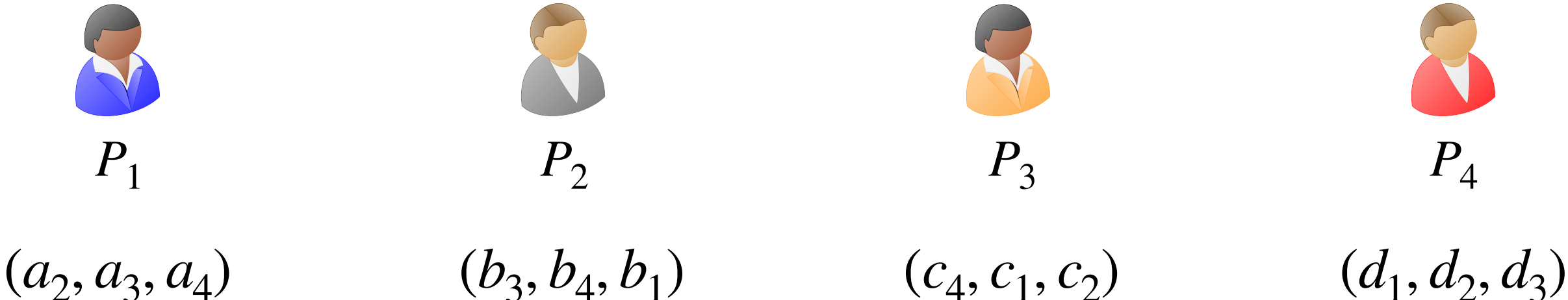


P_4

Perfect HMPC - VSS with Party Elimination

- Round 1

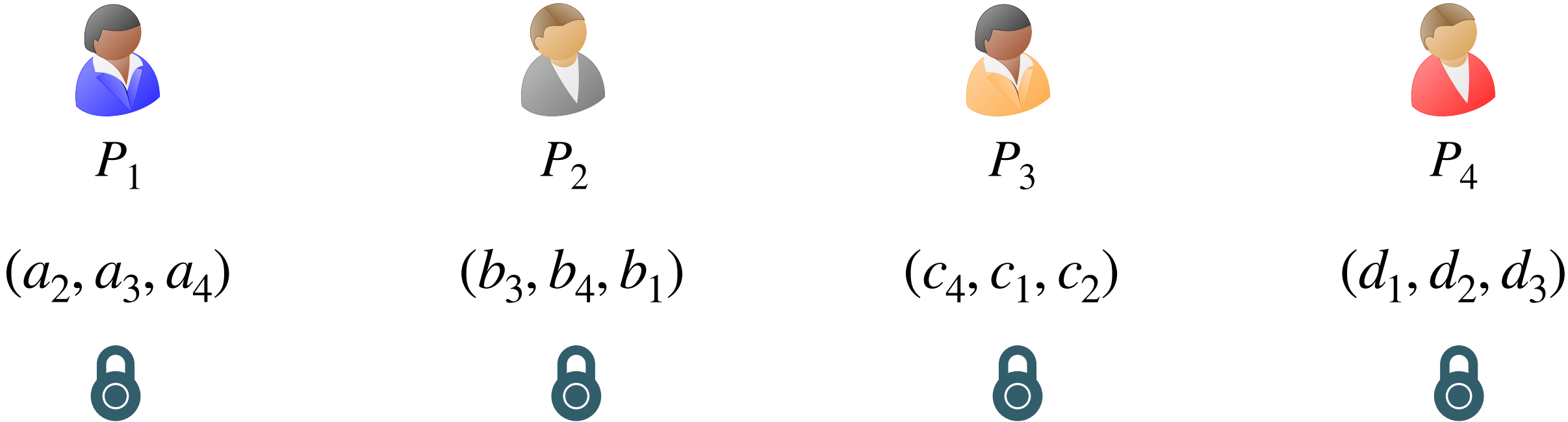
- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



Perfect HMPC - VSS with Party Elimination

- Round 1

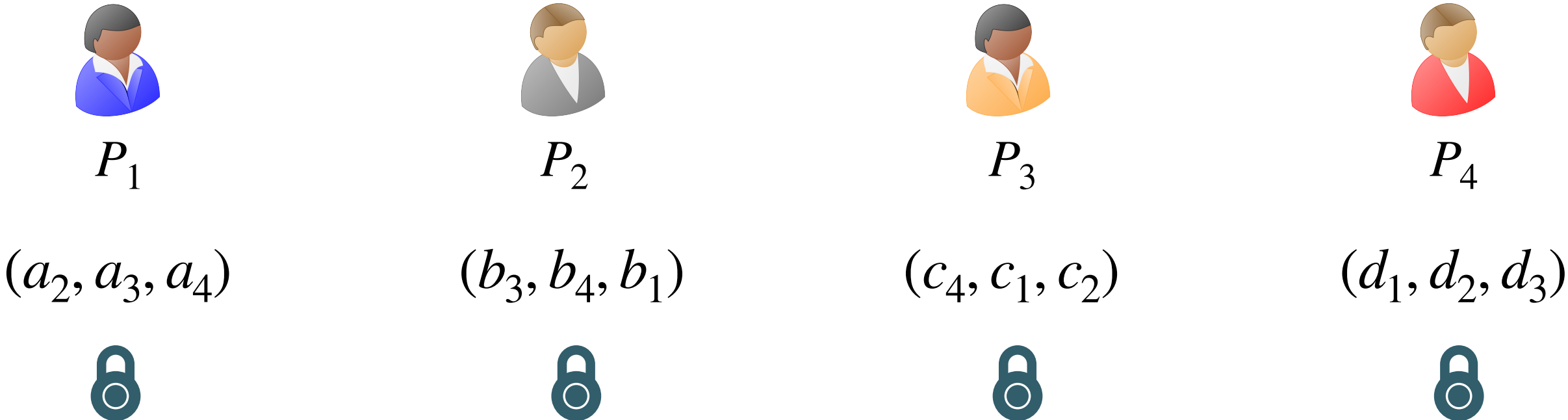
- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



Perfect HMPC - VSS with Party Elimination

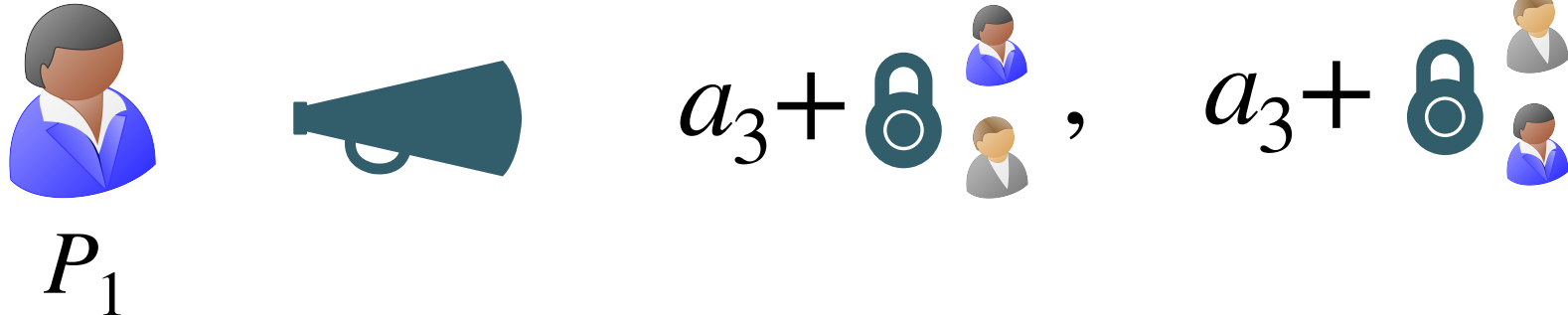
- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



- Round 2

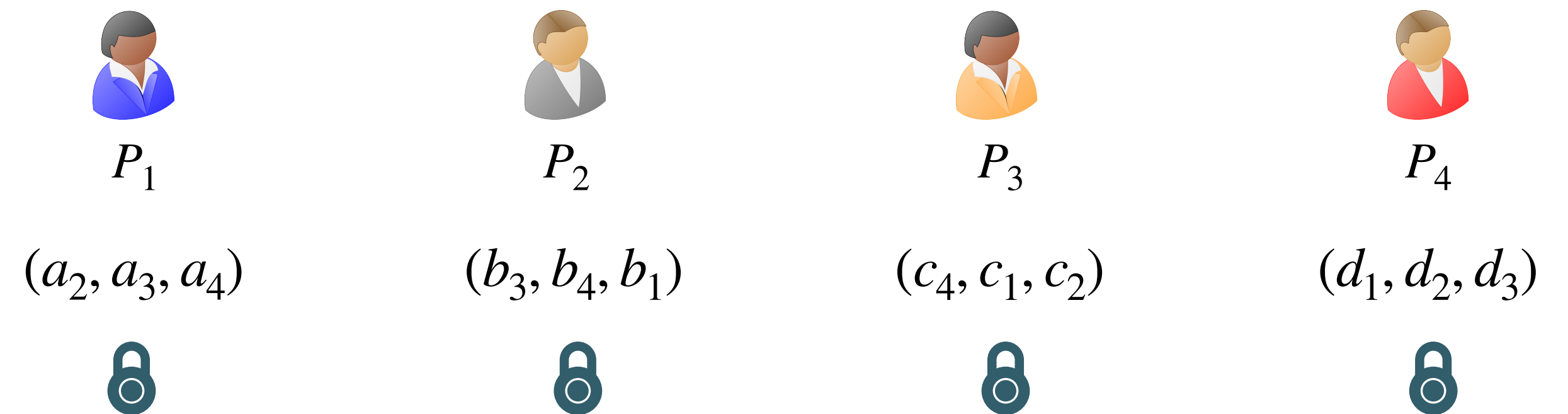
- Parties broadcast **masked** shares



Perfect HMPC - VSS with Party Elimination

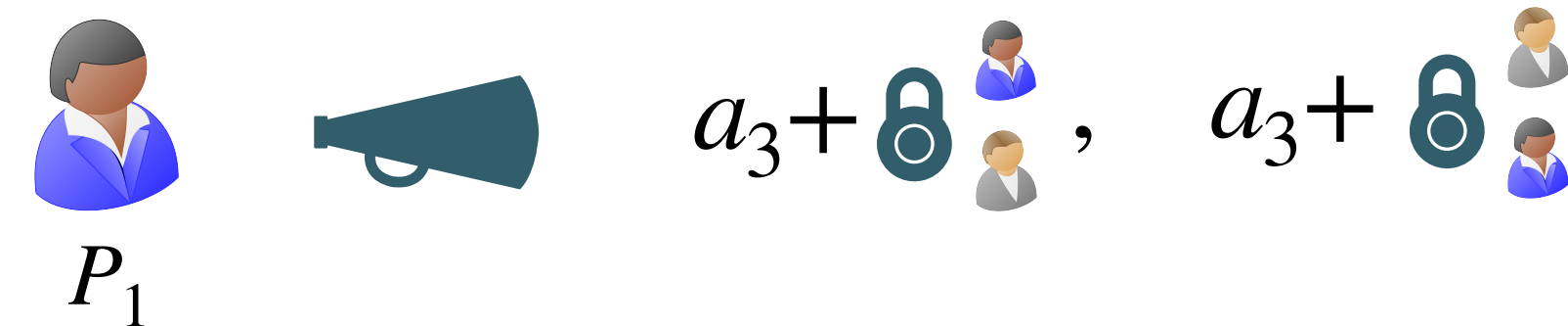
- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



- Round 2

- Parties broadcast **masked** shares



- Local computation

- If shares inconsistent, output dispute set
- Else output with secret shares

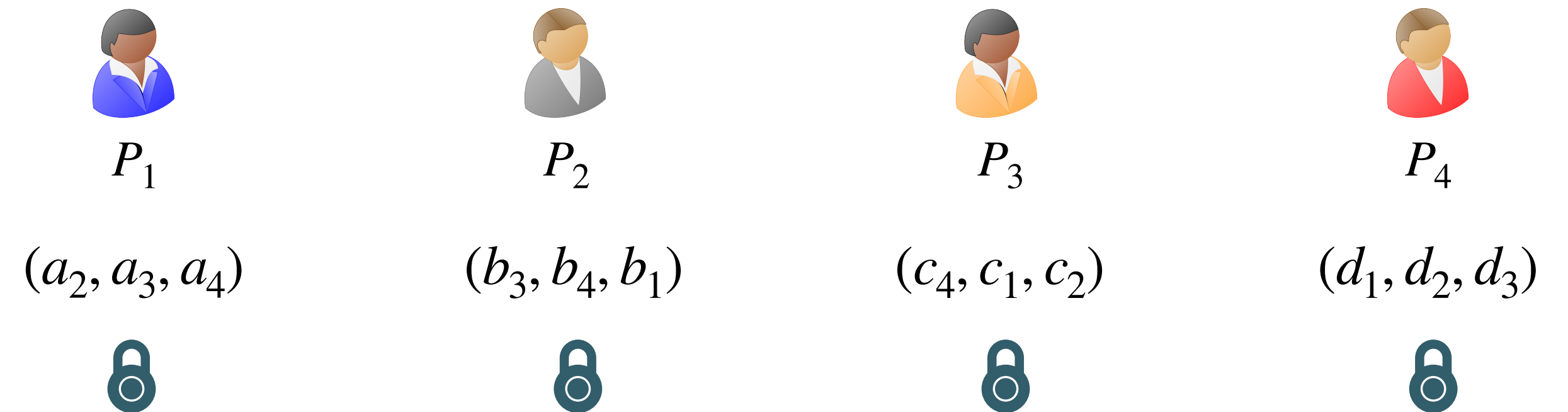
$$a_3 + \text{lock} + \text{party} \neq b_3 + \text{lock} + \text{party}$$

$$a_3 + \text{lock} + \text{party} \neq b_3 + \text{lock} + \text{party}$$

Perfect HMPC - VSS with Party Elimination

- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



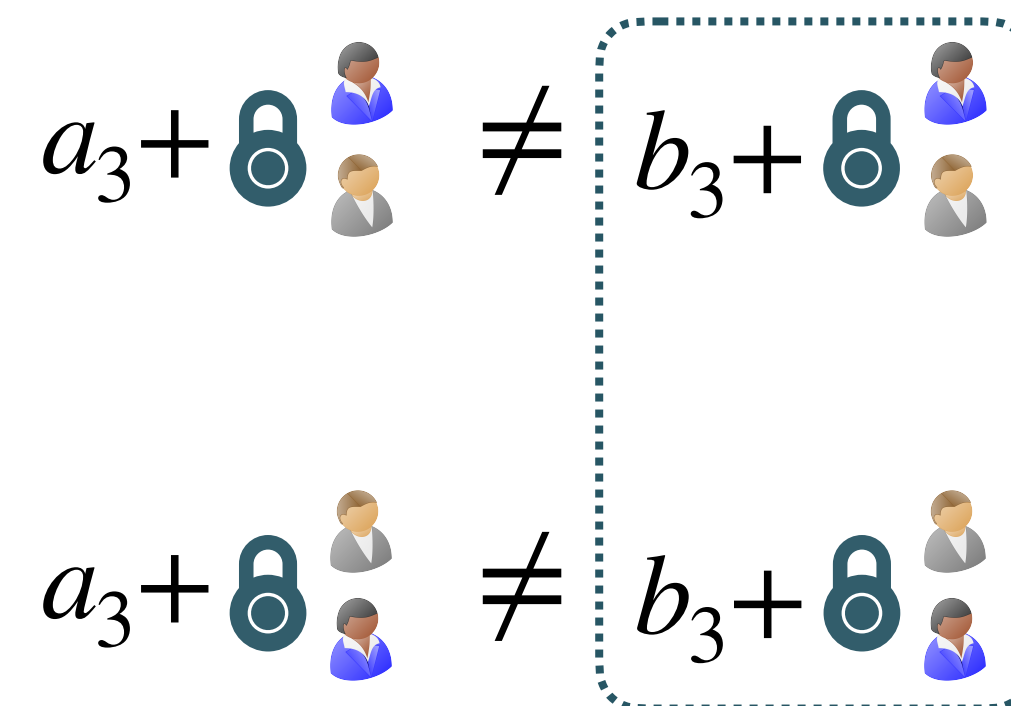
- Round 2

- Parties broadcast **masked** shares



- Local computation

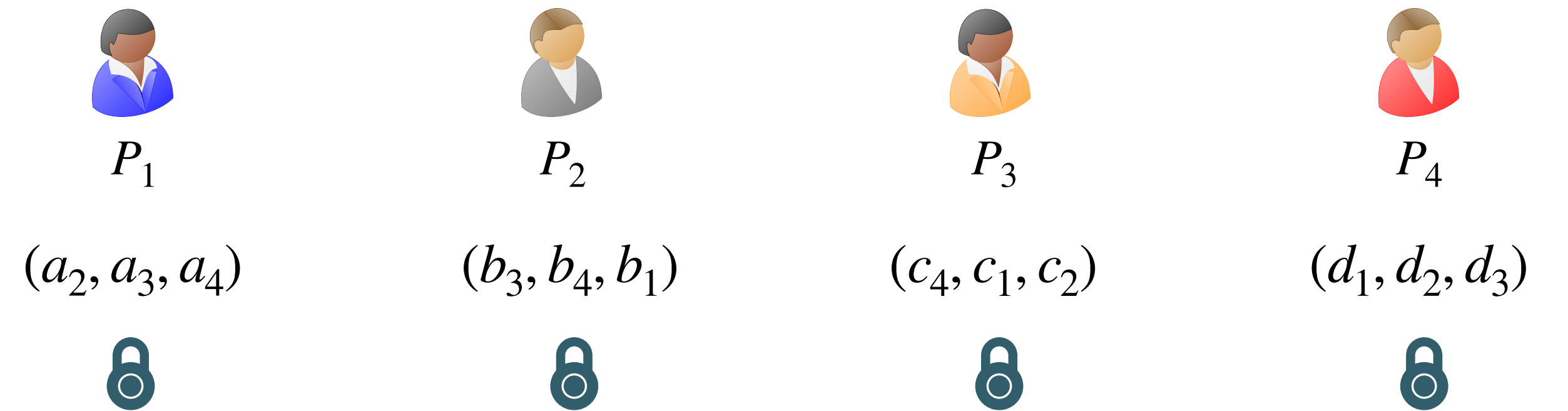
- If shares inconsistent, output dispute set
- Else output with secret shares



Perfect HMPC - VSS with Party Elimination

- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



- Round 2

- Parties broadcast **masked** shares



- Local computation

- If shares inconsistent, output dispute set
- Else output with secret shares

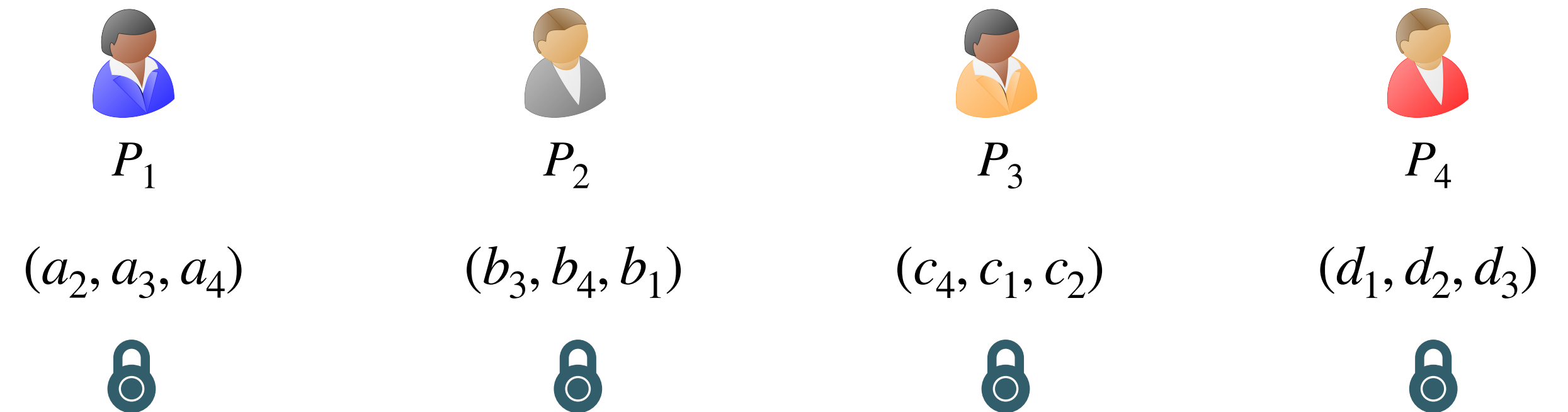
$$a_3 + \text{lock} + P_1 \neq b_3 + \text{lock} + P_1$$

$$a_3 + \text{lock} + P_2 \neq b_3 + \text{lock} + P_2$$

Perfect HMPC - VSS with Party Elimination

- Round 1

- D sends share to each party
- Parties exchange **random pad** for each **common element** in share



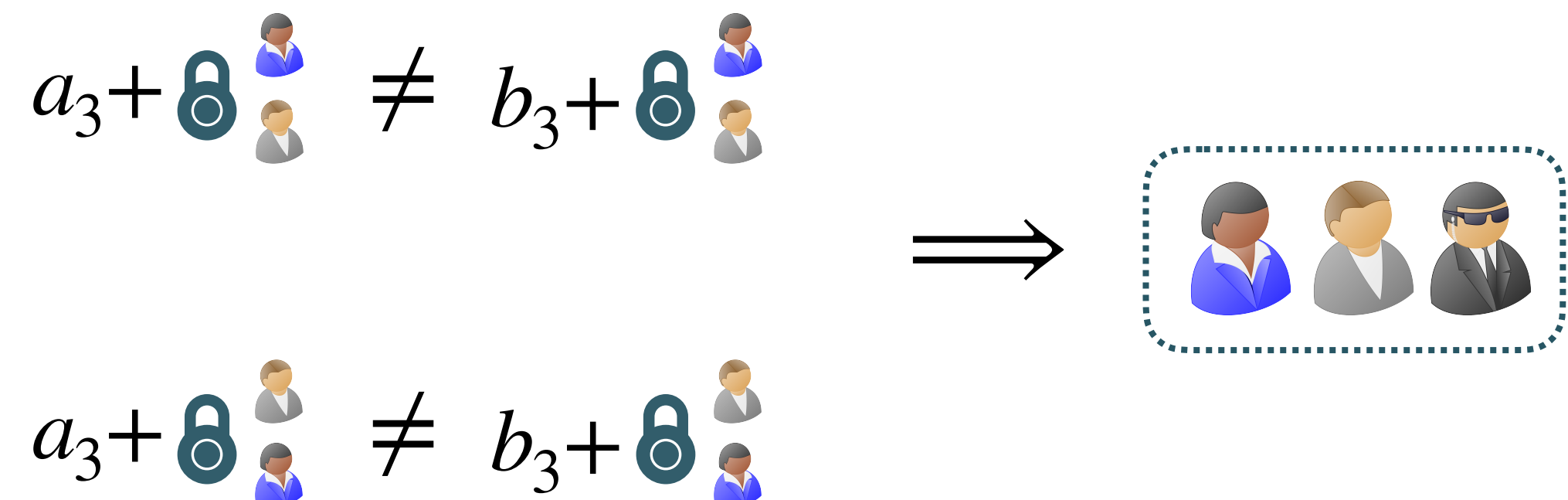
- Round 2

- Parties broadcast **masked** shares

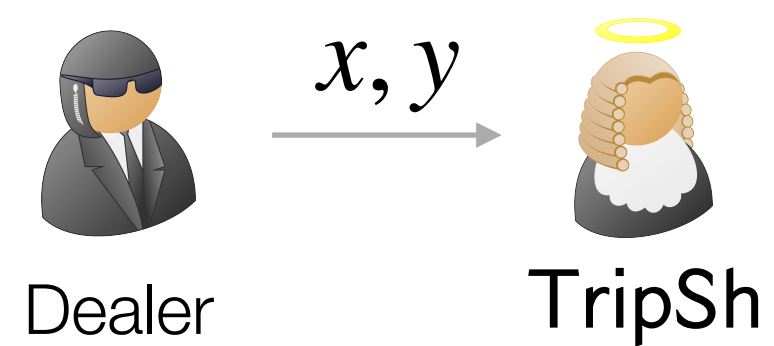


- Local computation

- If shares inconsistent, output dispute set
- Else output with secret shares

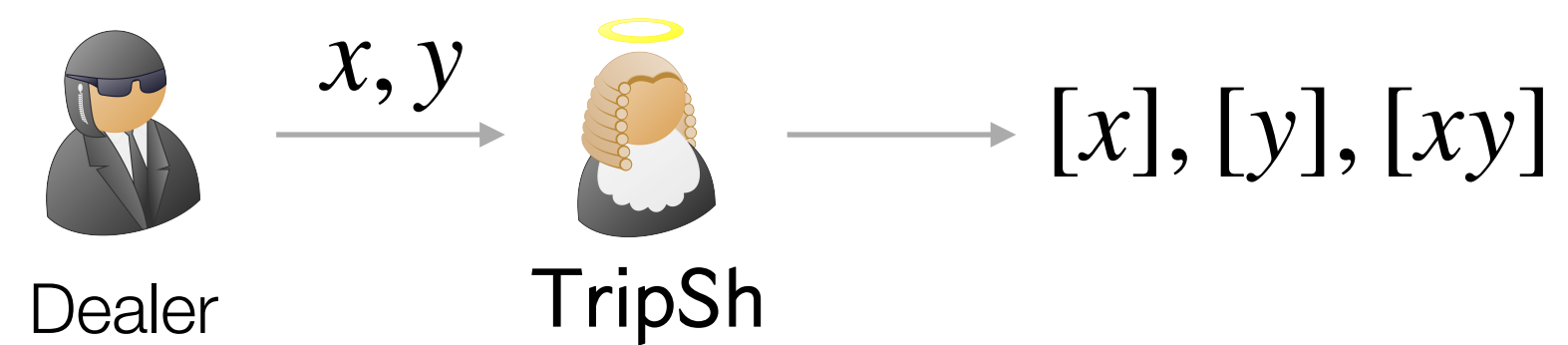


Perfect HMPC - Triple Sharing with Party Elimination Functionality



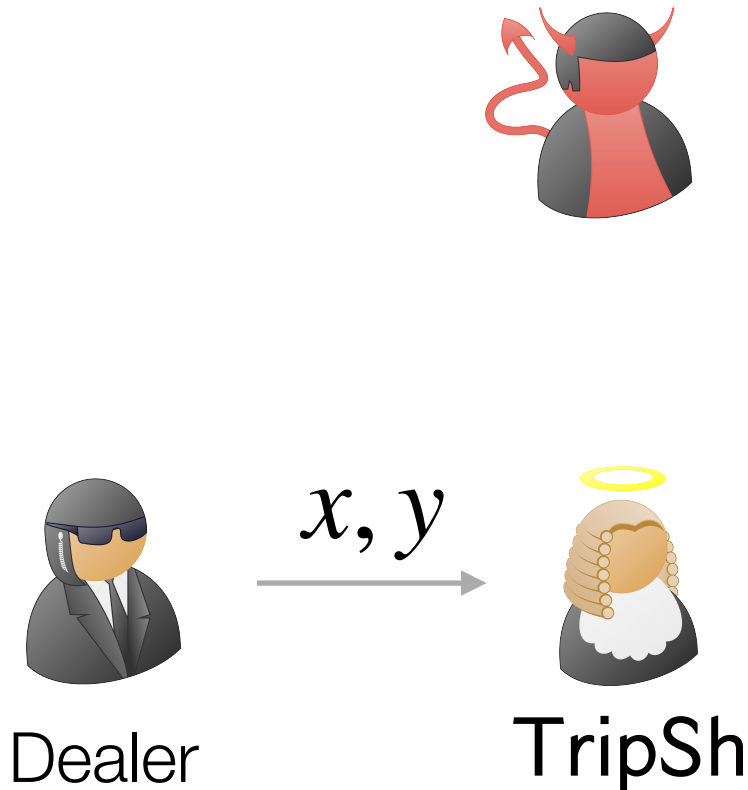
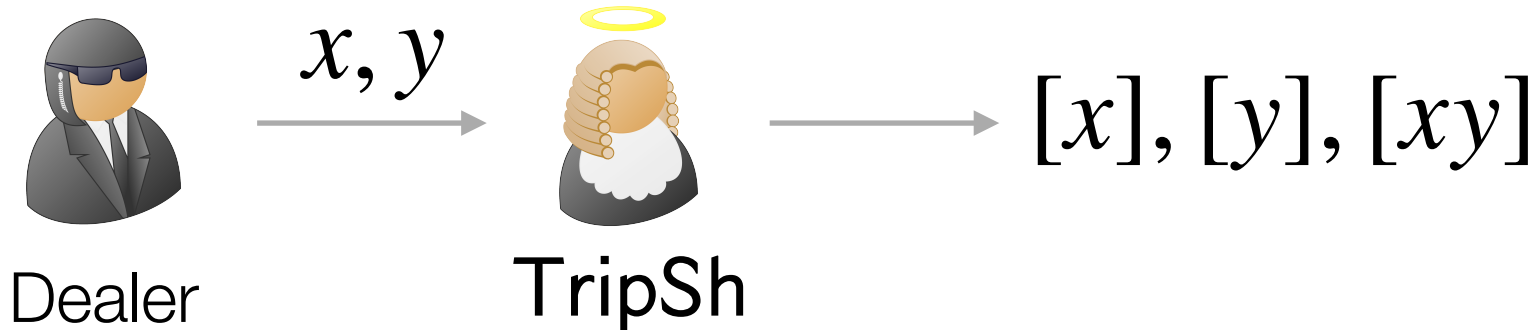
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Sharing with Party Elimination Functionality



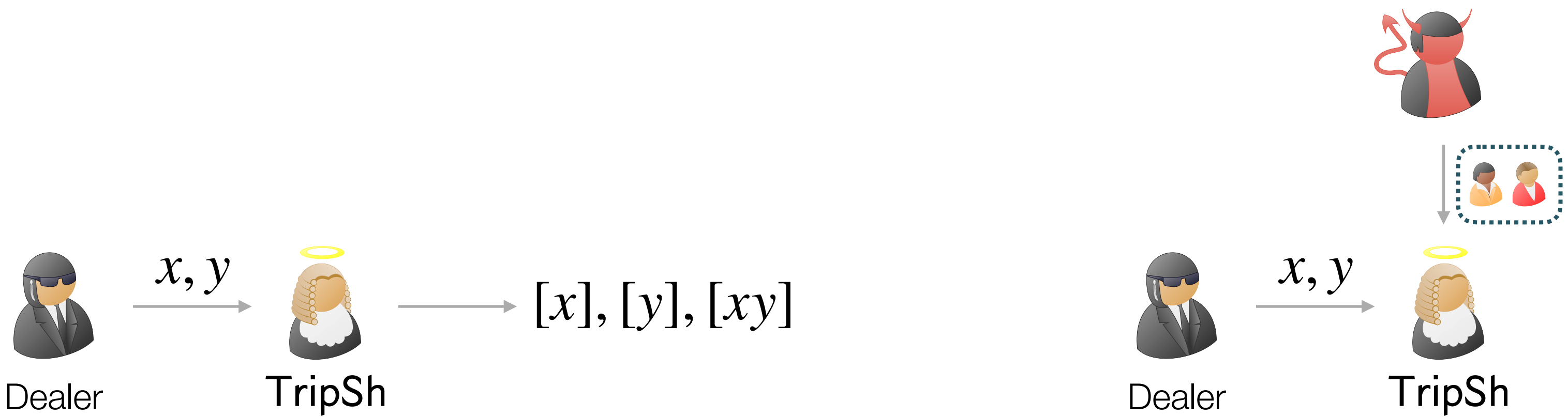
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Sharing with Party Elimination Functionality



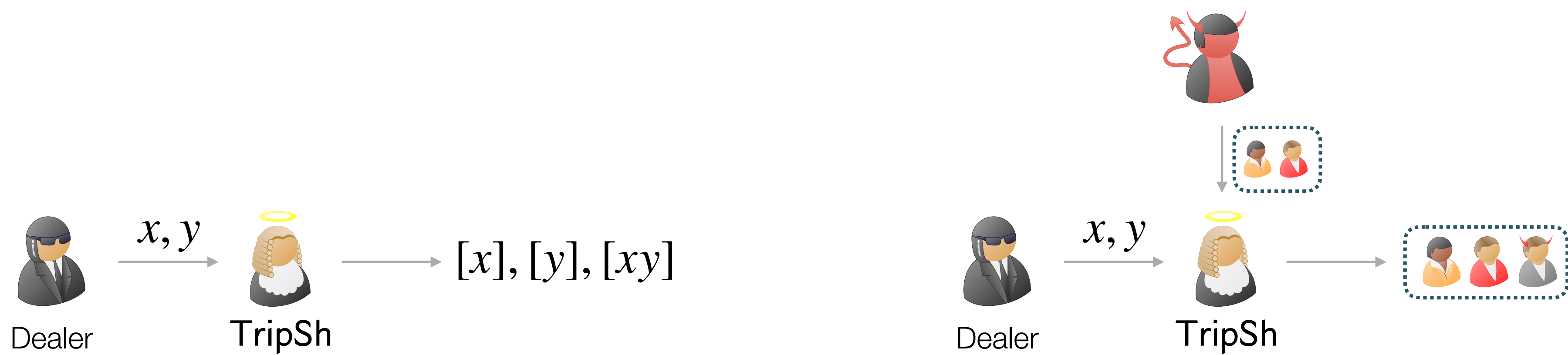
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Sharing with Party Elimination Functionality



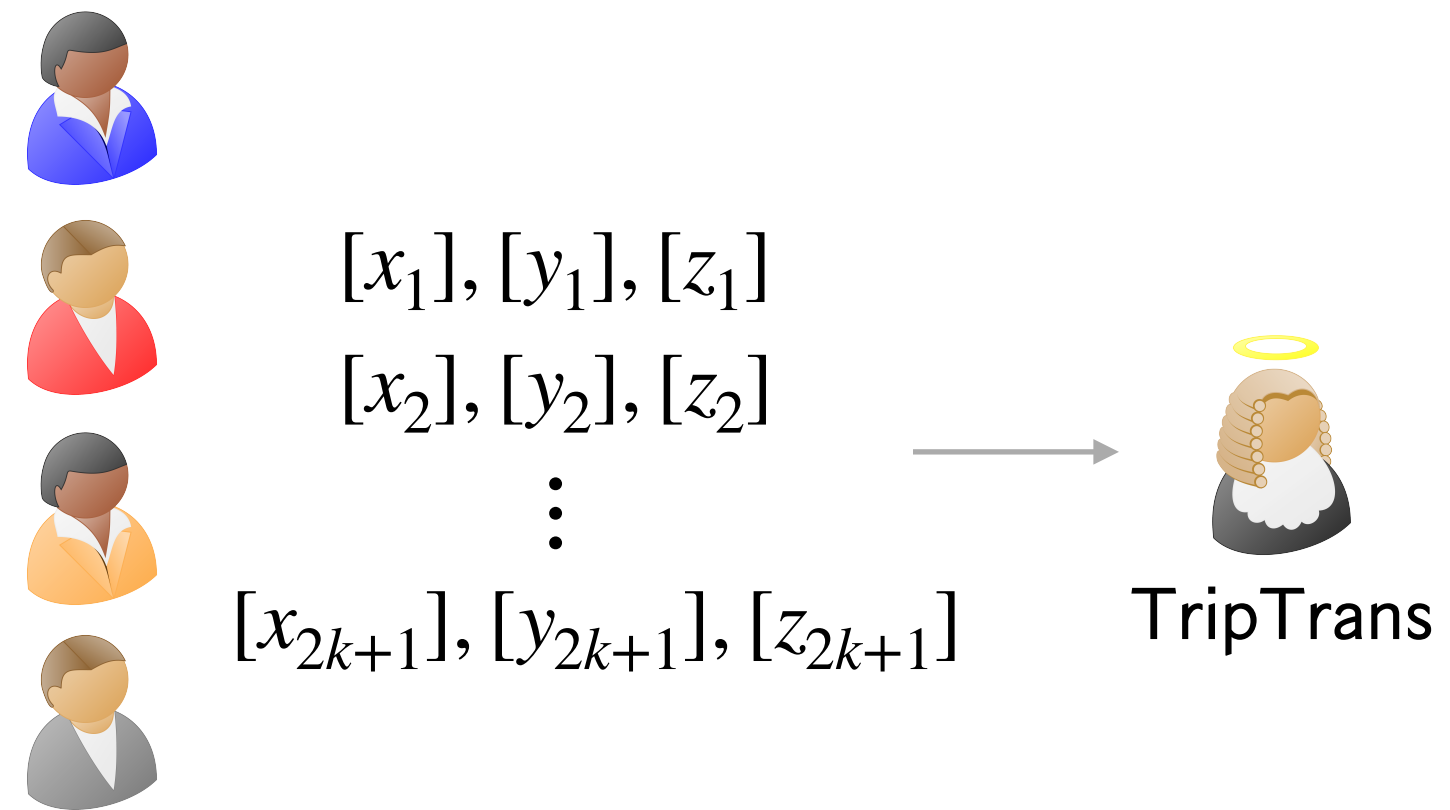
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Sharing with Party Elimination Functionality



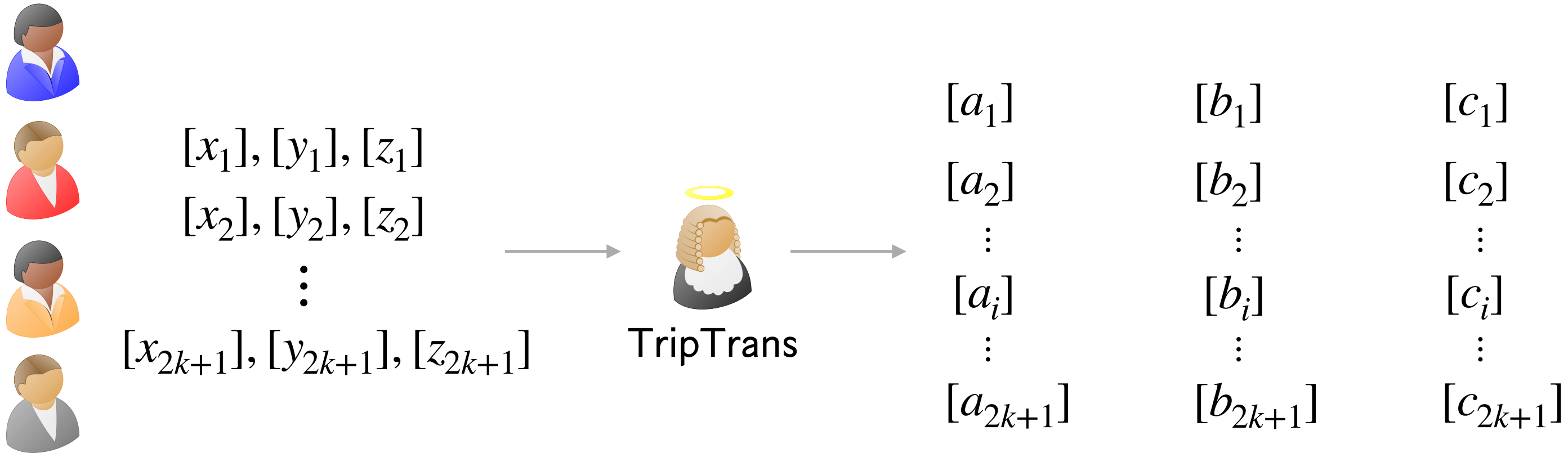
- Triple sharing with Party Elimination
 - Verified multiplication triple or dispute set

Perfect HMPC - Triple Transform Functionality



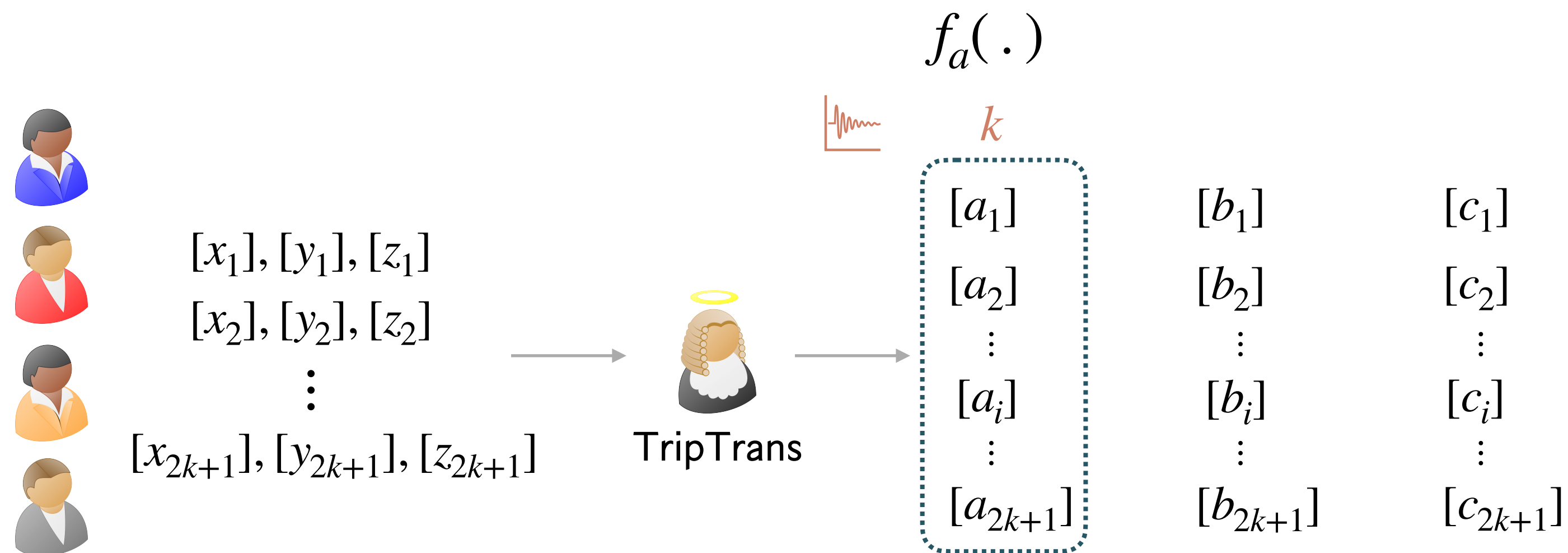
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



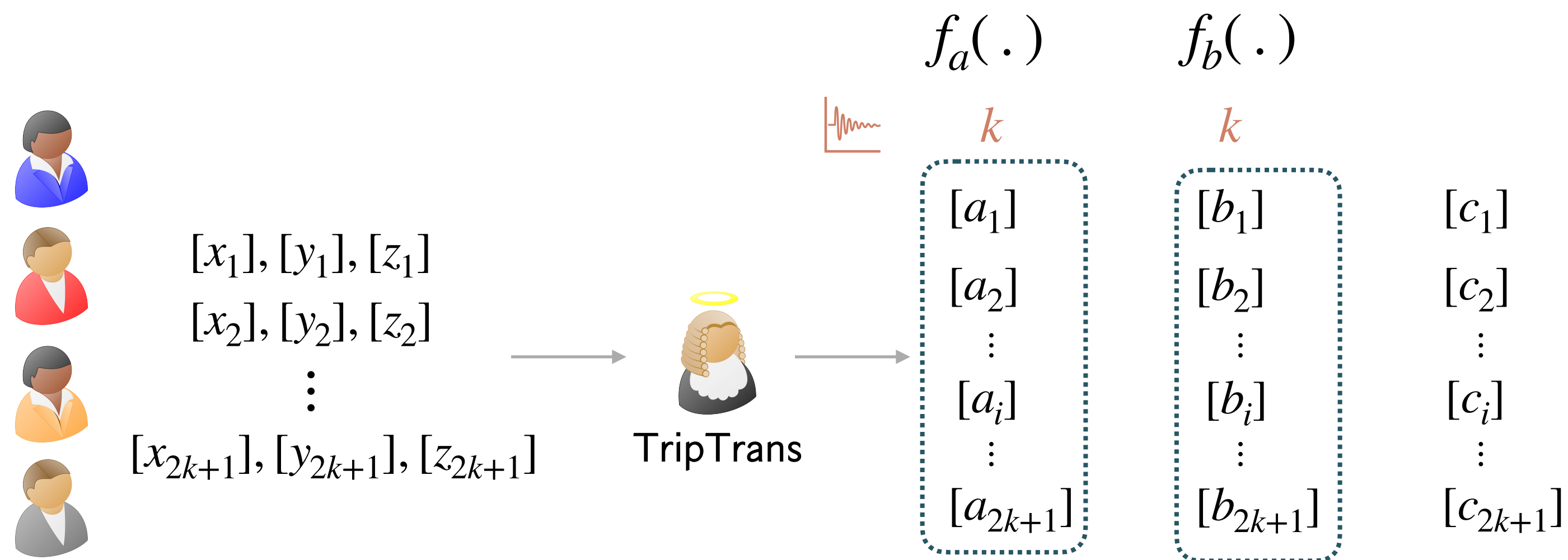
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



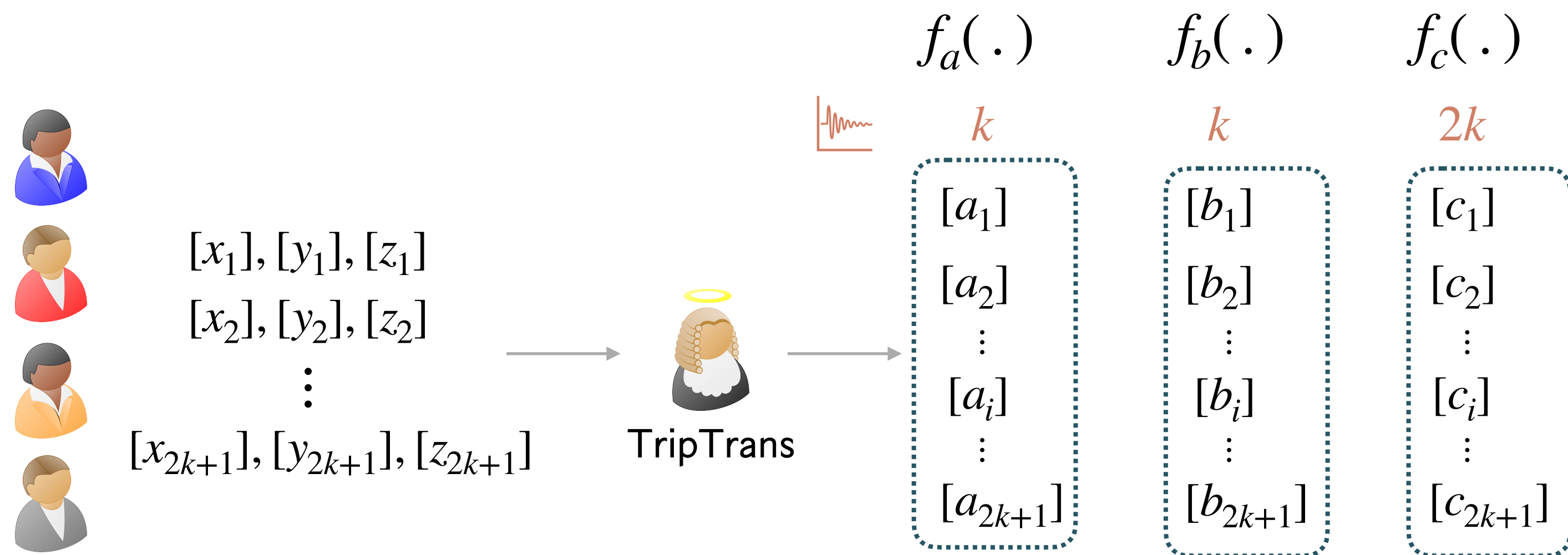
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



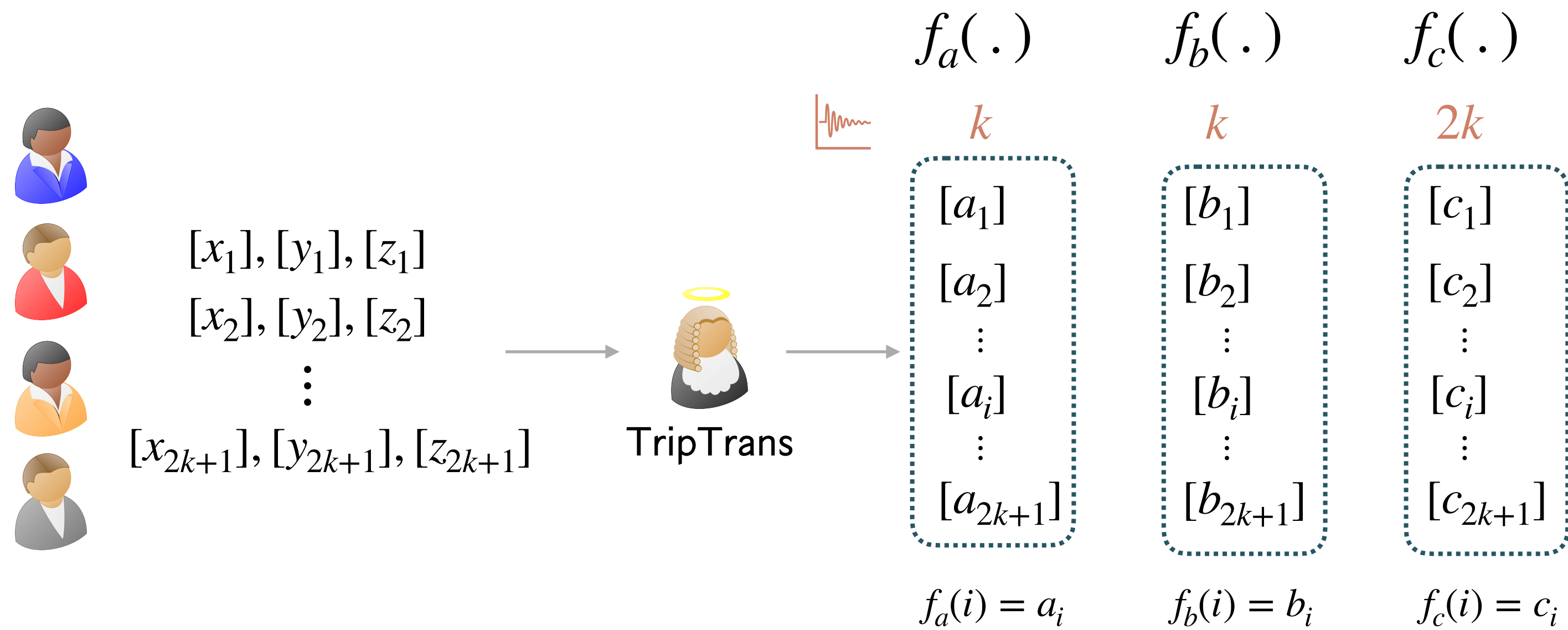
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



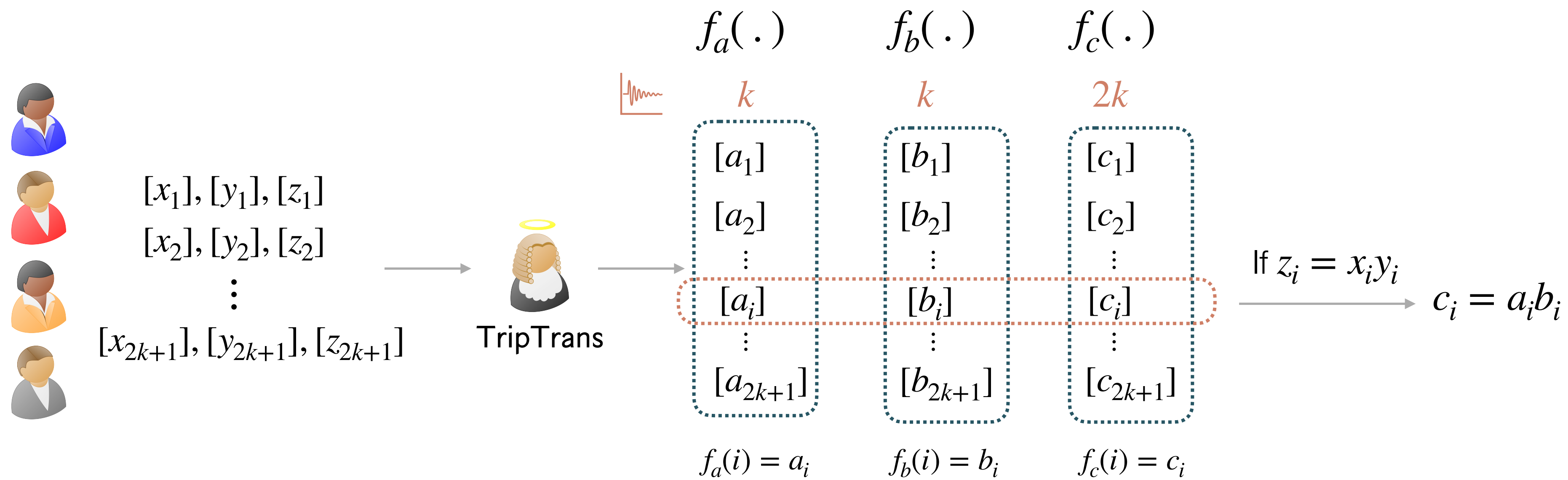
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



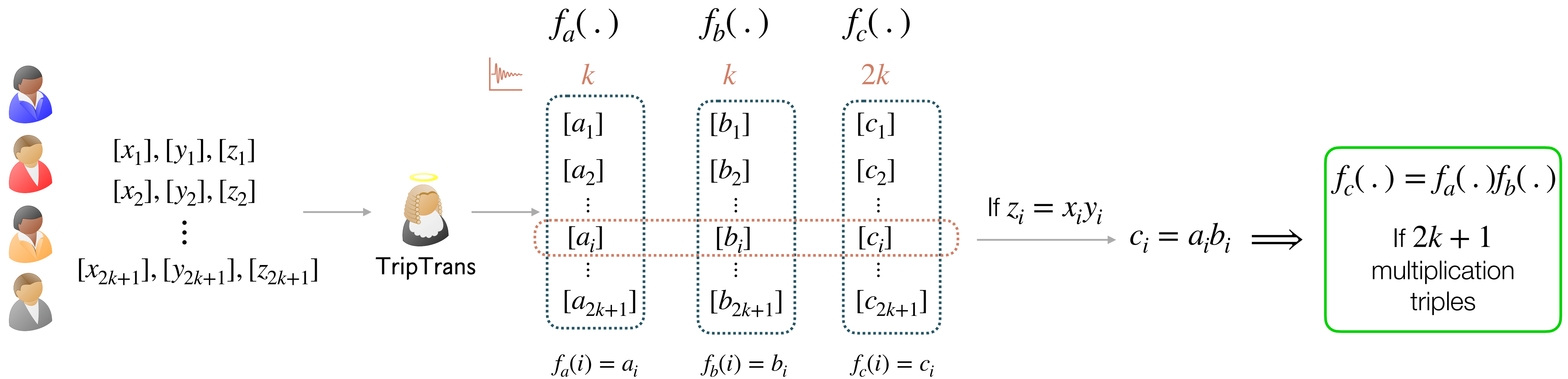
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



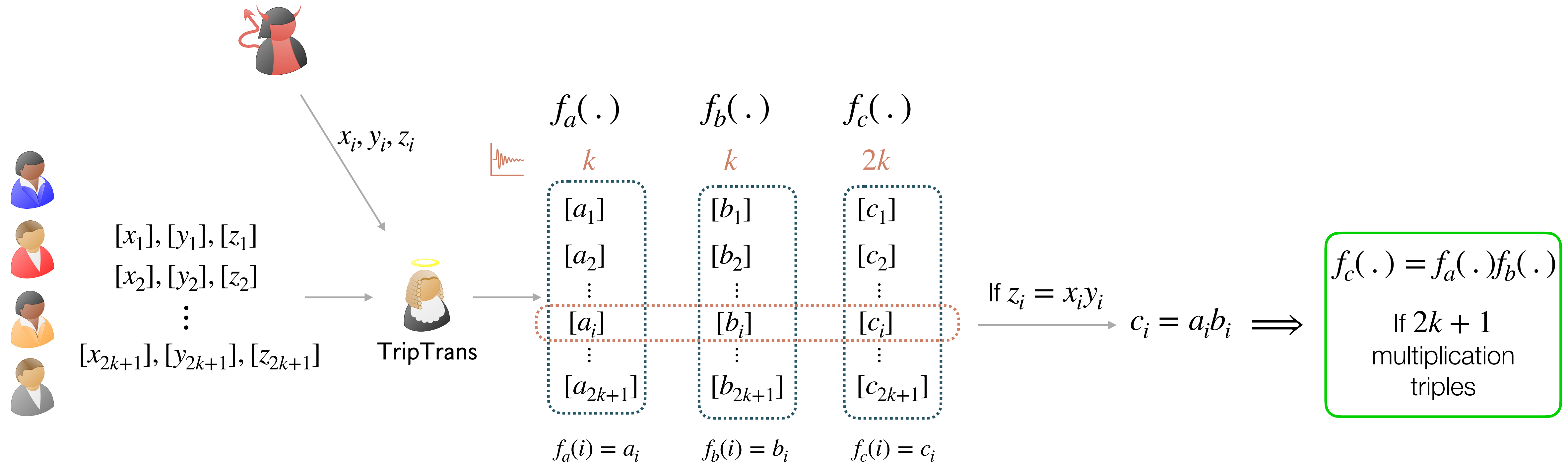
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



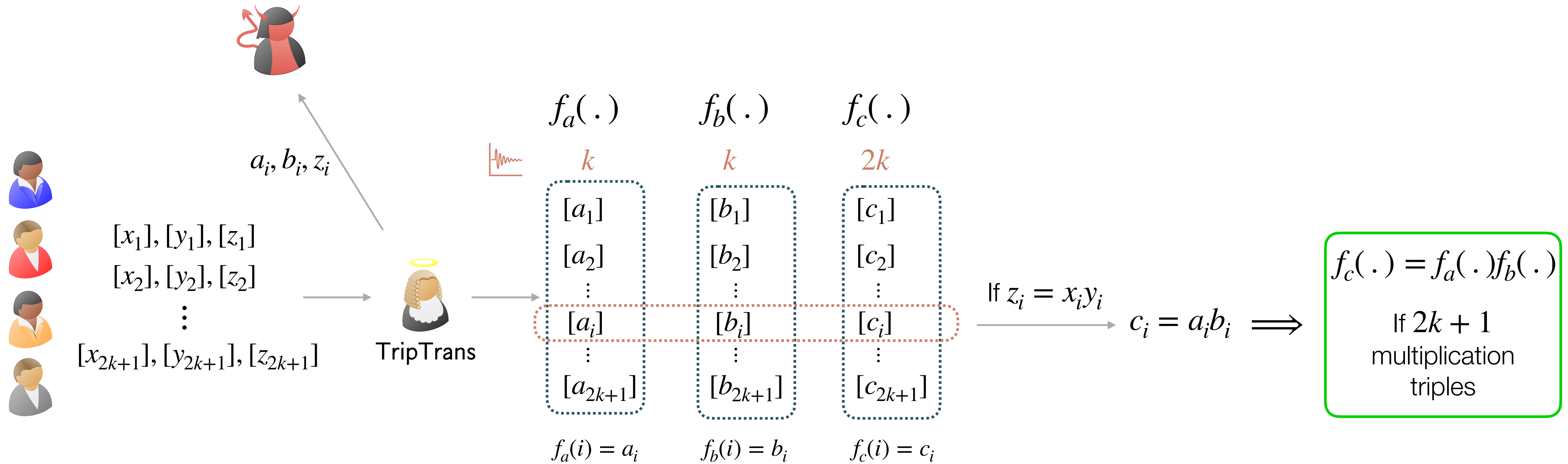
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



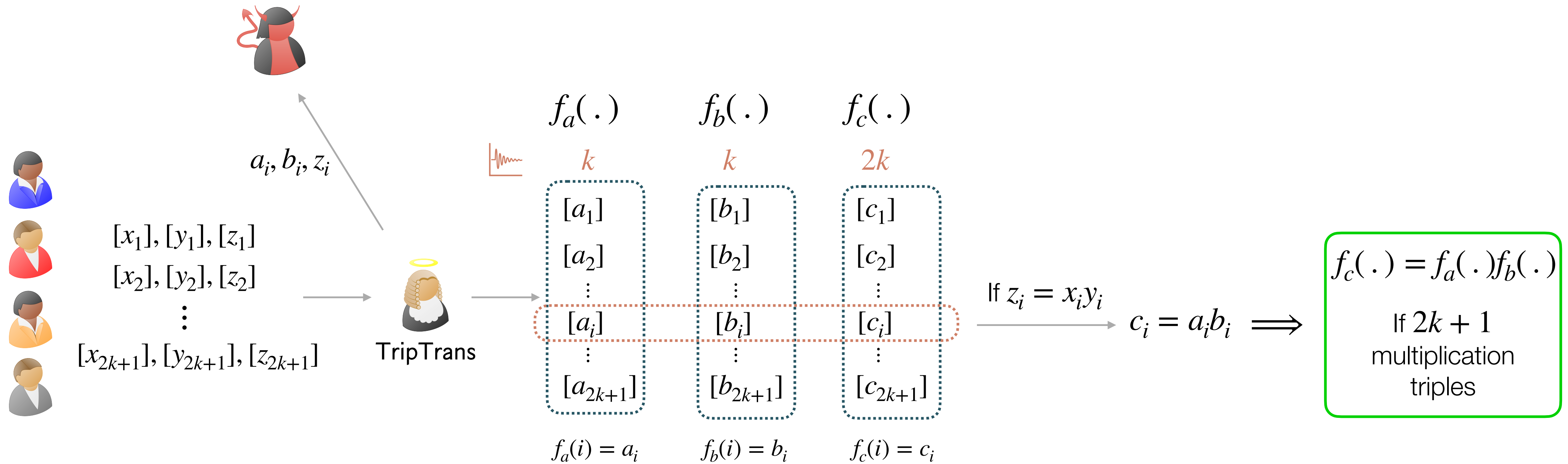
- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality



- Random triples \rightarrow correlated random triples

Perfect HMPC - Triple Transform Functionality

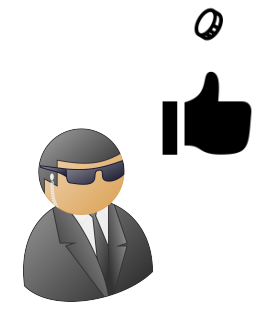


- Random triples \rightarrow correlated random triples
- Completely asynchronous instantiation in [CP17]

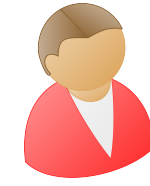
Perfect HMPC - Triple Sharing with Party Elimination Protocol



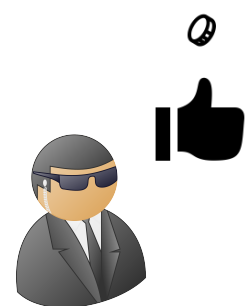
Perfect HMPC - Triple Sharing with Party Elimination Protocol



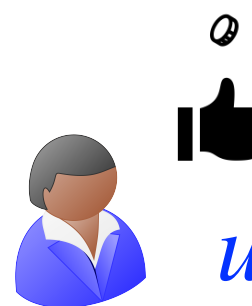
| | | |
|-------|-------|----------|
| x_1 | y_1 | x_1y_1 |
| x_2 | y_2 | x_2y_2 |
| x_3 | y_3 | x_3y_3 |



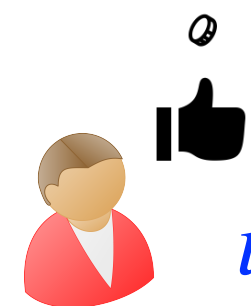
Perfect HMPC - Triple Sharing with Party Elimination Protocol



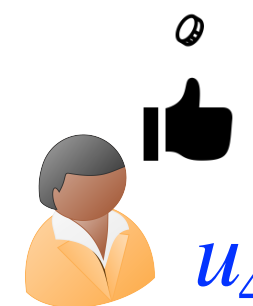
x_1 y_1 x_1y_1
 x_2 y_2 x_2y_2
 x_3 y_3 x_3y_3
 u_1 v_1 u_1v_1



u_2 v_2 u_2v_2

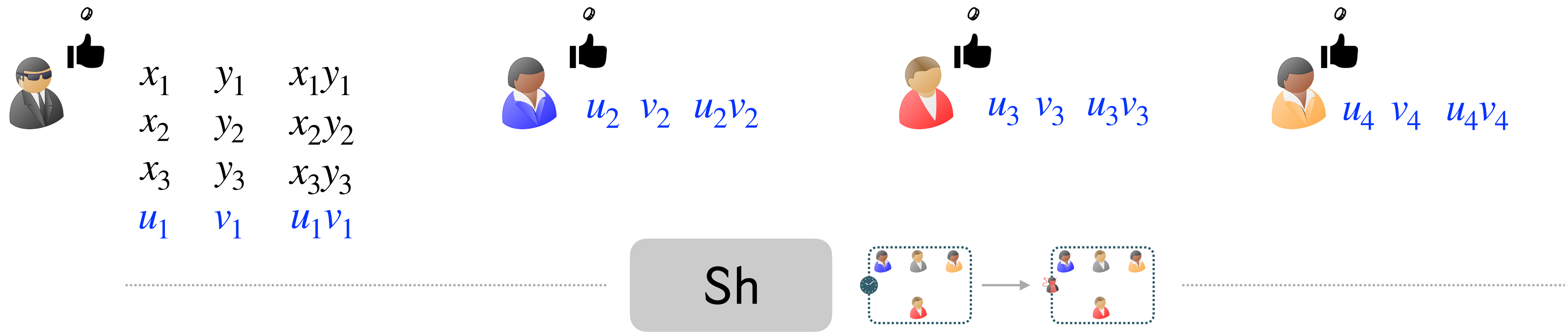


u_3 v_3 u_3v_3



u_4 v_4 u_4v_4

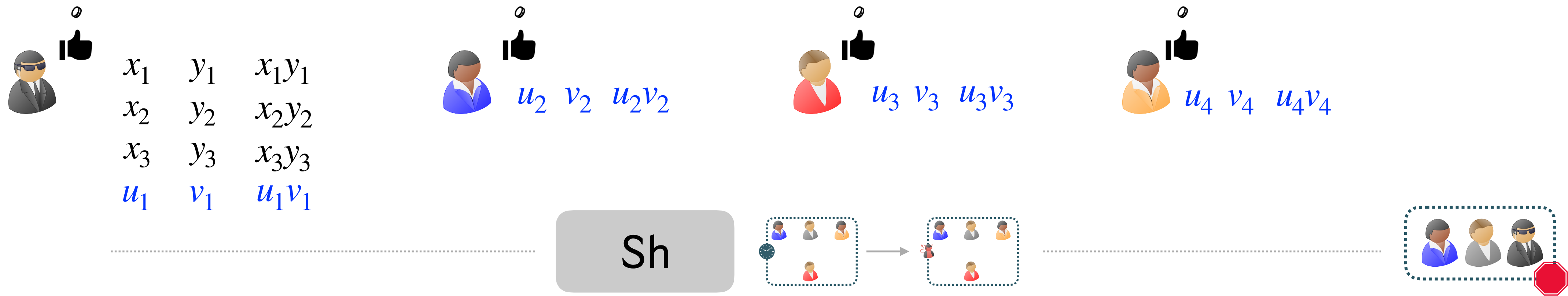
Perfect HMPC - Triple Sharing with Party Elimination Protocol



$[x_1] \ [y_1] \ [z_1]$
 $[x_2] \ [y_2] \ [z_2]$
 $[x_3] \ [y_3] \ [z_3]$

$[u_1] \ [v_1] \ [w_1]$
 $[u_2] \ [v_2] \ [w_2]$
 $[u_3] \ [v_3] \ [w_3]$
 $[u_4] \ [v_4] \ [w_4]$

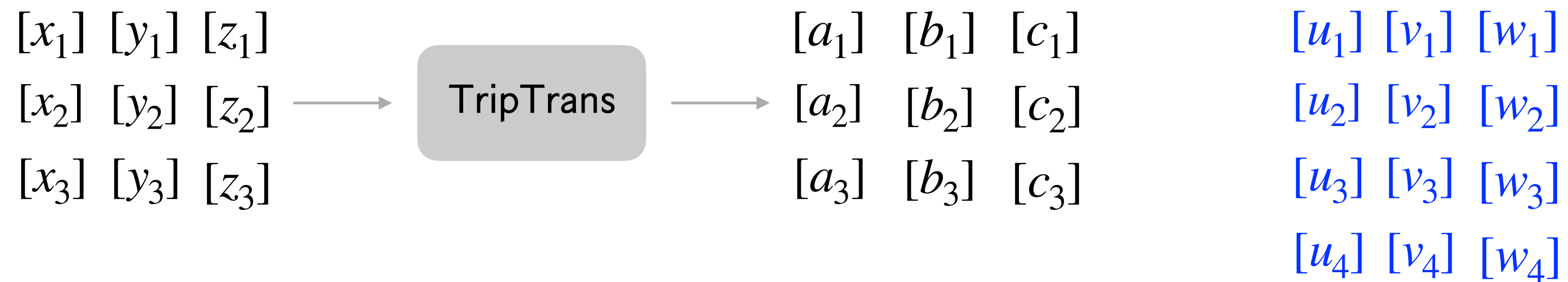
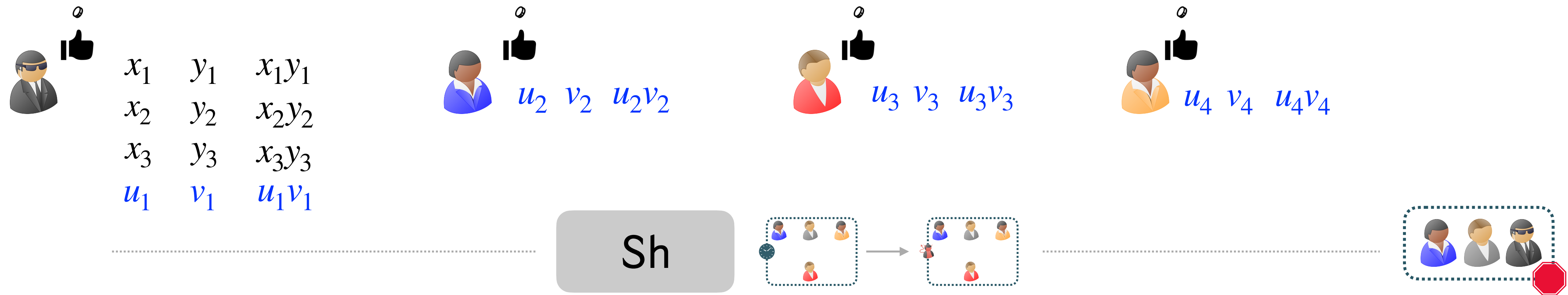
Perfect HMPC - Triple Sharing with Party Elimination Protocol



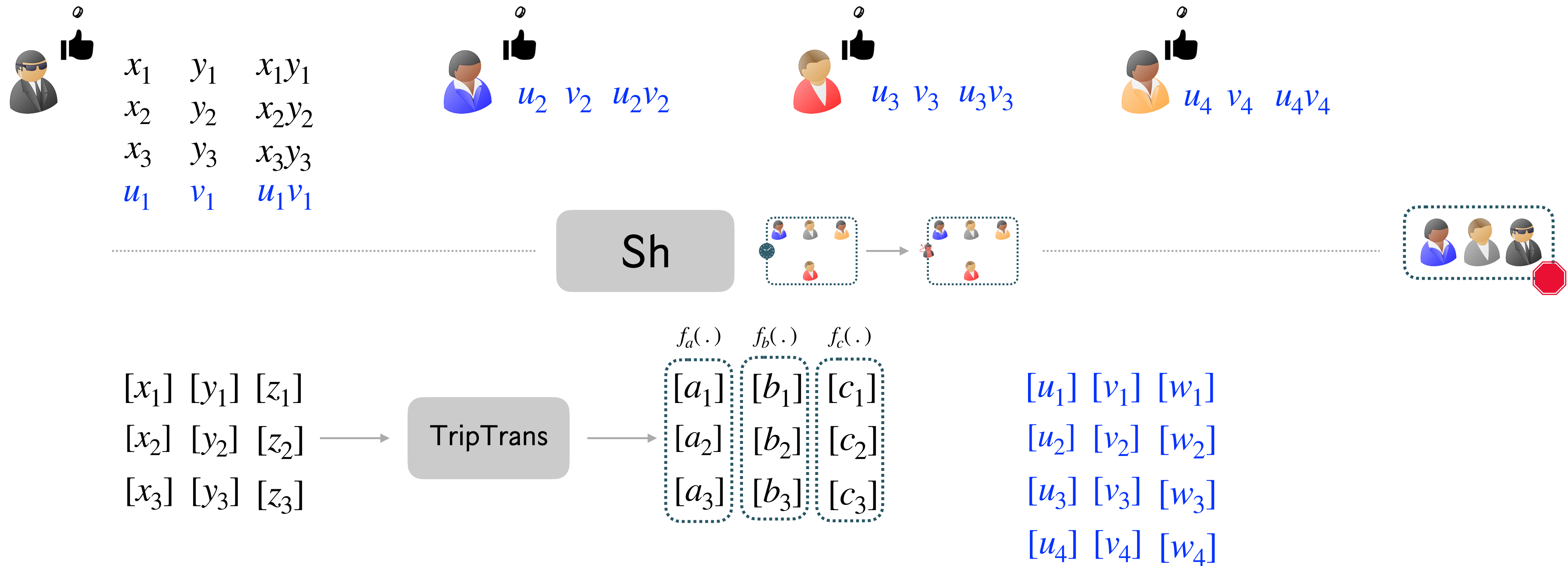
$[x_1] [y_1] [z_1]$
 $[x_2] [y_2] [z_2]$
 $[x_3] [y_3] [z_3]$

$[u_1] [v_1] [w_1]$
 $[u_2] [v_2] [w_2]$
 $[u_3] [v_3] [w_3]$
 $[u_4] [v_4] [w_4]$

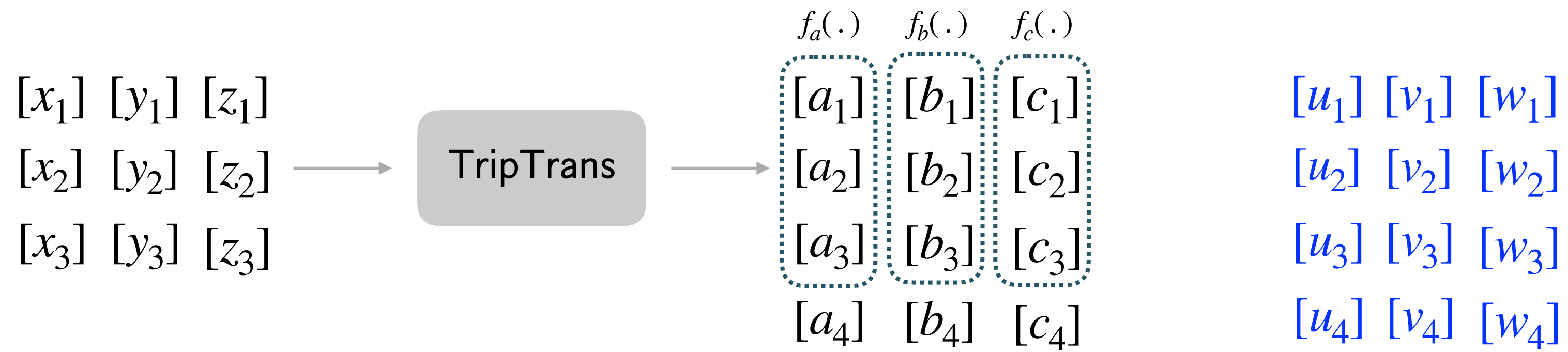
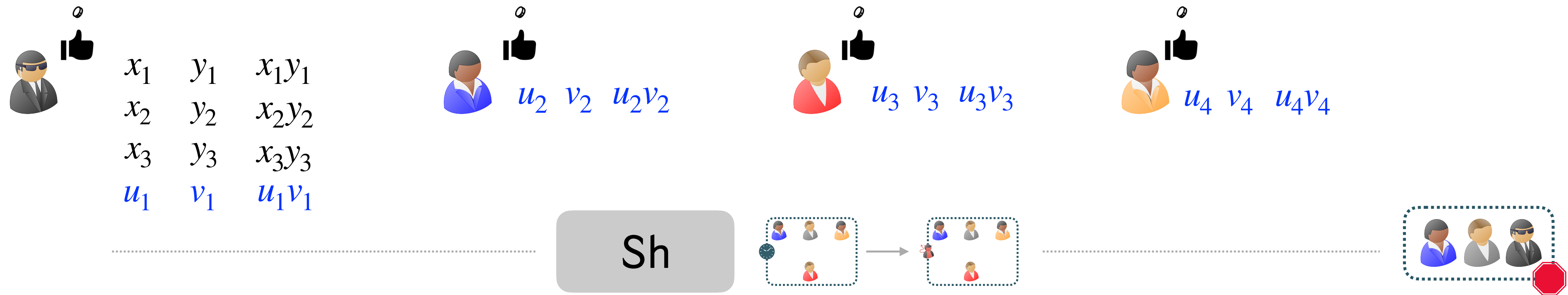
Perfect HMPC - Triple Sharing with Party Elimination Protocol



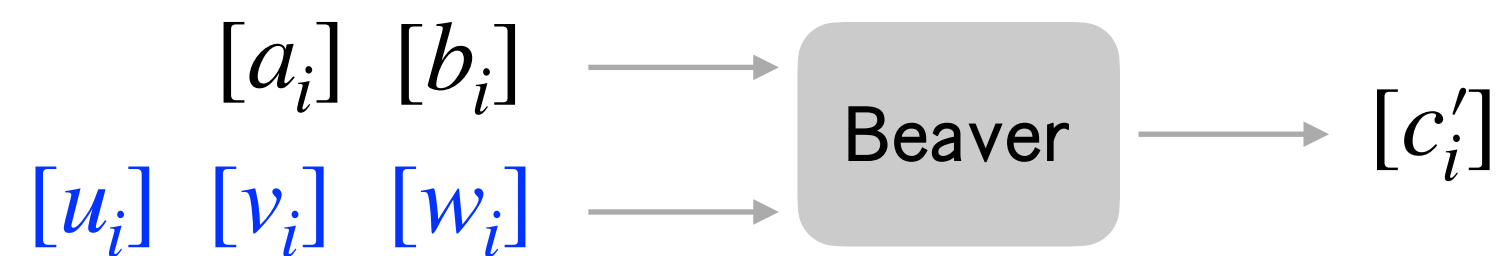
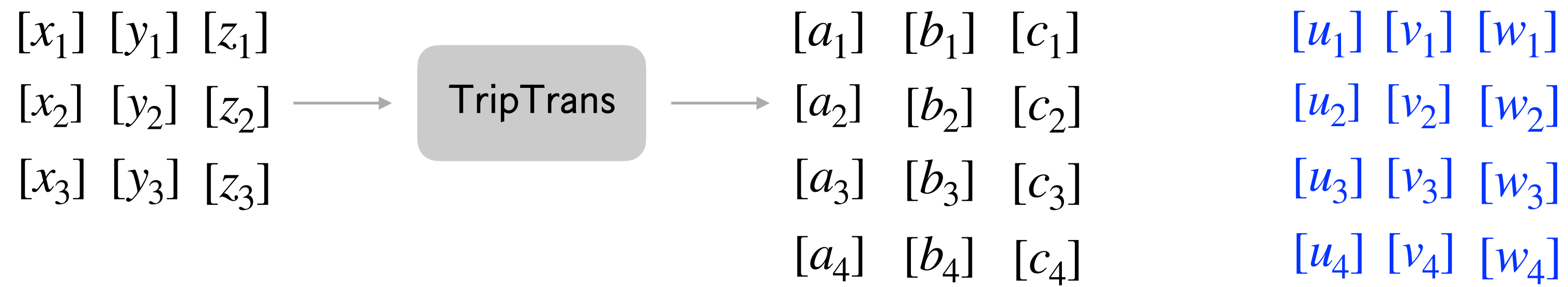
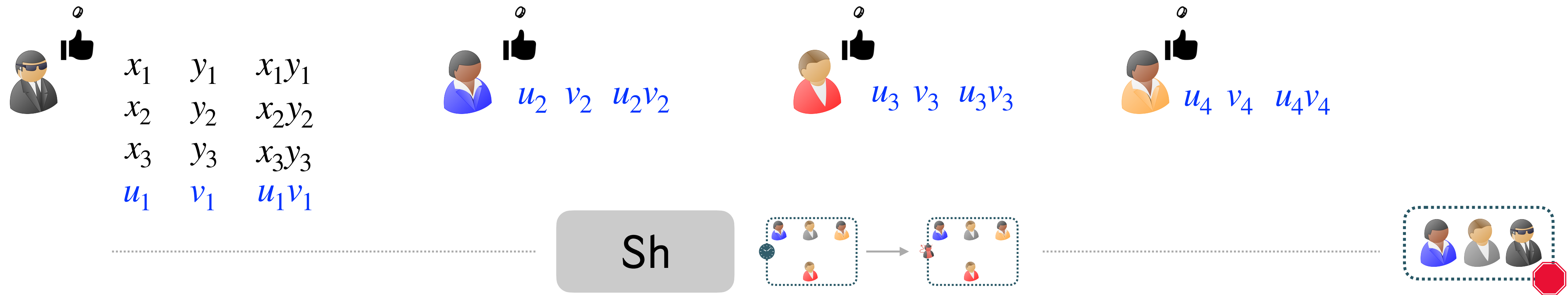
Perfect HMPC - Triple Sharing with Party Elimination Protocol



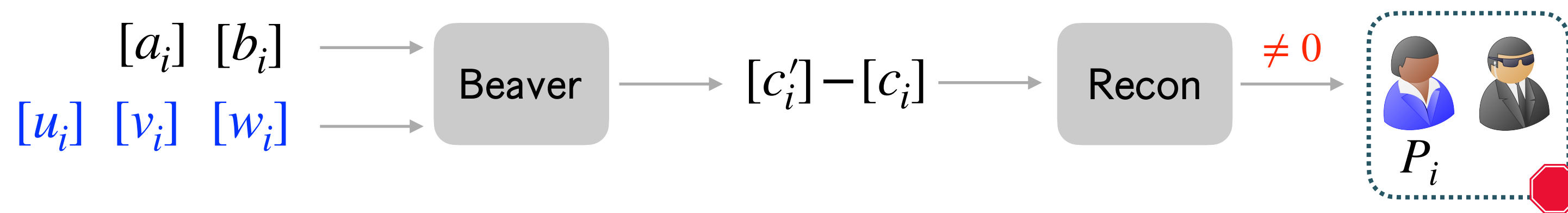
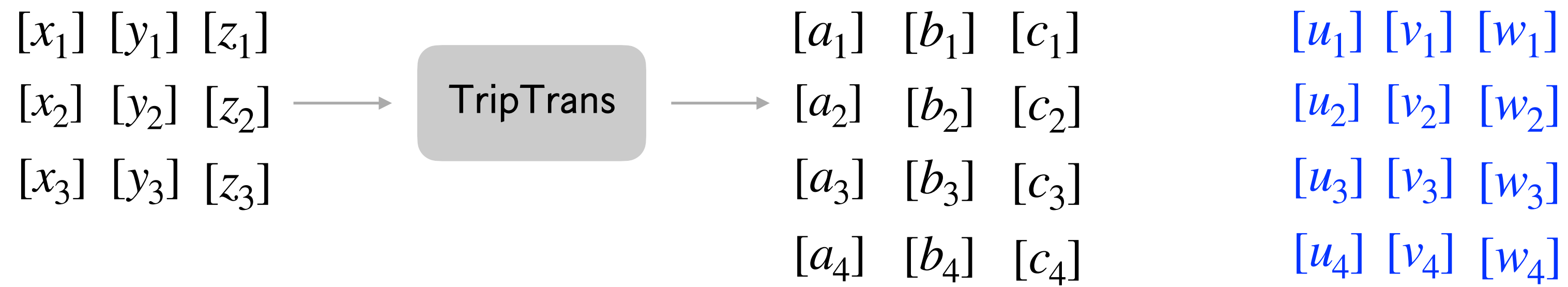
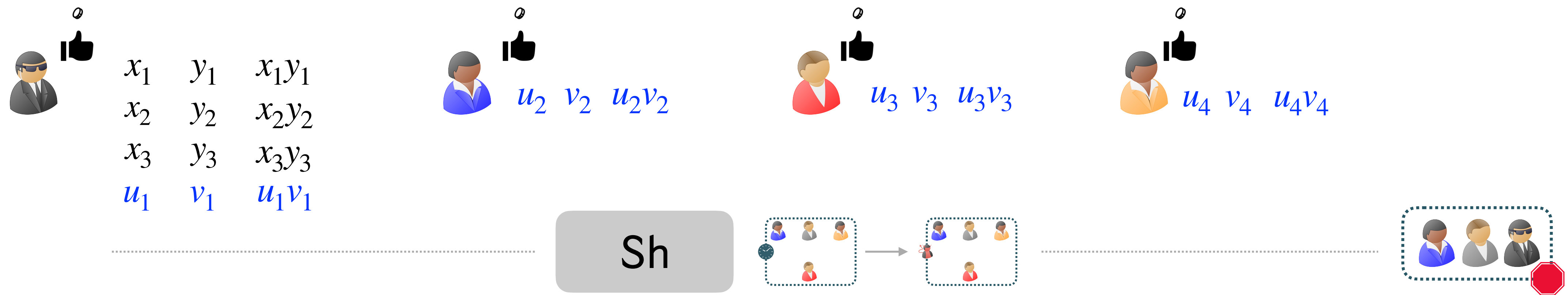
Perfect HMPC - Triple Sharing with Party Elimination Protocol



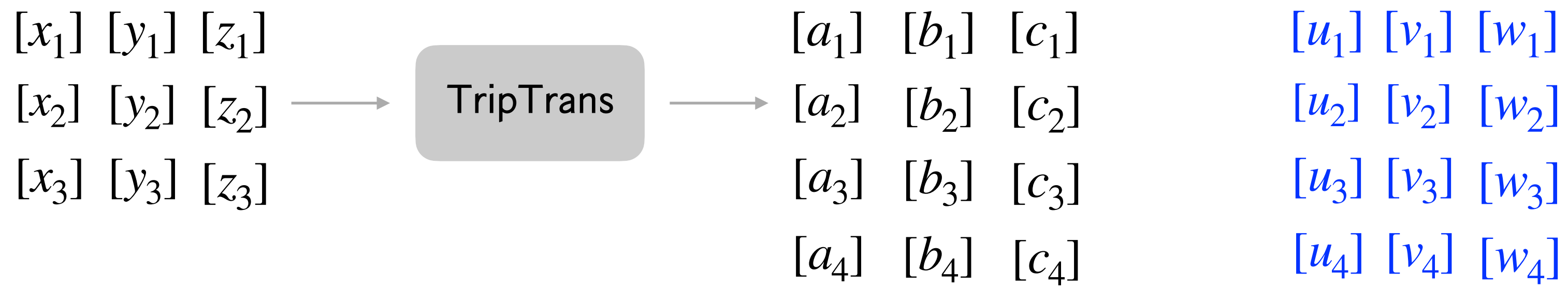
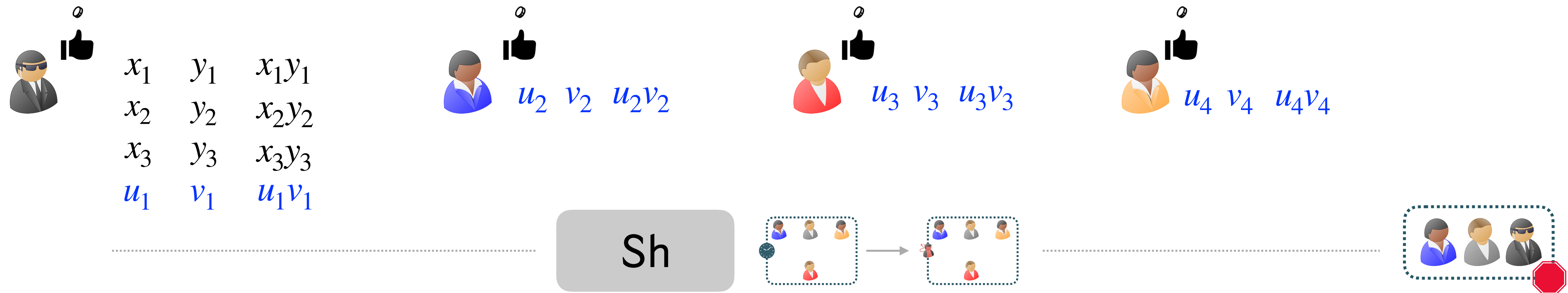
Perfect HMPC - Triple Sharing with Party Elimination Protocol



Perfect HMPC - Triple Sharing with Party Elimination Protocol

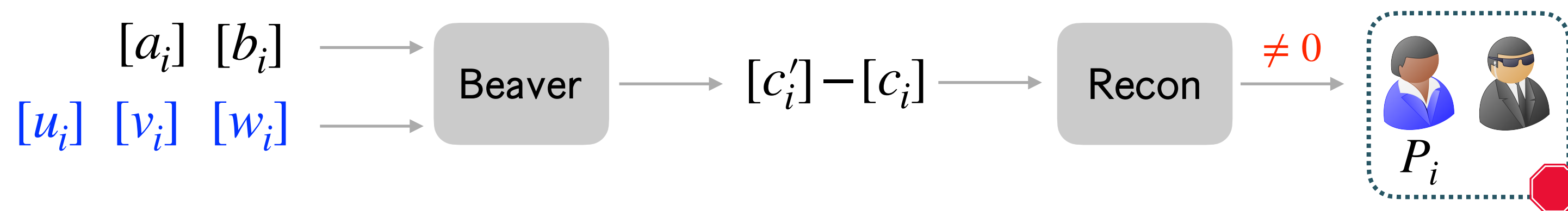


Perfect HMPC - Triple Sharing with Party Elimination Protocol

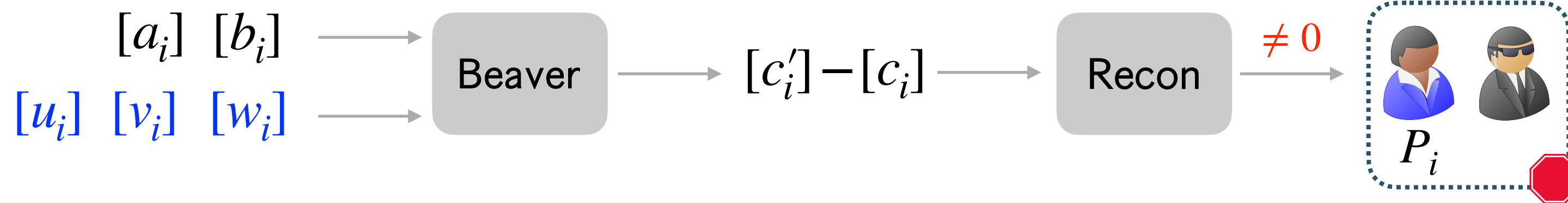
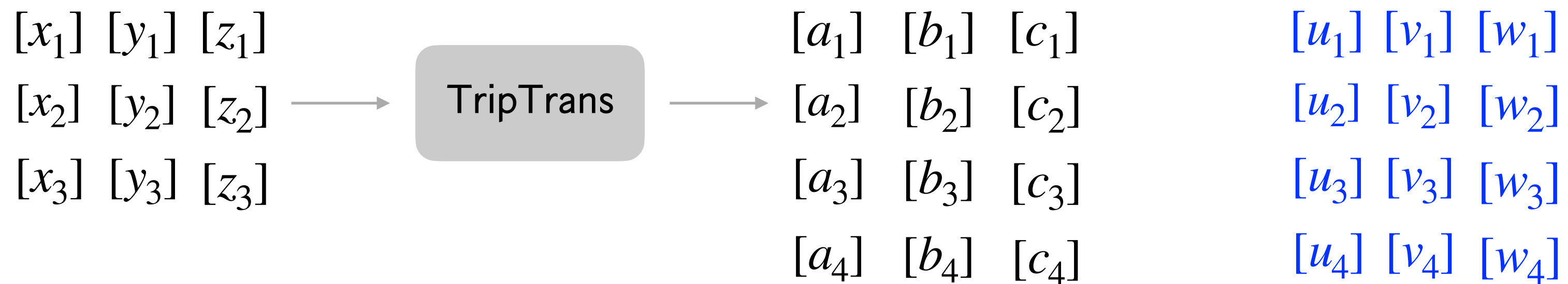
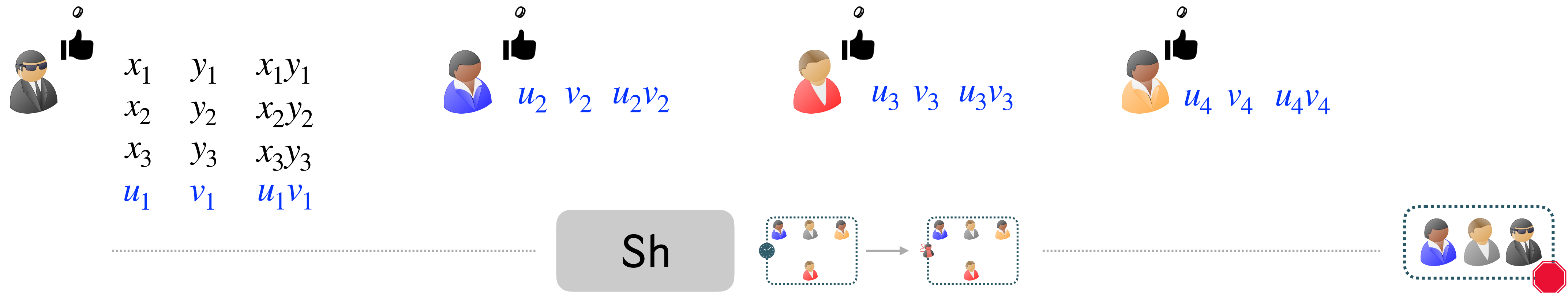


$$f_c(\cdot) = f_a(\cdot)f_b(\cdot)$$

if all checks hold



Perfect HMPC - Triple Sharing with Party Elimination Protocol

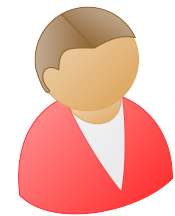
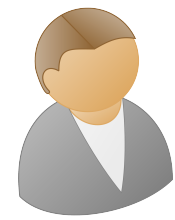


$$f_c(\cdot) = f_a(\cdot)f_b(\cdot)$$

if all checks hold

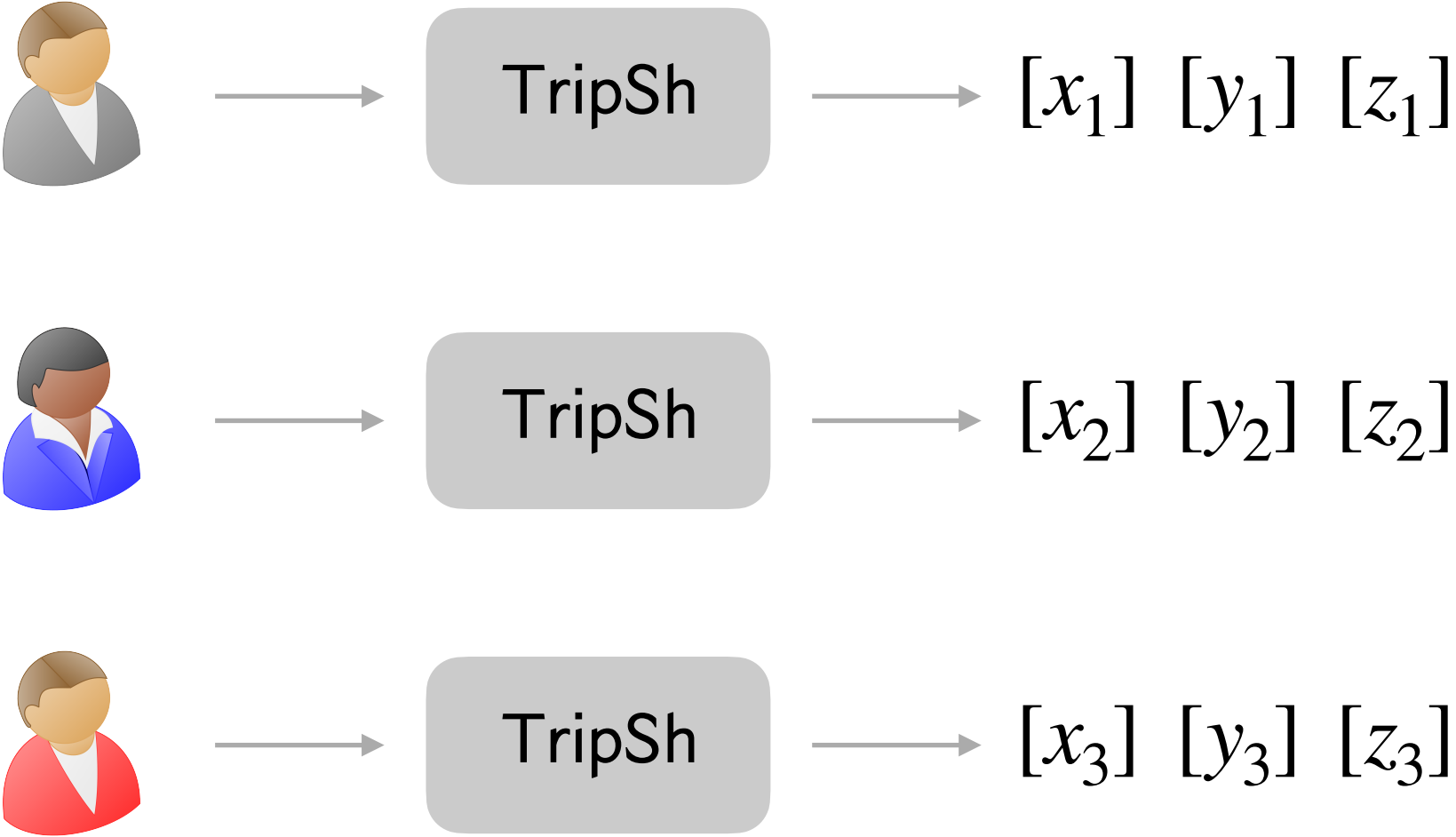
$[f_a(5)], [f_b(5)], [f_c(5)]$

Perfect HMPC - Triple Generation with Party Elimination



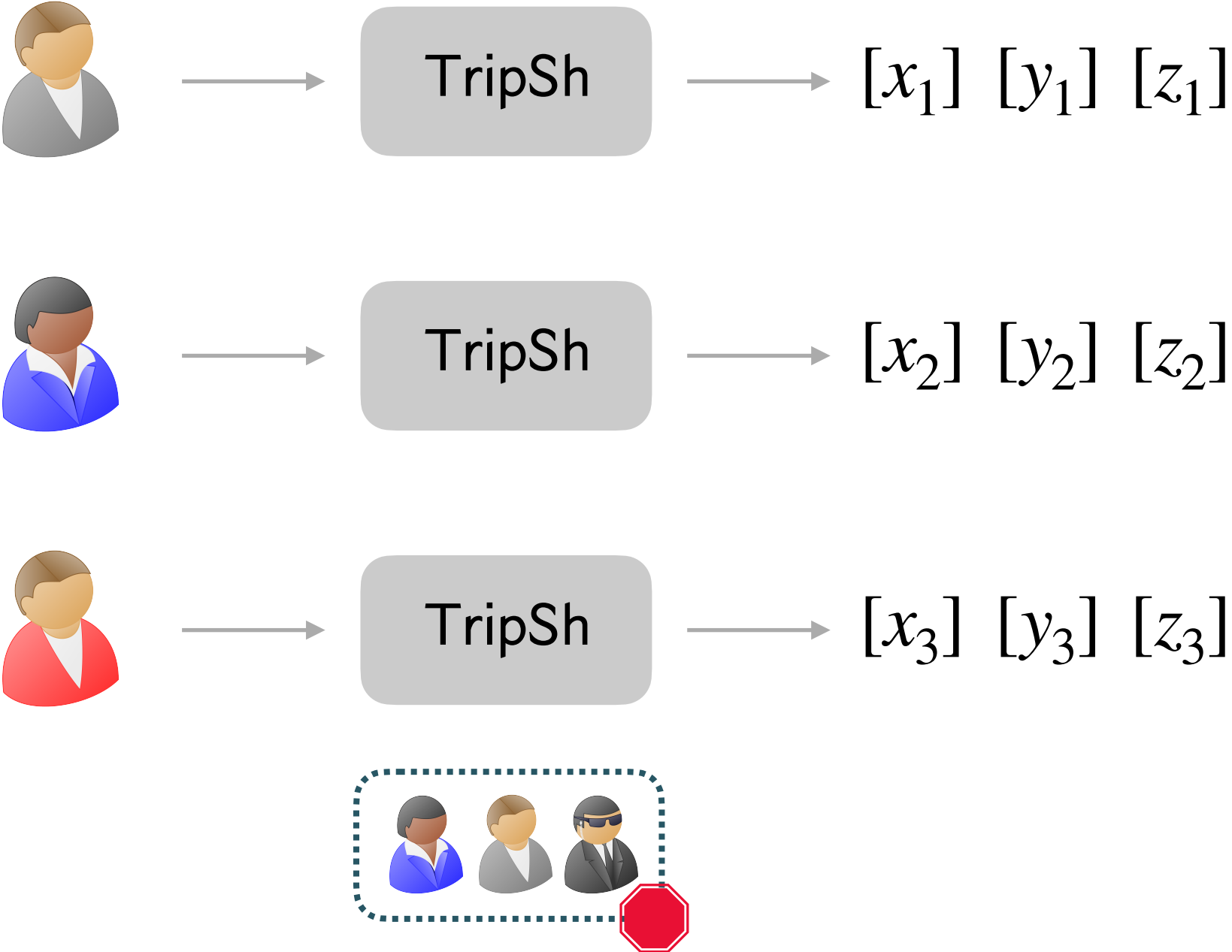
TripGen

Perfect HMPC - Triple Generation with Party Elimination



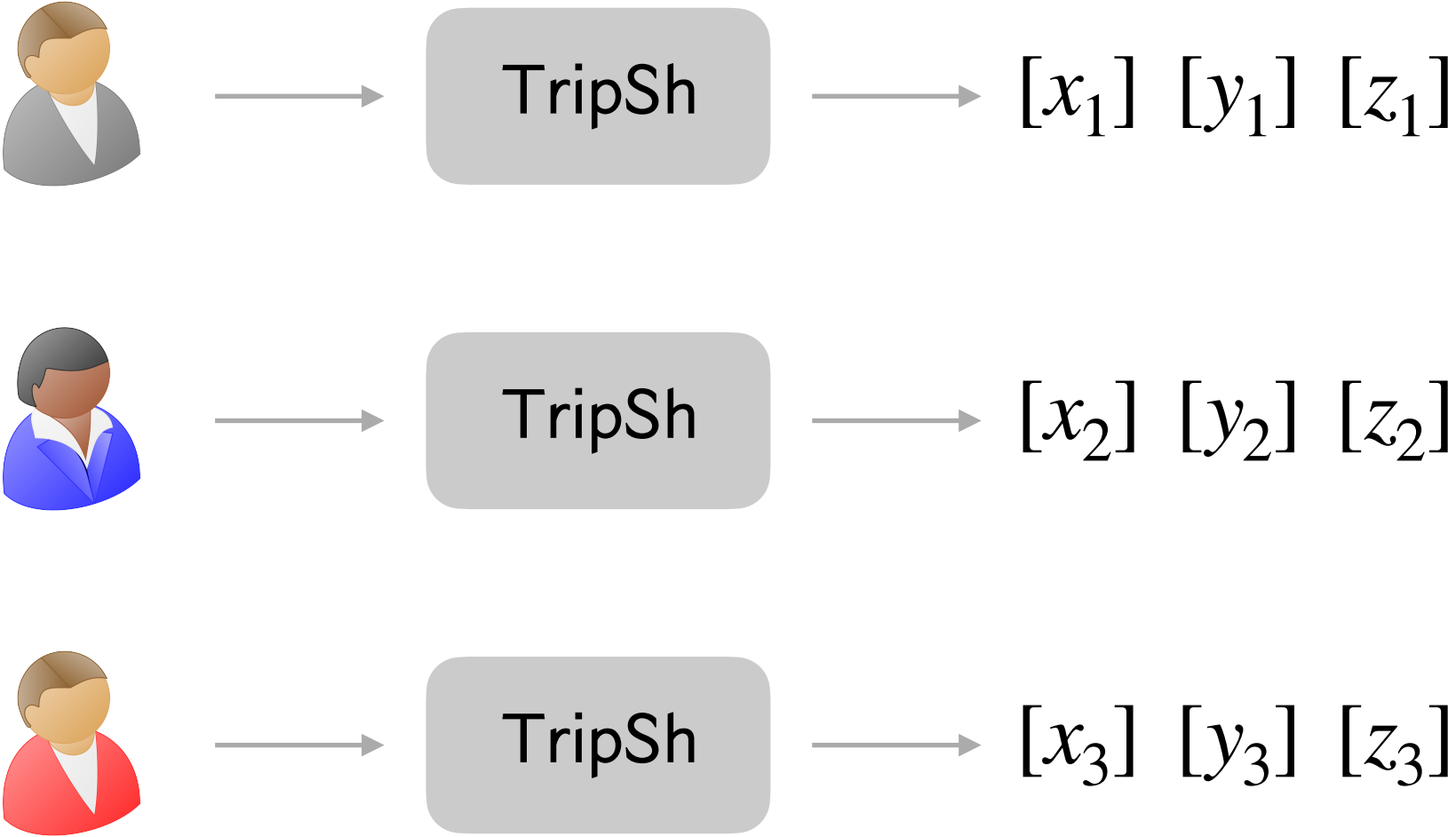
TripGen

Perfect HMPC - Triple Generation with Party Elimination



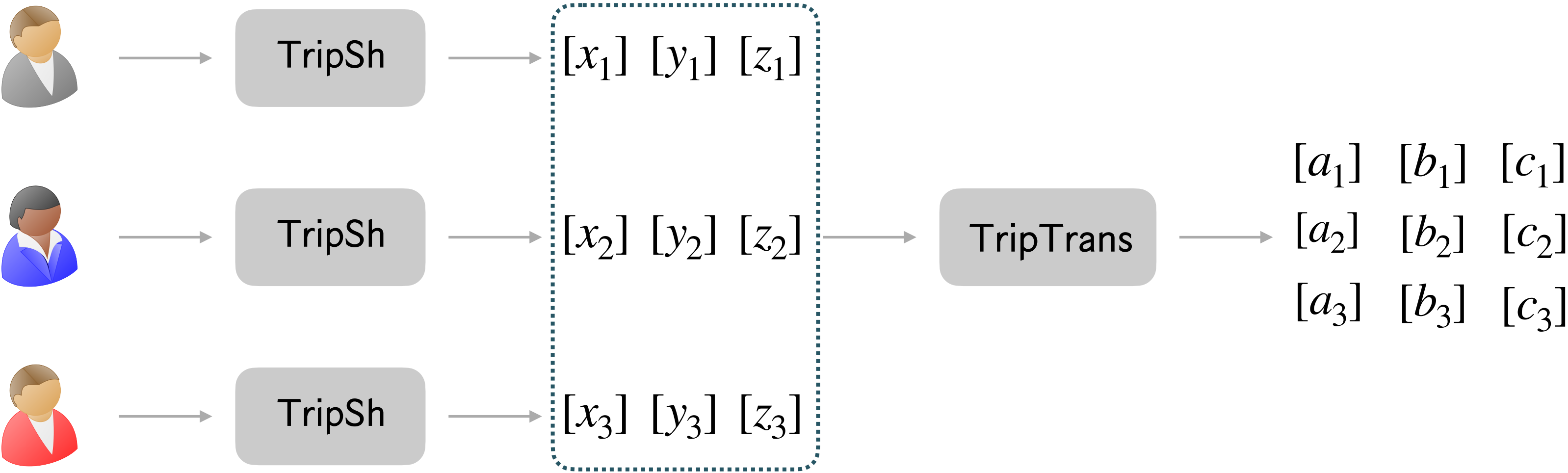
TripGen

Perfect HMPC - Triple Generation with Party Elimination



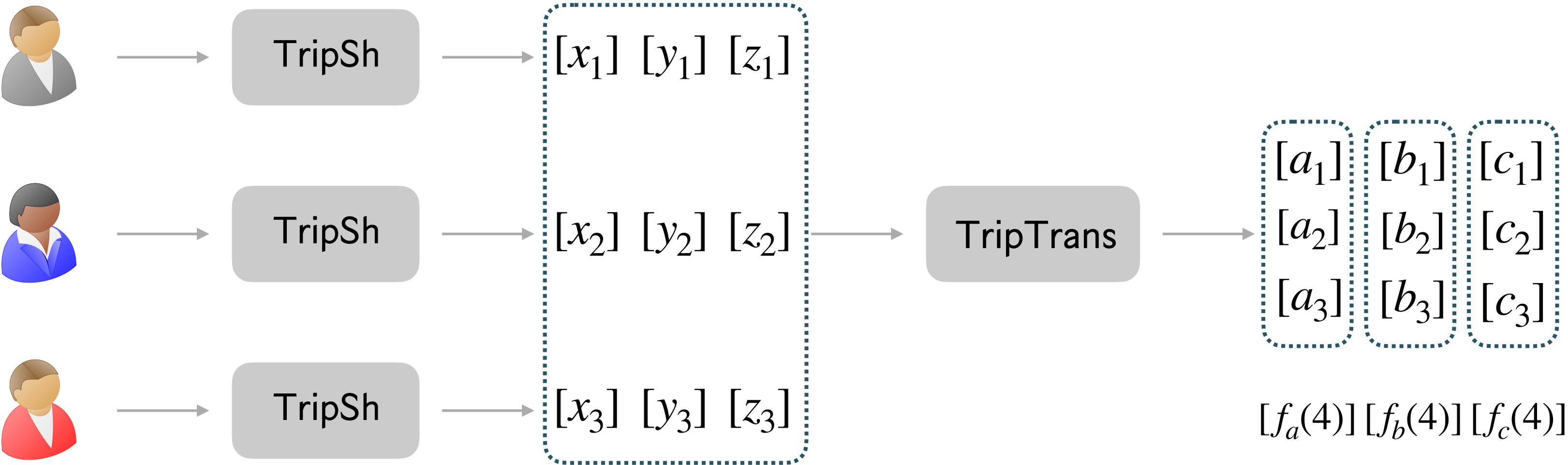
TripGen

Perfect HMPC - Triple Generation with Party Elimination



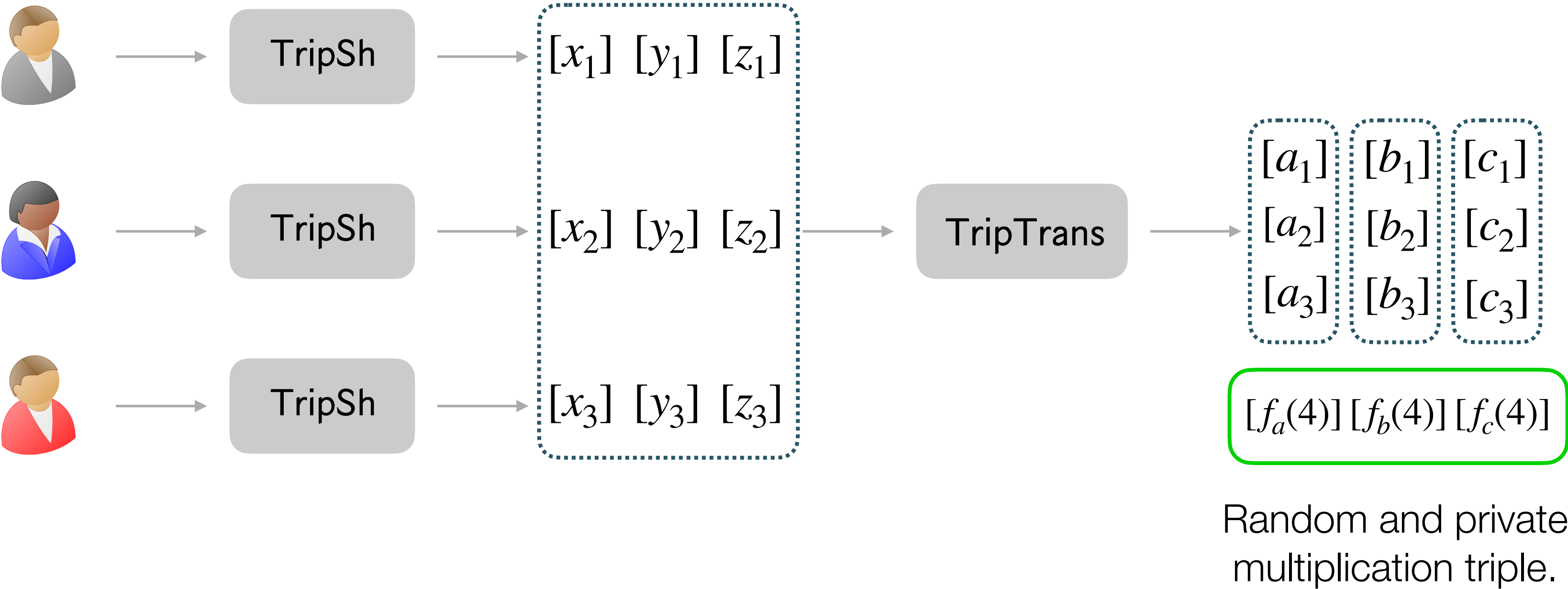
TripGen

Perfect HMPC - Triple Generation with Party Elimination



TripGen

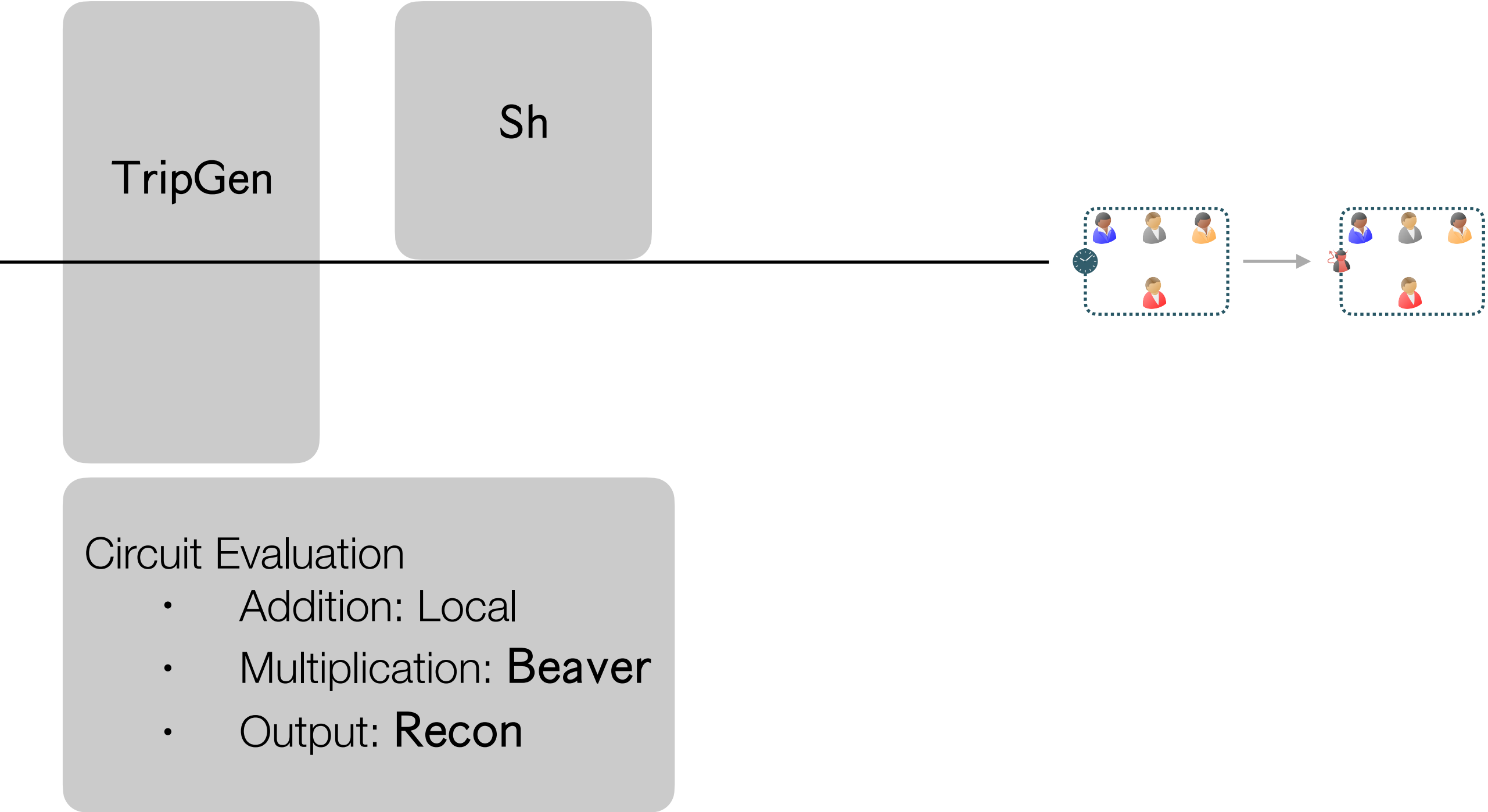
Perfect HMPC - Triple Generation with Party Elimination



TripGen

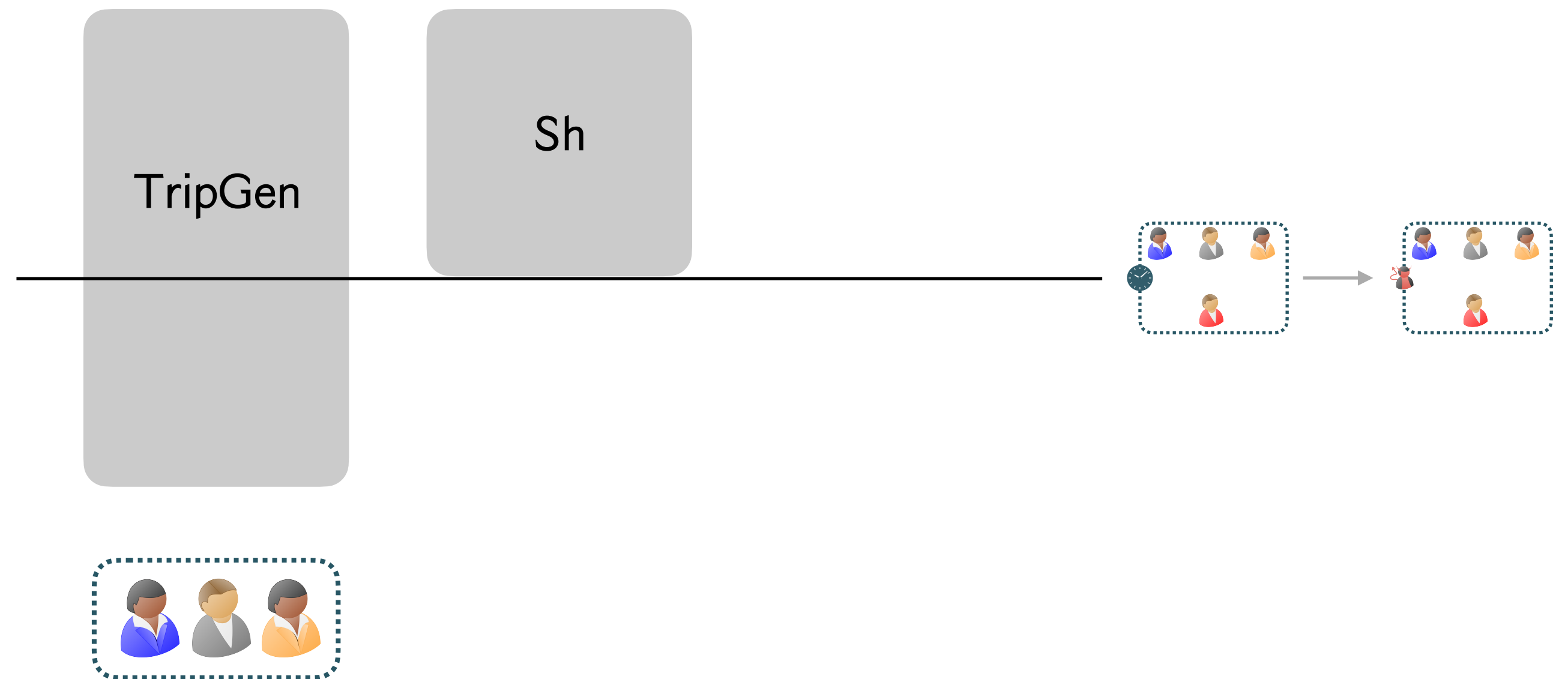
Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase



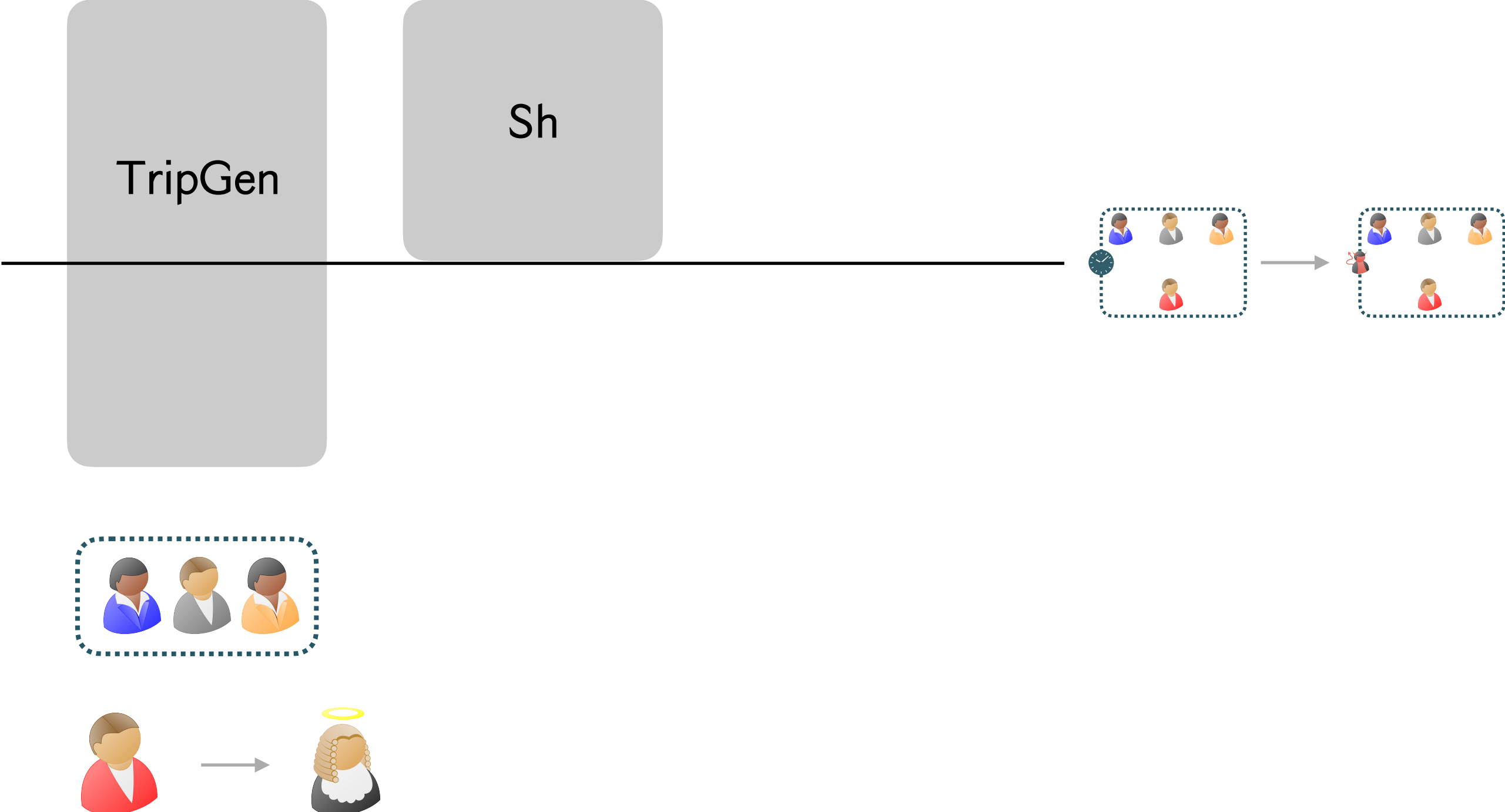
Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase



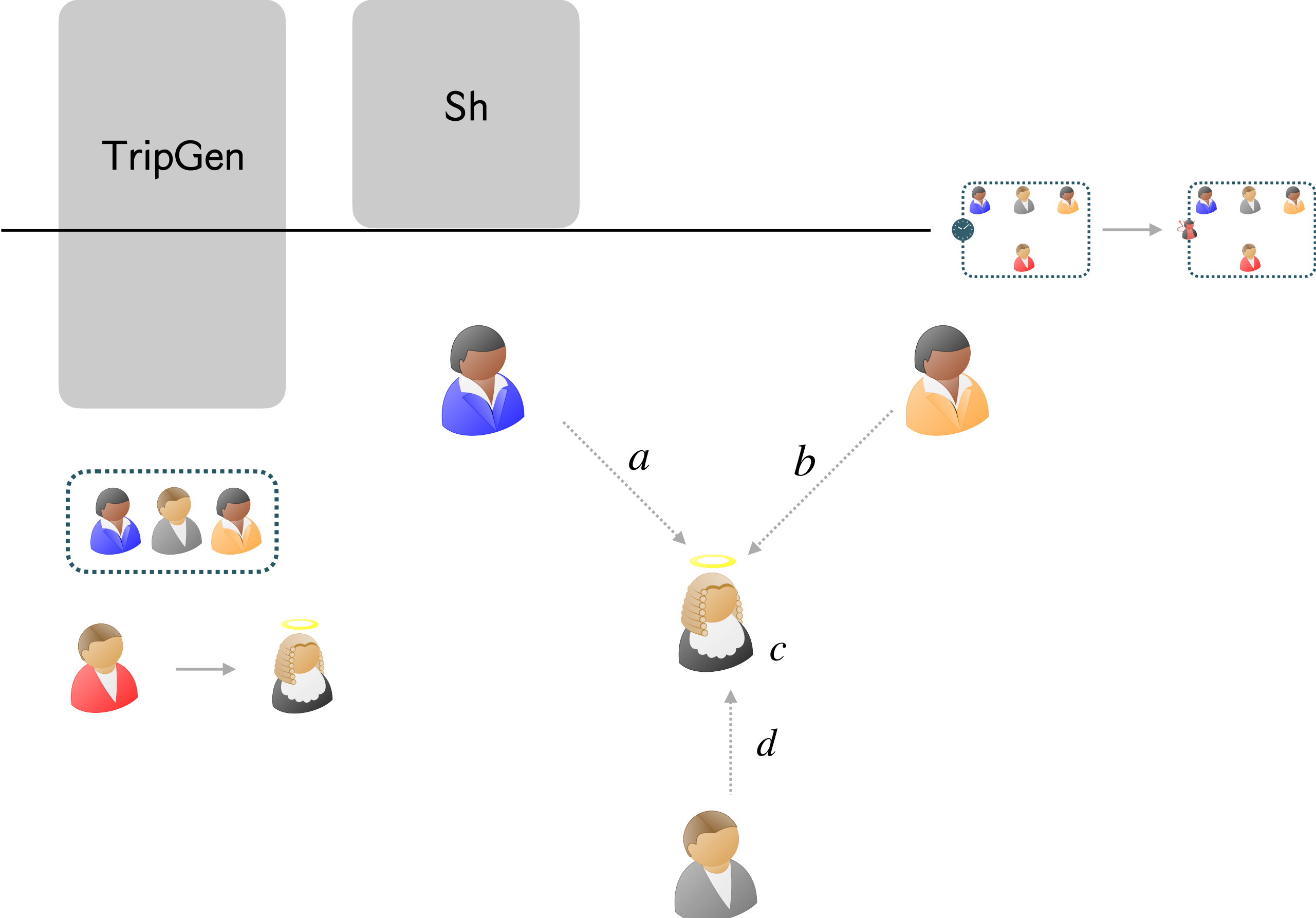
Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase



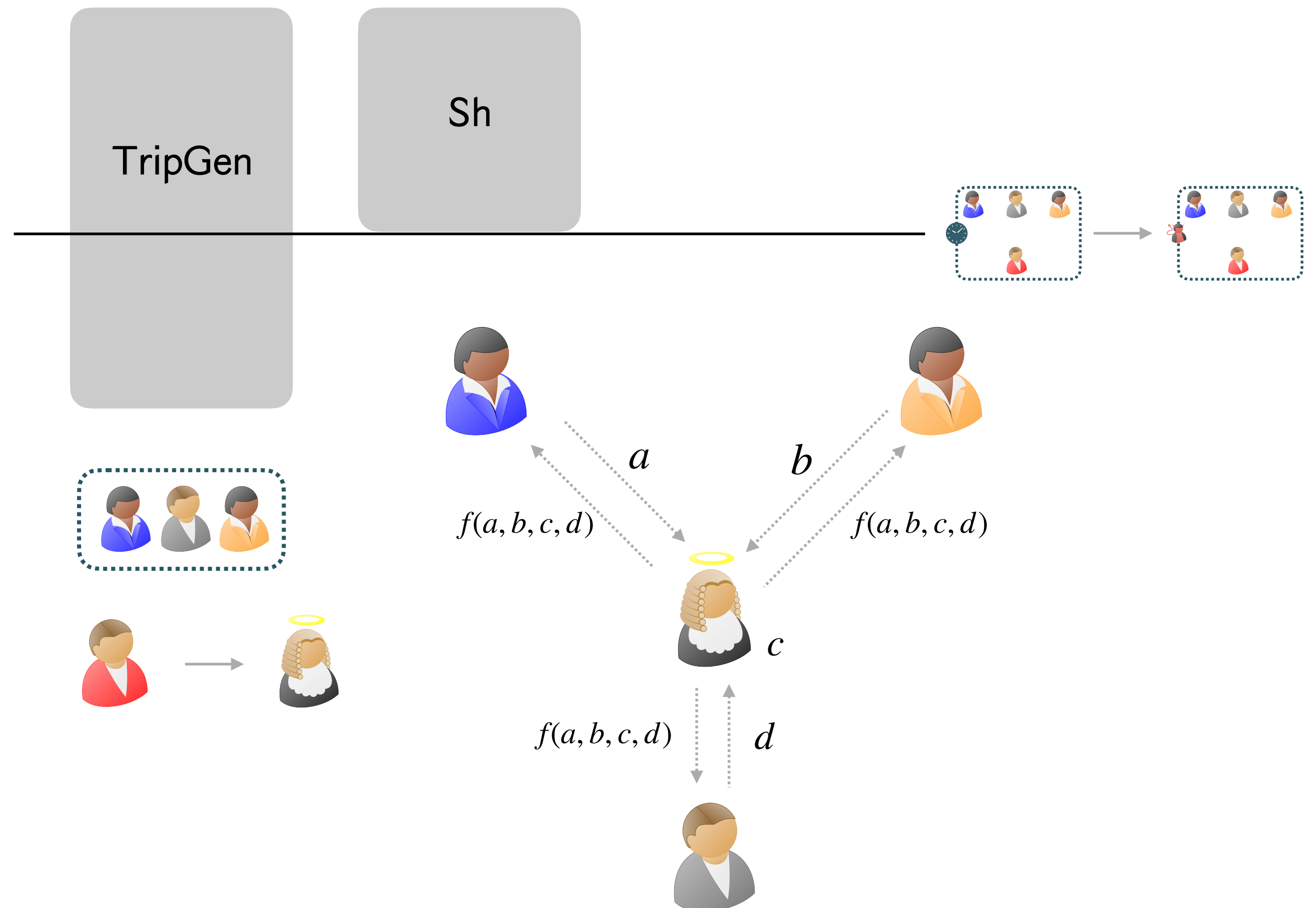
Perfect HMPC

- 3 phases
 - Triple generation phase
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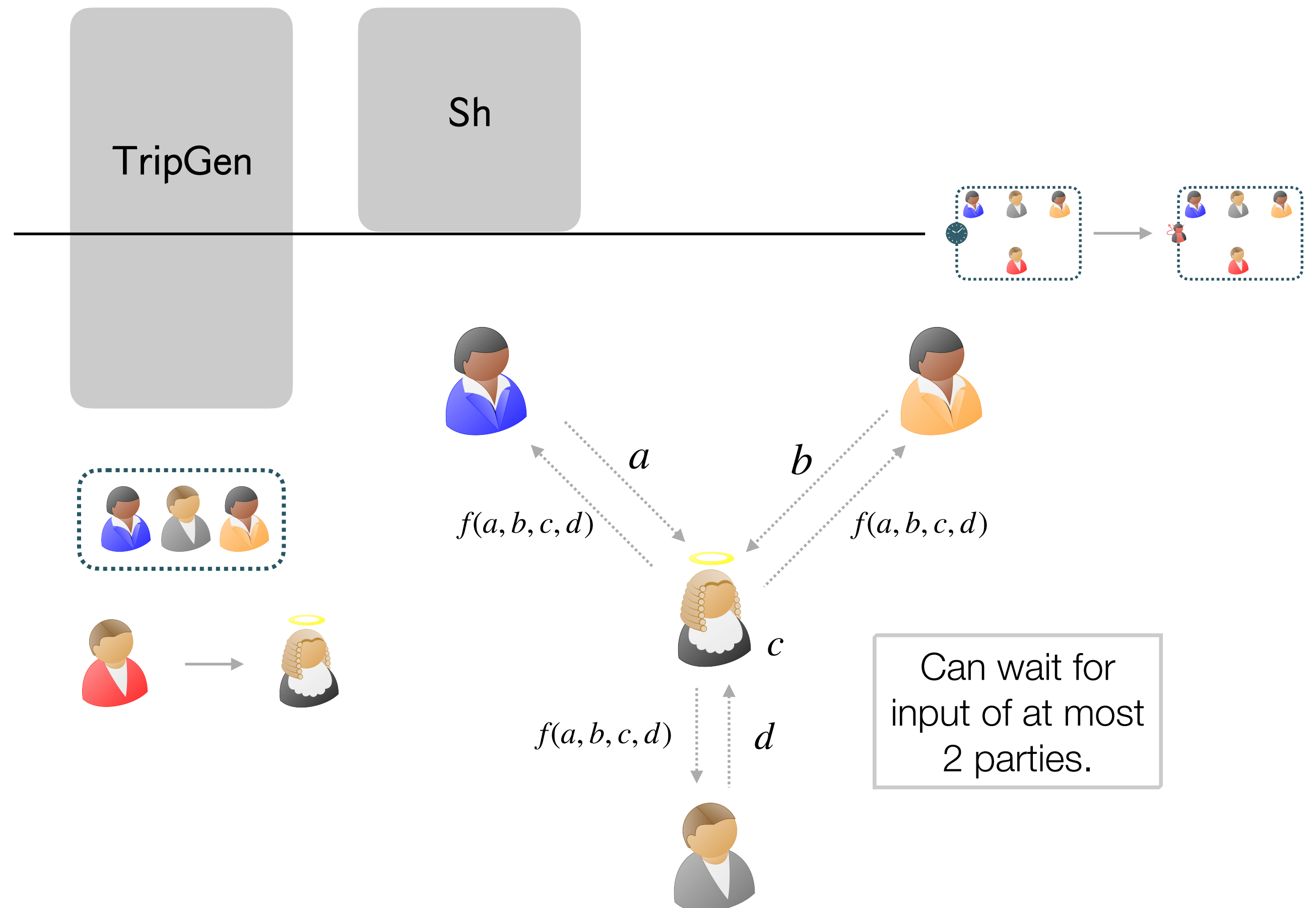
Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase

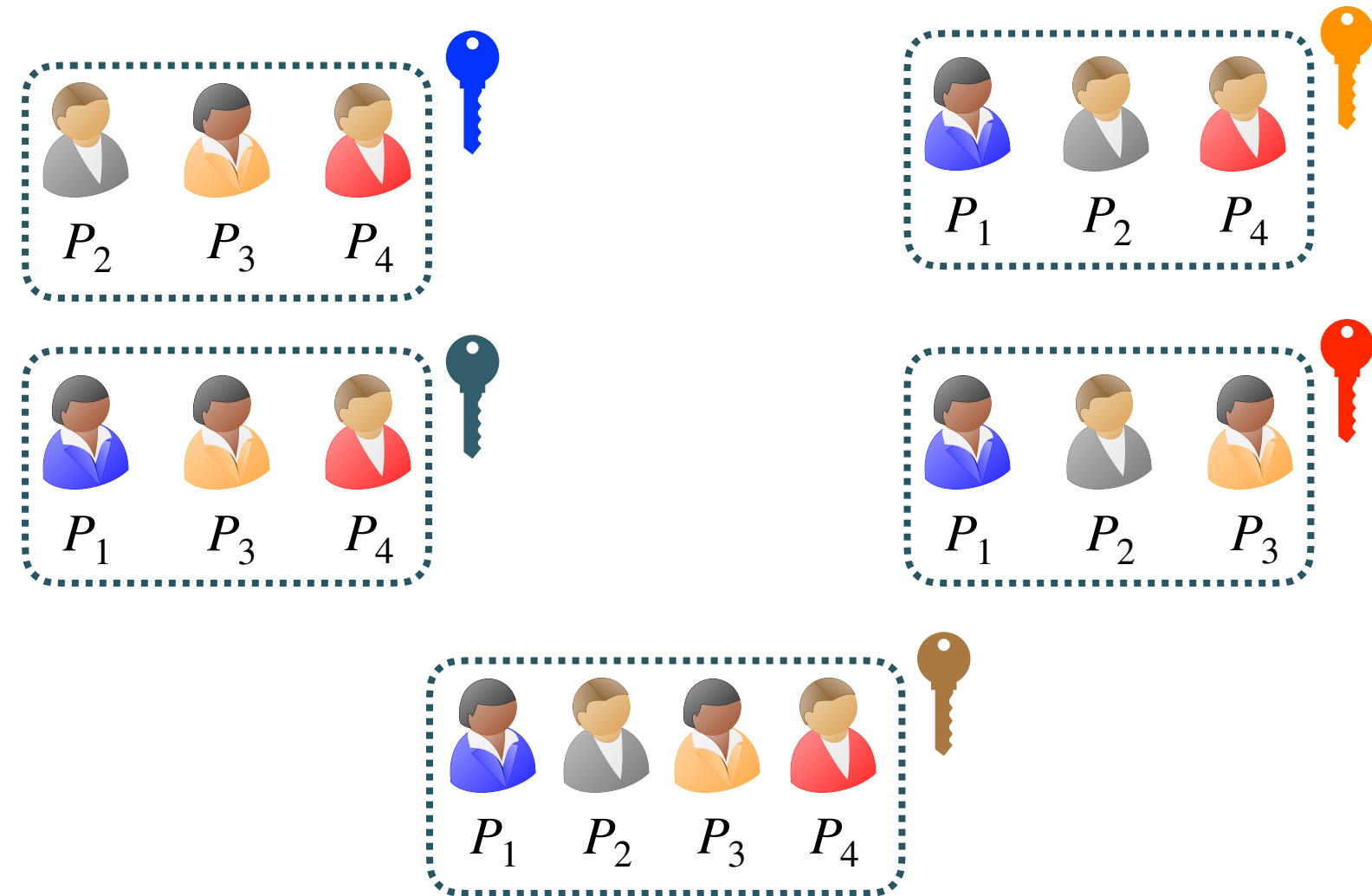


Perfect HMPC

- 3 phases
 - Triple generation phase
 - Input phase
 - Circuit evaluation and output phase

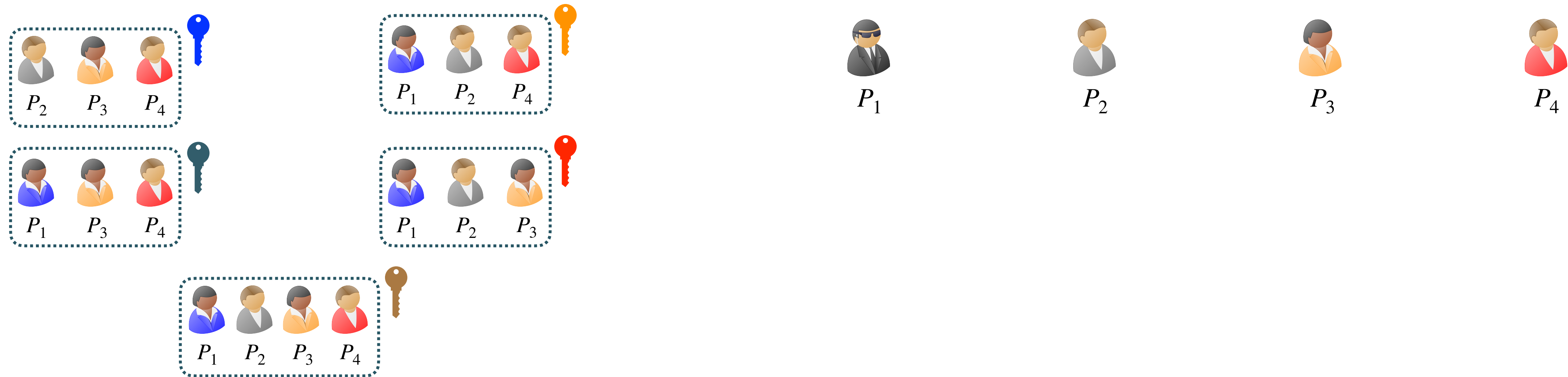


Cryptographically Secure HMPC - Secret Sharing and Reconstruction



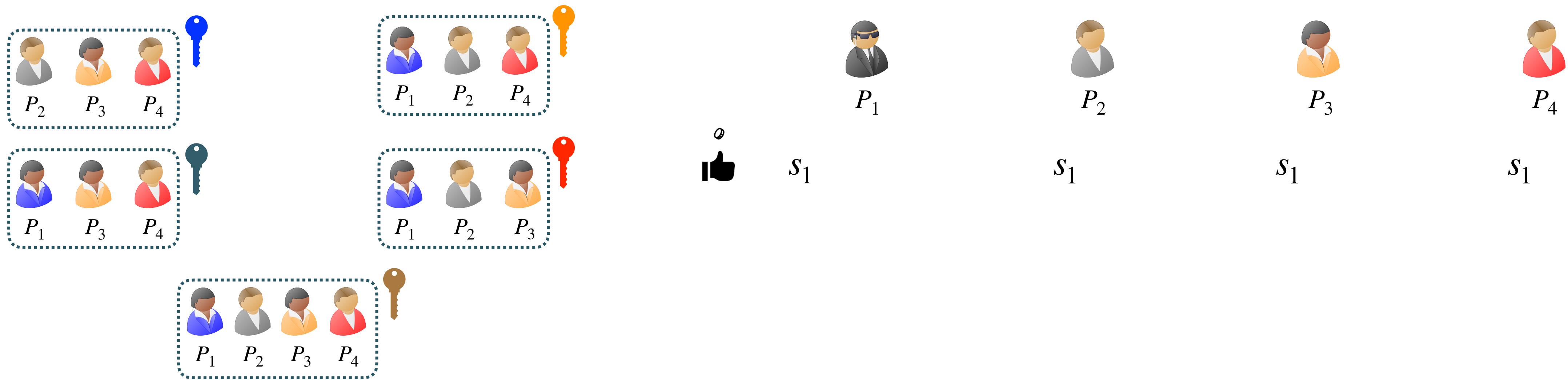
- Assume symmetric-key setup for PRF [AFL+16,CCP+19,MR18]

Cryptographically Secure HMPC - Secret Sharing and Reconstruction



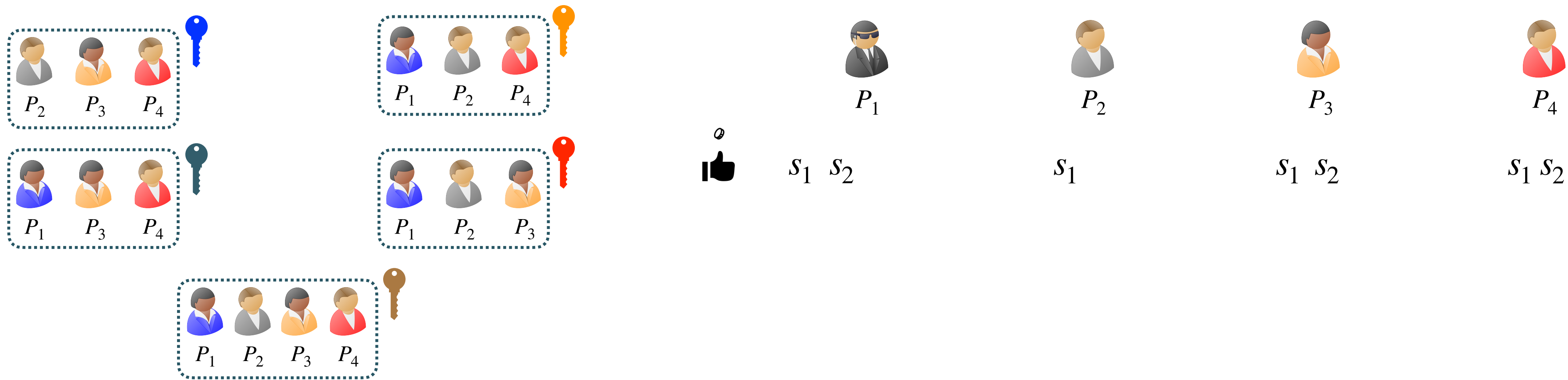
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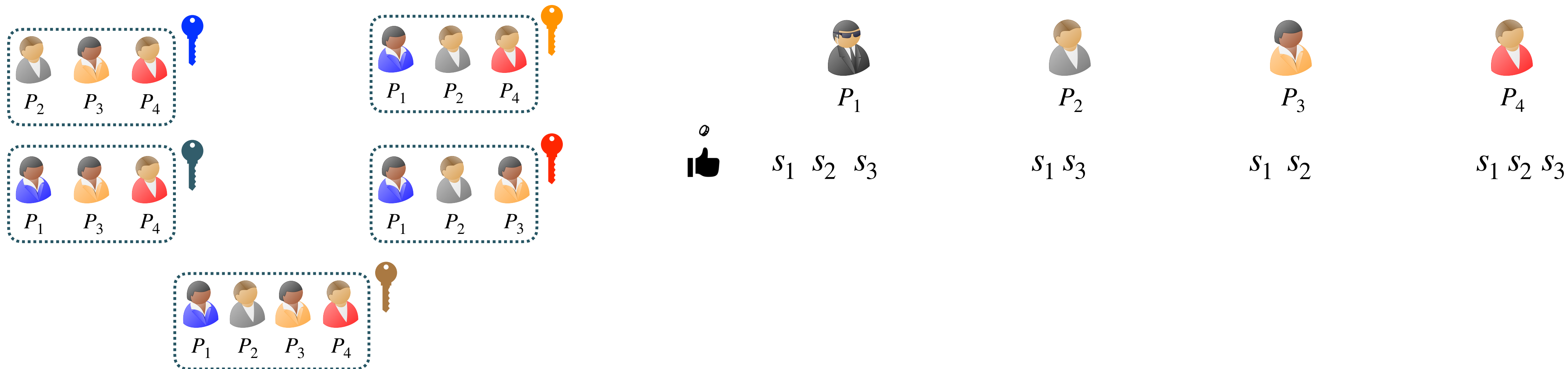
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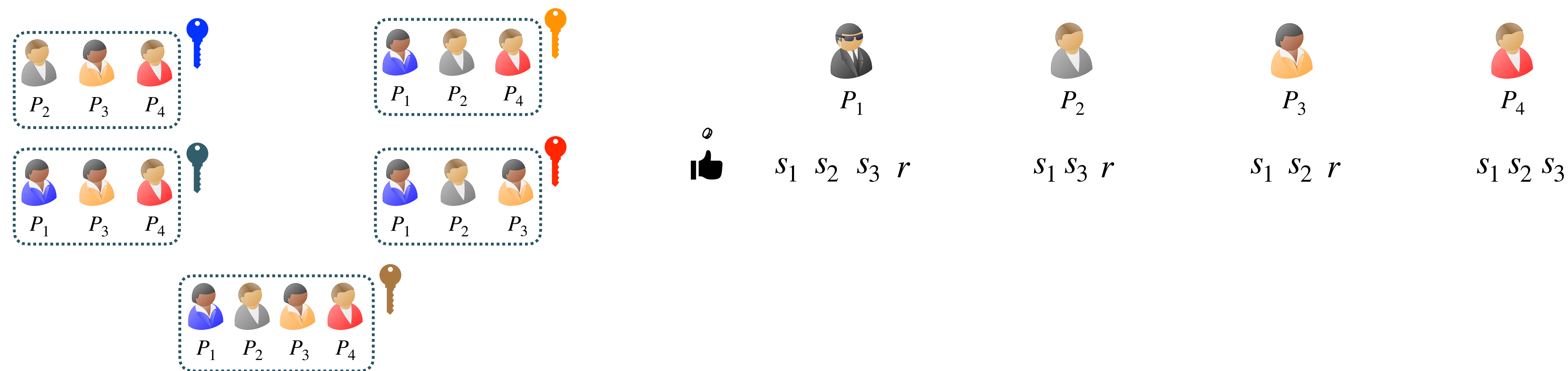
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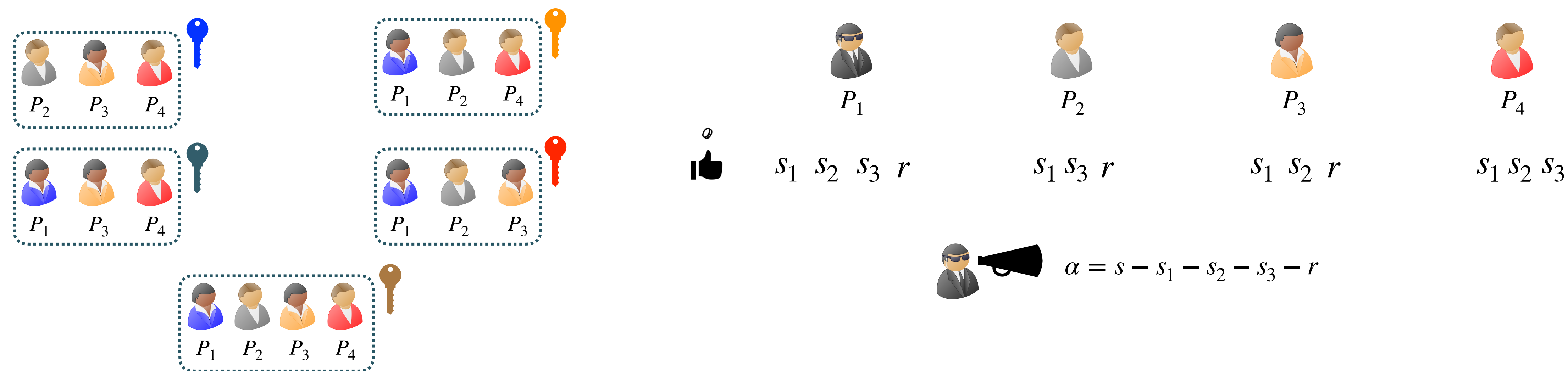
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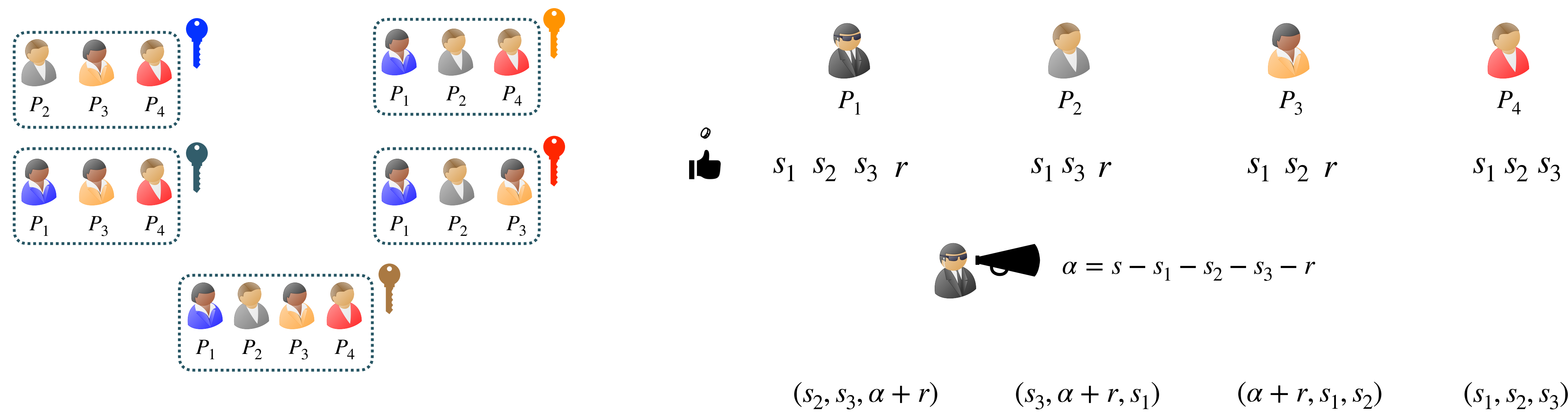
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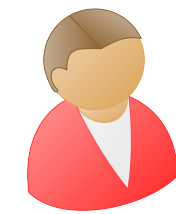
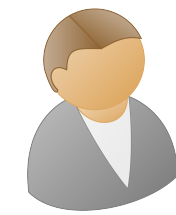
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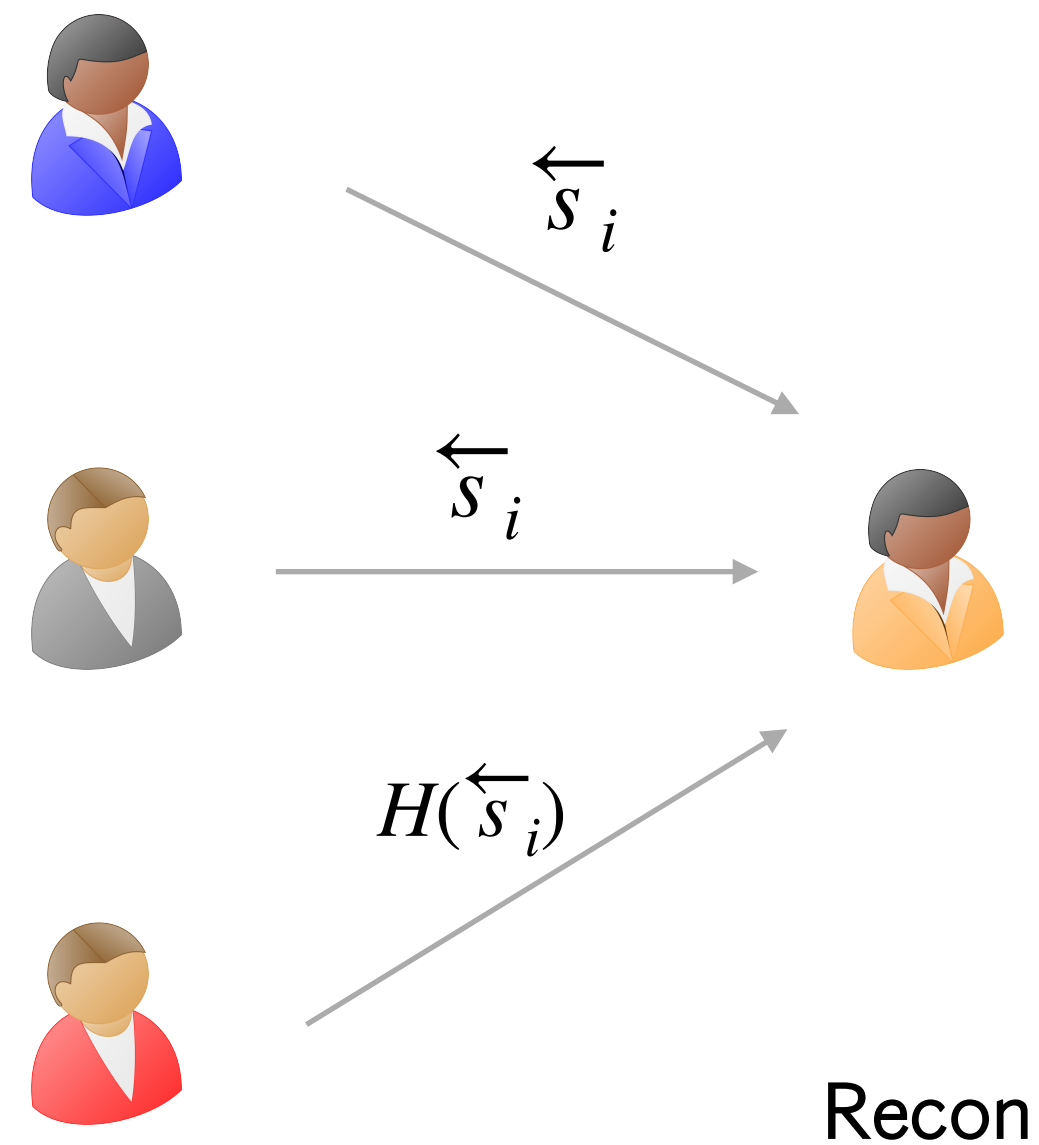
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Recon

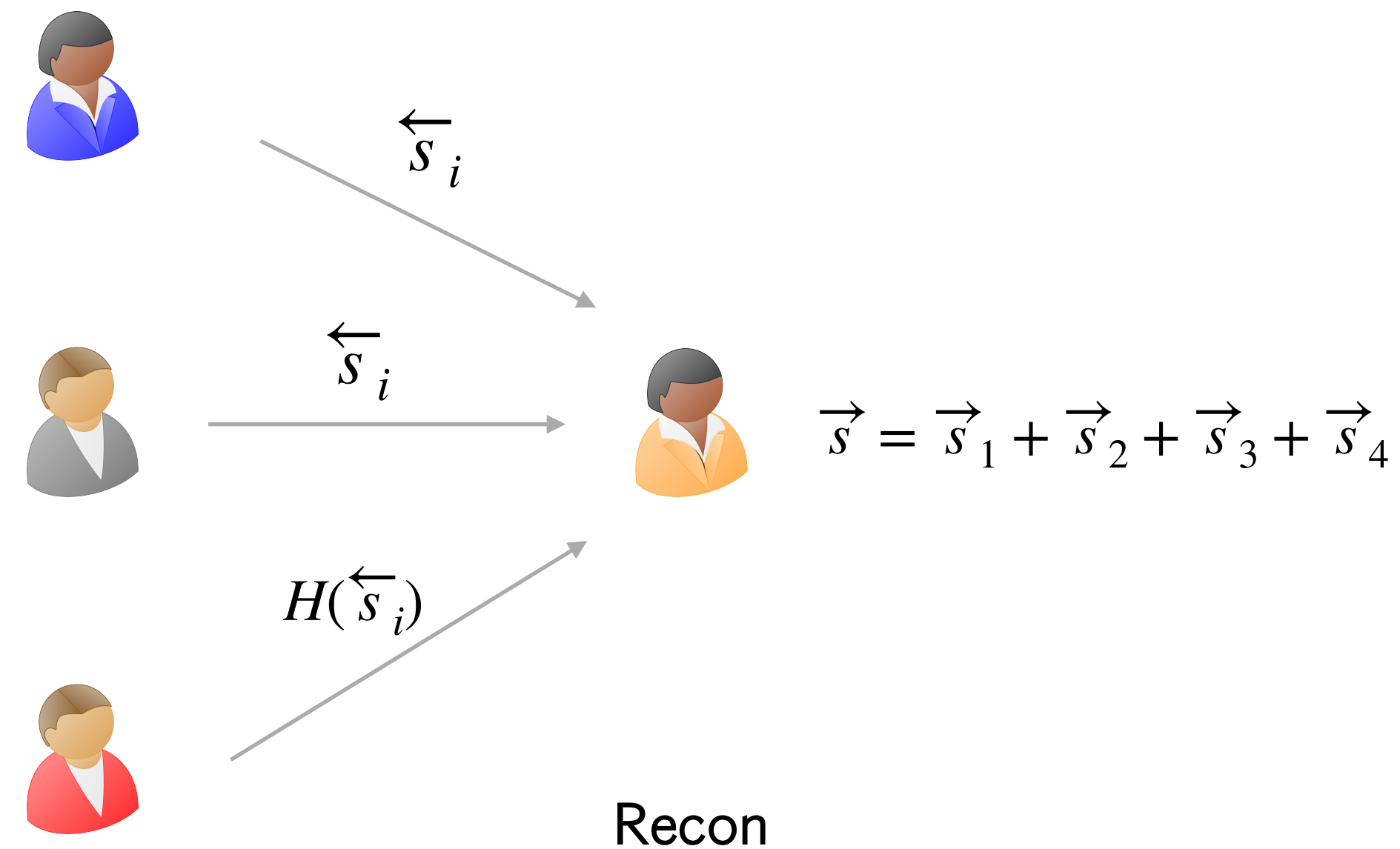
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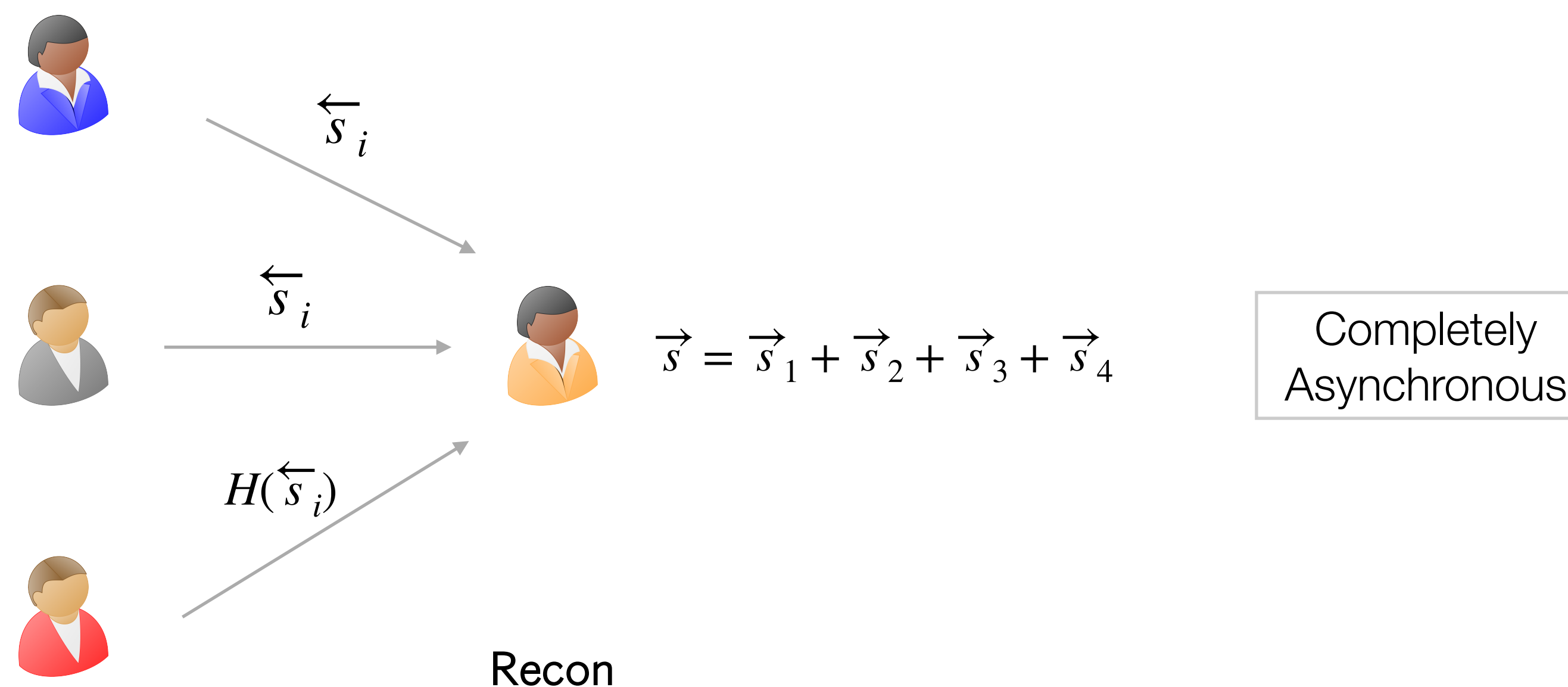
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Cryptographically Secure HMPC - Triple Generation Protocol

- Triple sharing similar to **TripSh**
 - Dealer shares $2l + 1$ triples instead of 3 triples
 - Other parties don't share triples

Cryptographically Secure HMPC - Triple Generation Protocol

- Triple sharing similar to **TripSh**

$[x_1], [y_1], [z_1]$

$[x_2], [y_2], [z_2]$

\vdots

$[x_{2l+1}], [y_{2l+1}], [z_{2l+1}]$

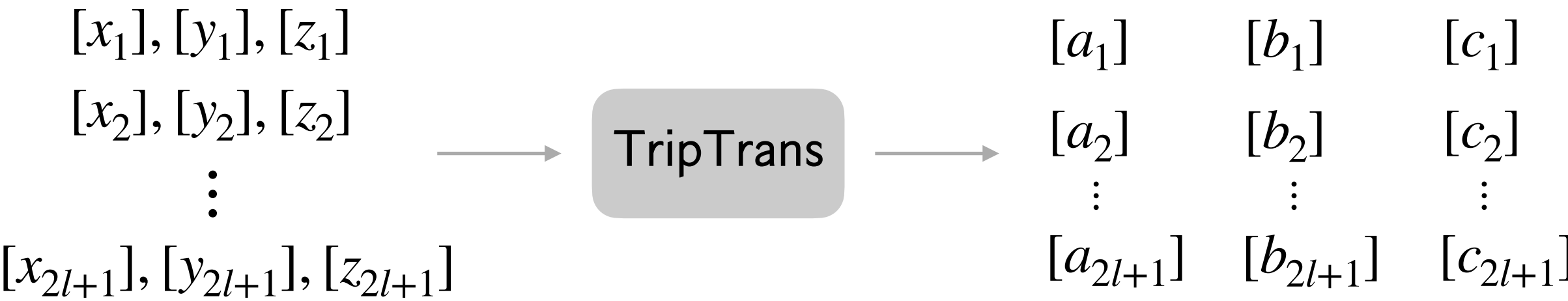
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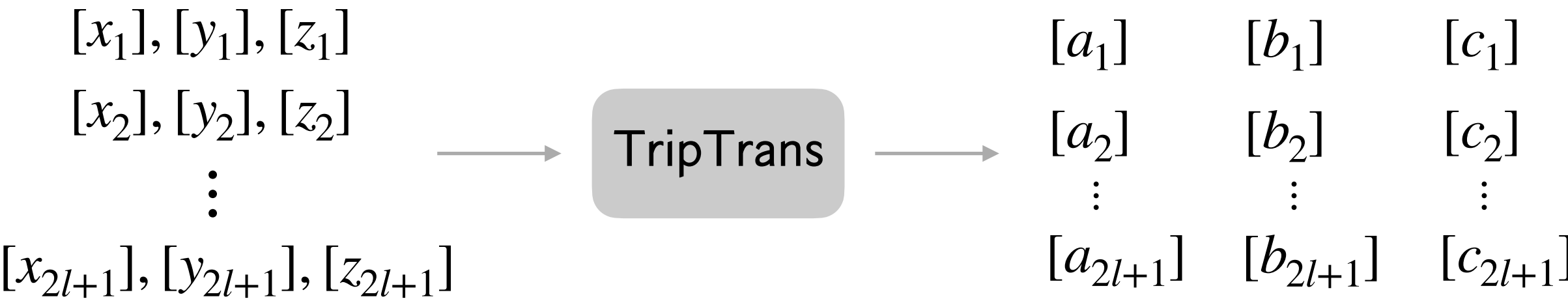


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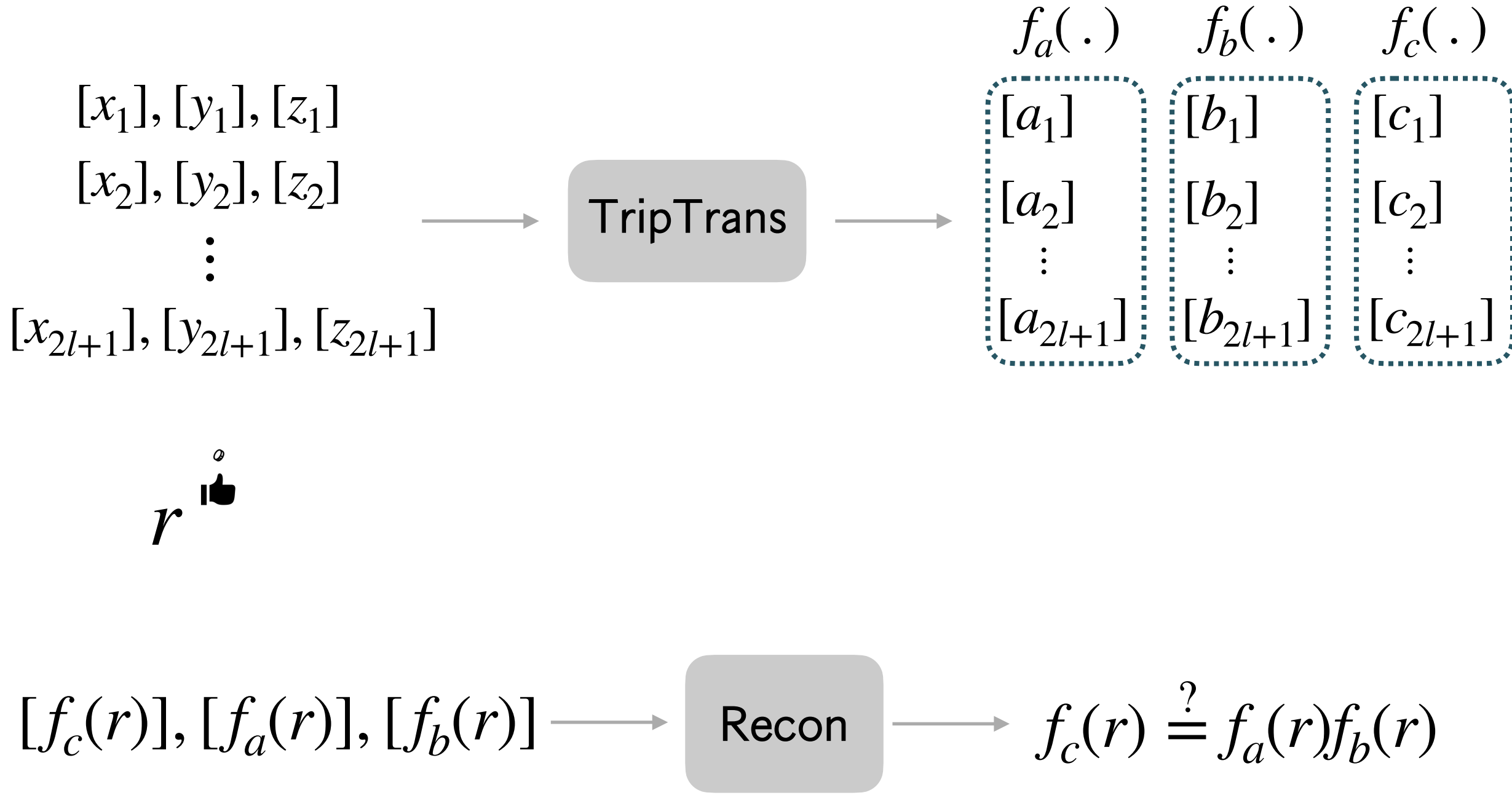
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r 

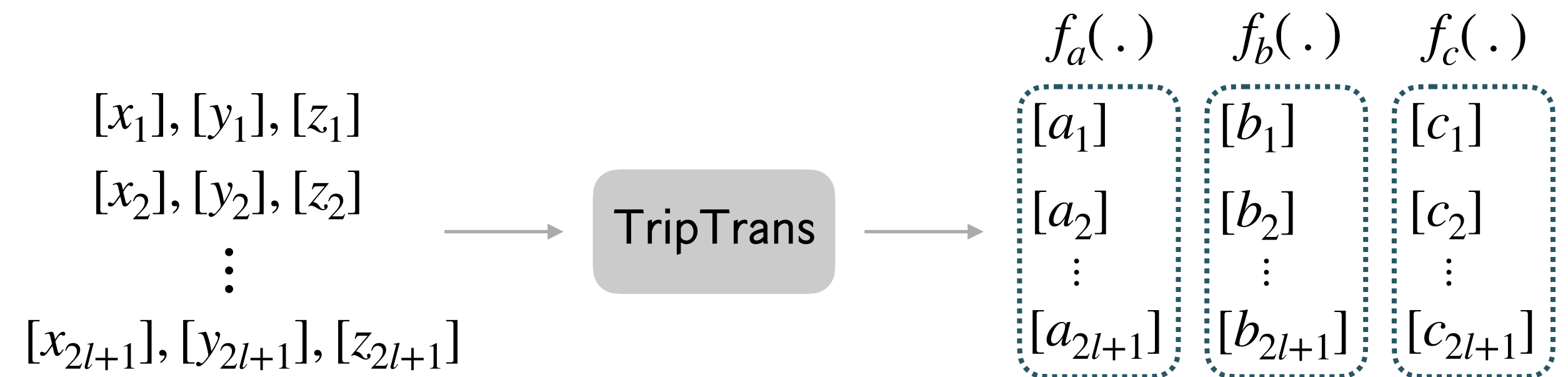
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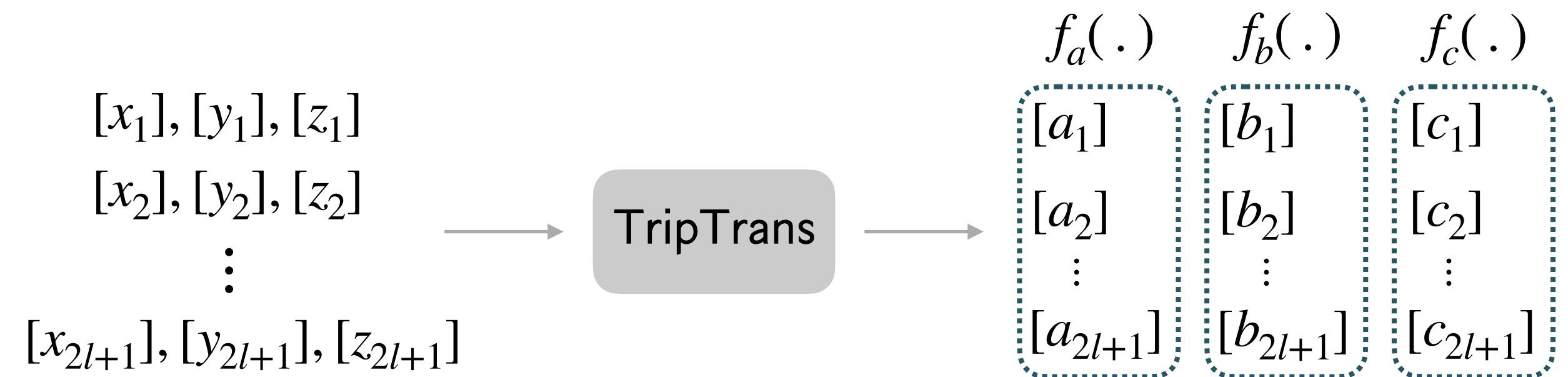
r



- ✓ l multiplication triples
- ✗ l shares of $([0], [0], [0])$

Cryptographically Secure HMPC - Triple Generation Protocol

- Triple sharing similar to **TripSh**
 - Dealer shares $2l + 1$ triples instead of 3 triples
 - Other parties don't share triples
- Triple generation similar to **TripGen**
 - Each instance outputs l triples



r



- ✓ l multiplication triples
- ✗ l shares of $([0], [0], [0])$

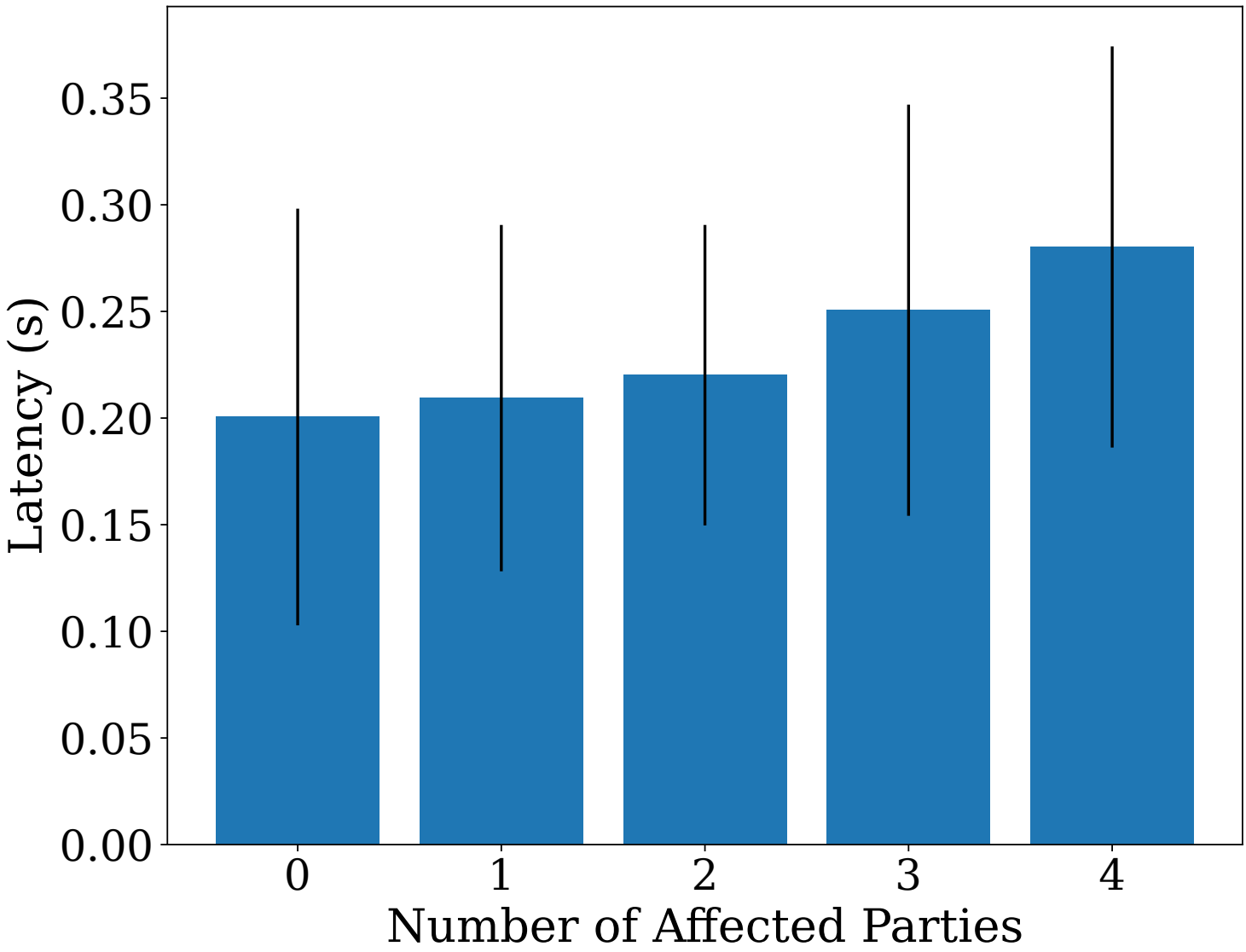
Cryptographically Secure HMPC and AMPC

- Cryptographically secure HMPC
 - Triple generation phase and input phase use 1 synchronous round
 - Circuit evaluation is completely asynchronous
 - Input provision

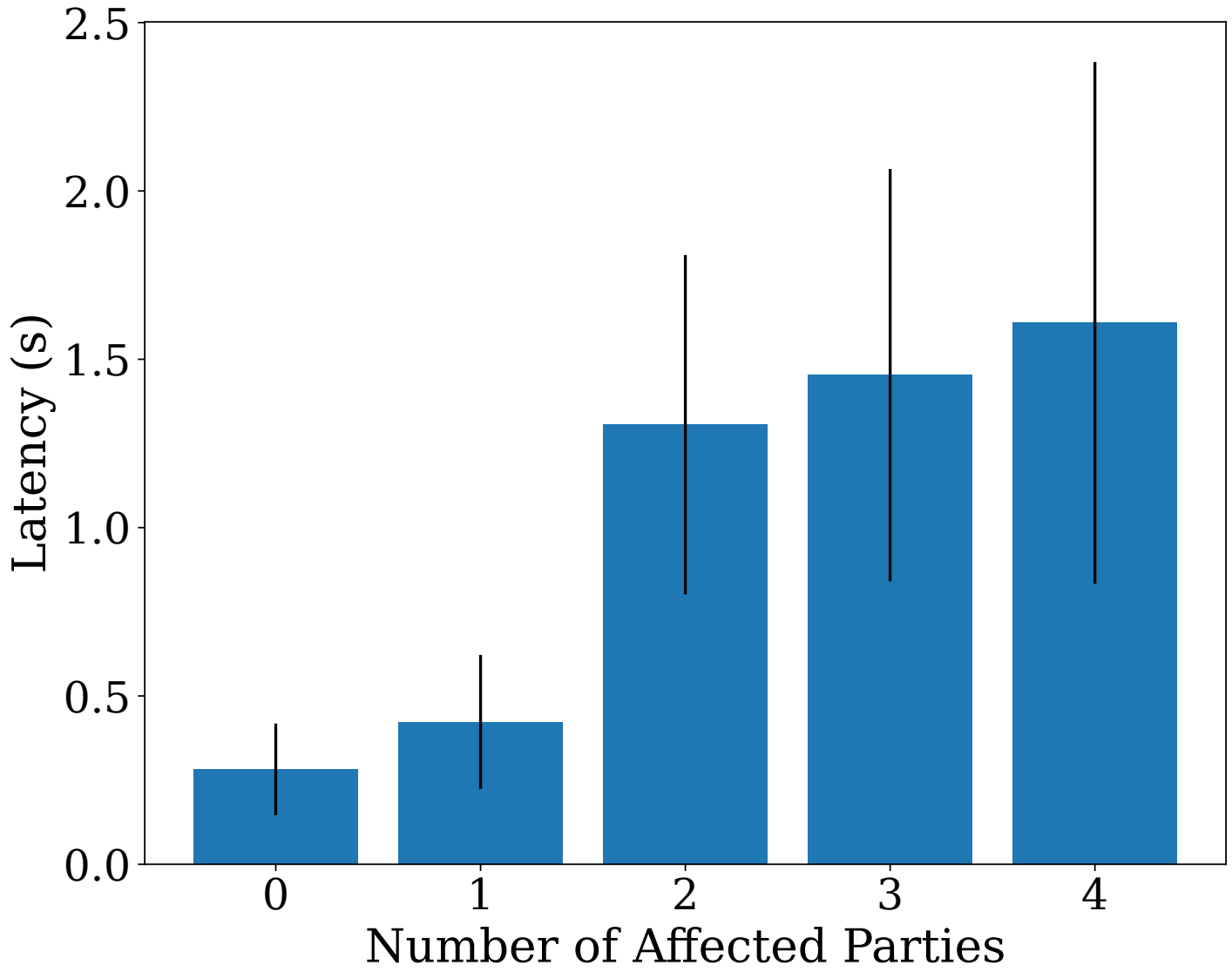
Cryptographically Secure HMPC and AMPC

- Cryptographically secure HMPC
 - Triple generation phase and input phase use 1 synchronous round
 - Circuit evaluation is completely asynchronous
 - Input provision
- Cryptographically secure AMPC
 - Similar to Cryptographically secure HMPC
 - No synchronous broadcast \implies **ACast** and **ACS**
 - No input provision

Conclusion

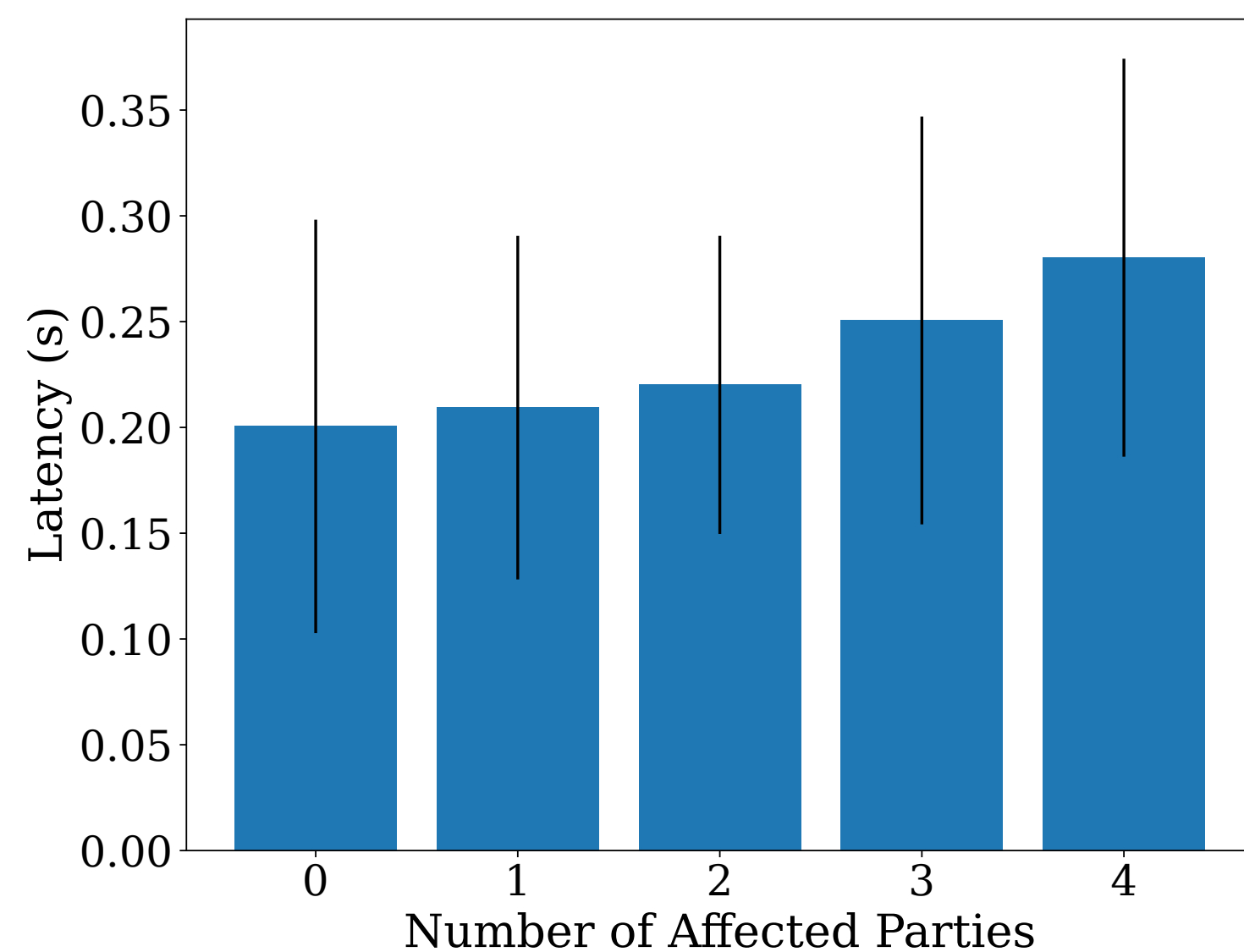


LAN

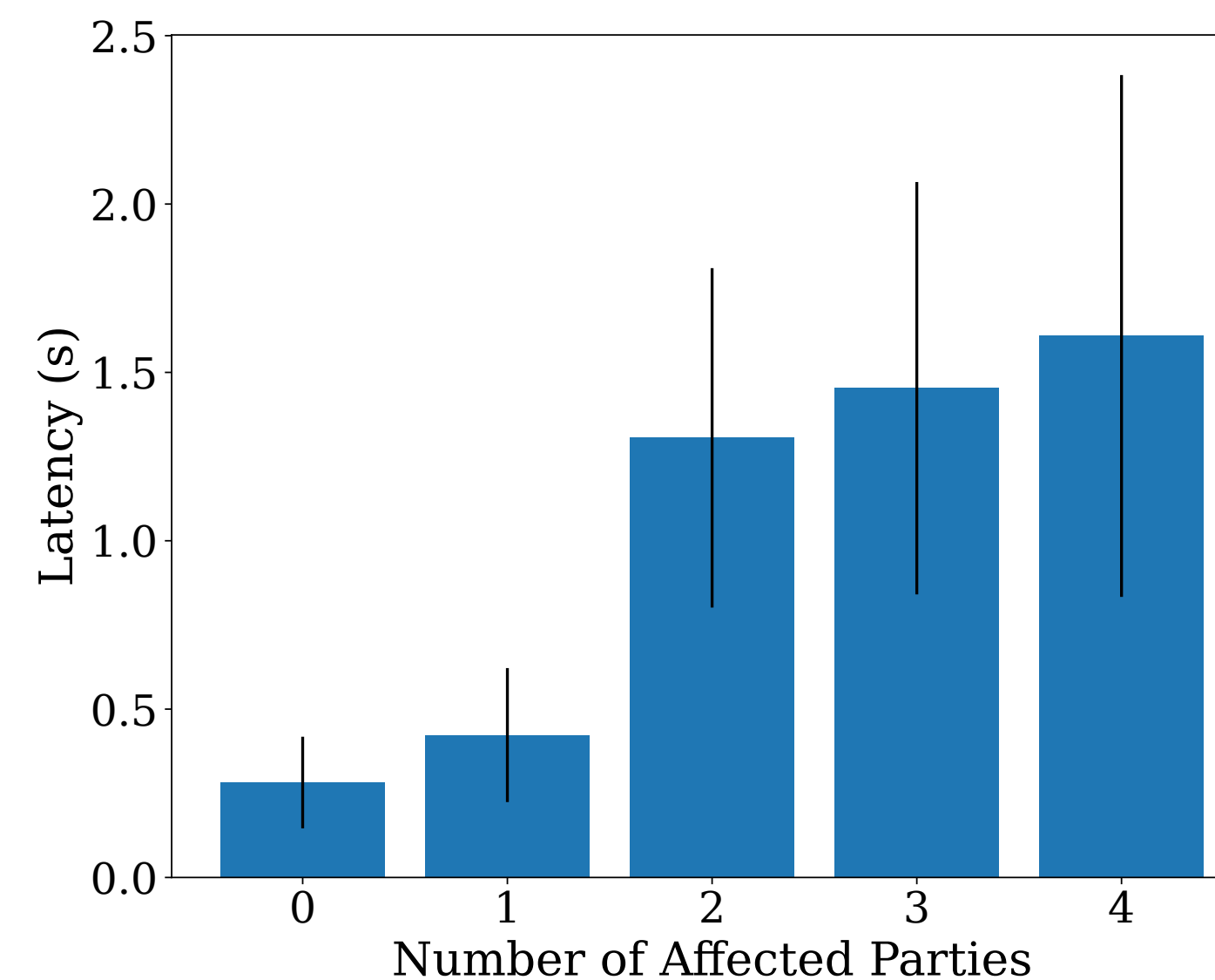


WAN

Conclusion



LAN



WAN

- Open problems
 - Perfect HMPC protocol for general case
 - Bridging the gap between synchronous and asynchronous MPC protocols

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Thank You